# Understanding Hard Disks and File Systems

Module 03

Designed by Cyber Crime Investigators. Presented by Professionals.











## Module Objectives





### After successfully completing this module, you will be able to:

- 1 Describe the different types of disk drives and their characteristics
- 2 Understand the physical and logical structure of a hard disk
- 3 Identify the types of hard disk interfaces and discuss the various hard disk components
- 4 Describe hard disk partitions
- 5 Summarize Windows, Mac, and Linux boot Processes
- 6 Understand various Windows, Linux and Mac OS X file systems
- 7 Differentiate between various RAID storage systems
- 8 Demonstrate file system analysis

## Disk Drive Overview



## HDD

#### **Hard Disk Drive (HDD)**

- The HDD is a non-volatile, random access digital data storage device used in any computer system
- It utilizes a mechanism that reads data from a disk and writes onto an another disk
- The hard disk record data magnetically

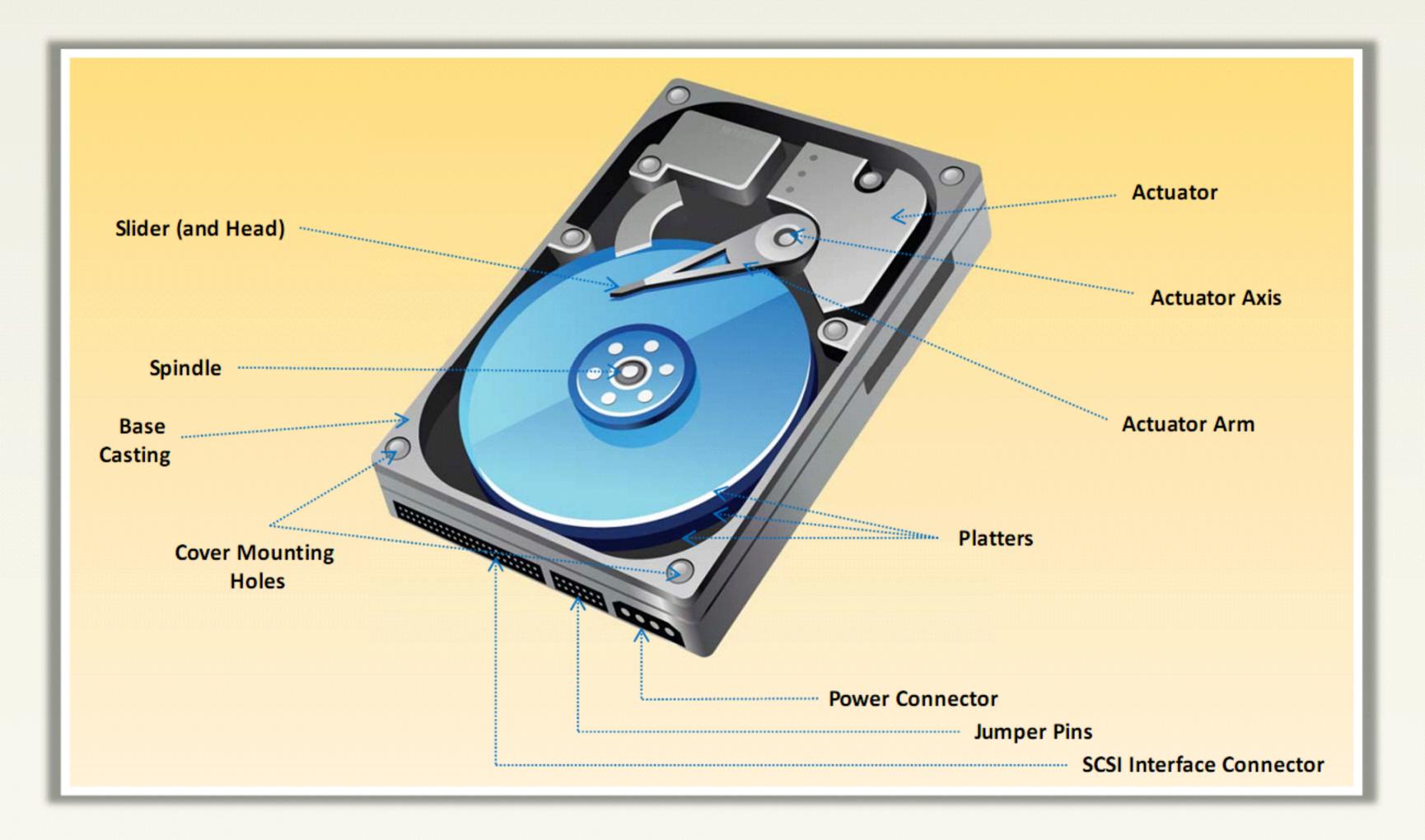
## SSD

#### Solid-state Drive (SSD)

- The SSD is a data storage device that uses solid-state memory to store data and provides access to the stored data in the same manner as a HDD
- It uses microchips to hold data in non-volatile memory chips and does not contain any moving parts
- It is very expensive per gigabyte (GB) and supports a restricted number of writes over the life of the device
- It uses two memories:
  - NAND-based flash memory: It retains memory even without power
  - Volatile RAM: It provides faster access

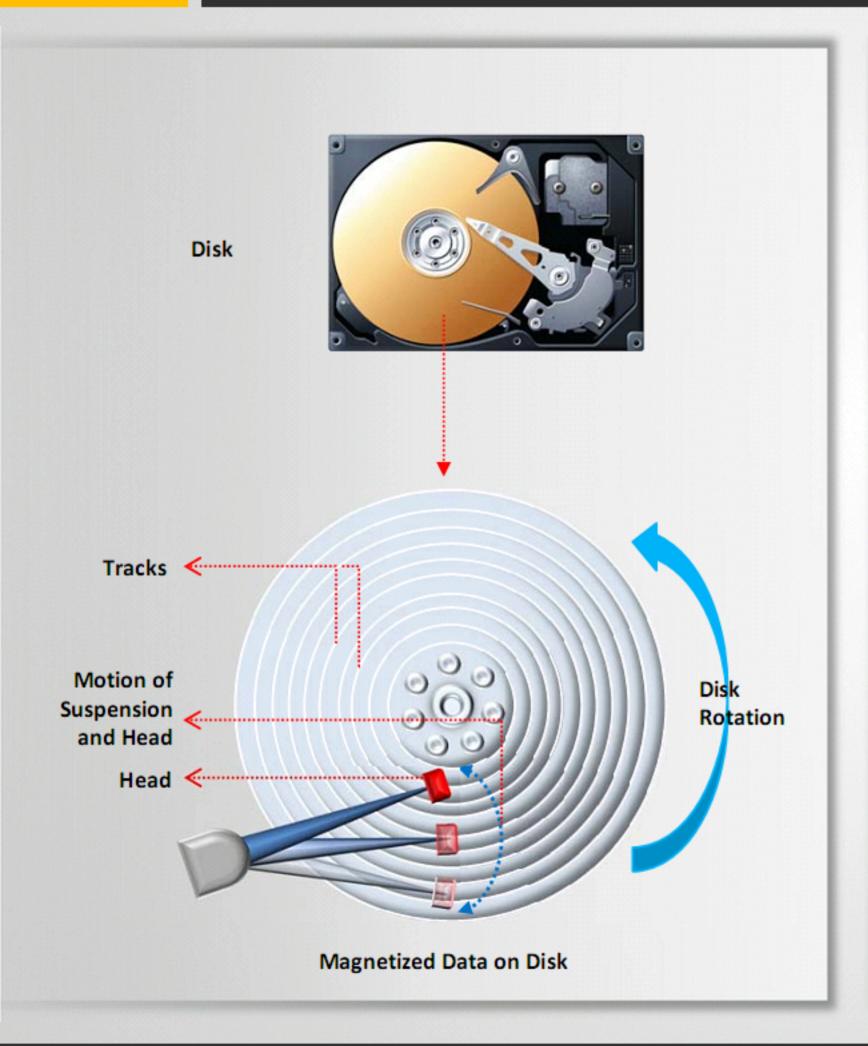
# Physical Structure of a Hard Disk

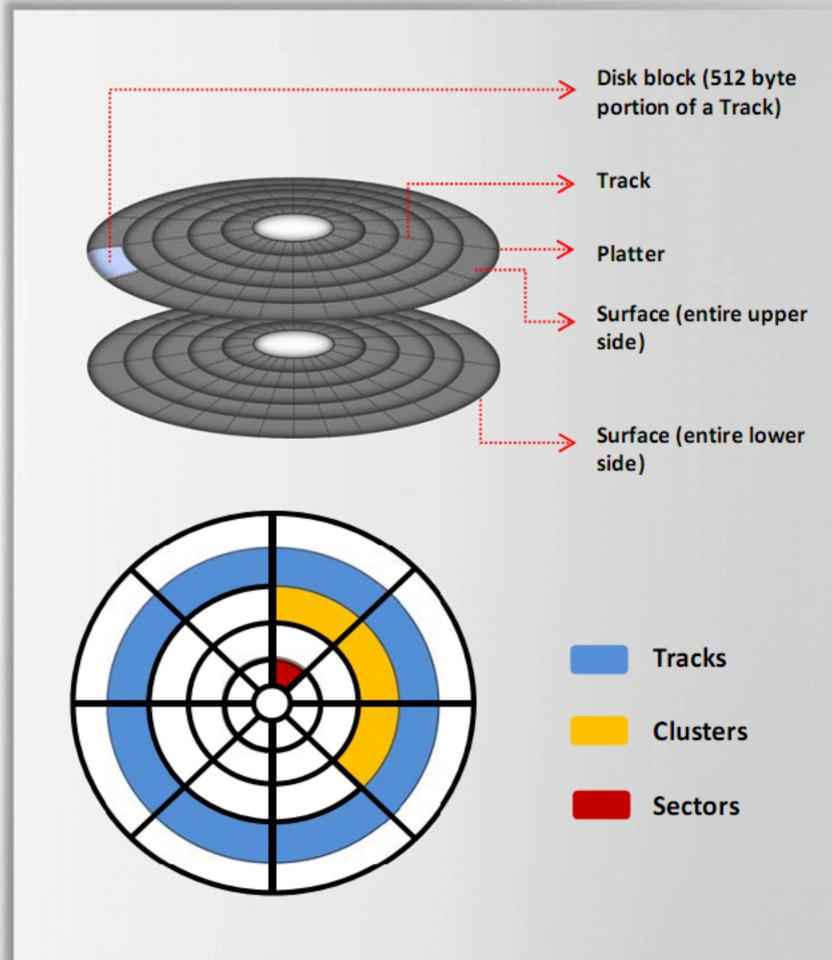




# Physical Structure of a Hard Disk (Cont'd)







## Logical Structure of Hard Disk





The logical structure of a hard disk is the file system and software utilized to control access to the storage on the disk



The hard disk logical structure has significant influence on the performance, consistency, expandability, and compatibility of the storage subsystem of the hard disk



Different operating systems have different file systems and use various ways of arranging and controlling access to data on the hard disk

## **Hard Disk Interfaces**





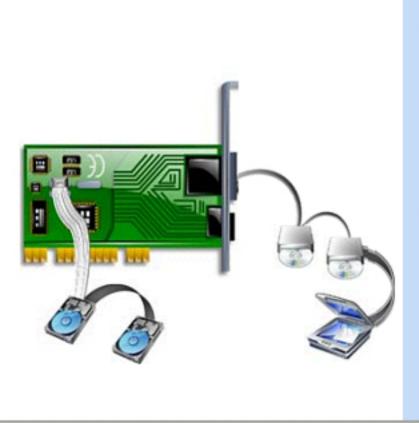
#### ATA/PATA (IDE/EIDE)

ATA (Advanced Technology Attachment) is the official ANSI name of Integrated Drive Electronics (IDE), a standard interface between a motherboard's data bus and storage discs



#### Serial ATA (SATA)

It is an advancement of ATA and uses serial signaling unlike IDE's parallel signaling



#### SCSI

SCSI (Small Computer System Interface) refers to a set of ANSI standard interfaces, based on the parallel bus structure and designed to connect multiple peripherals to a computer

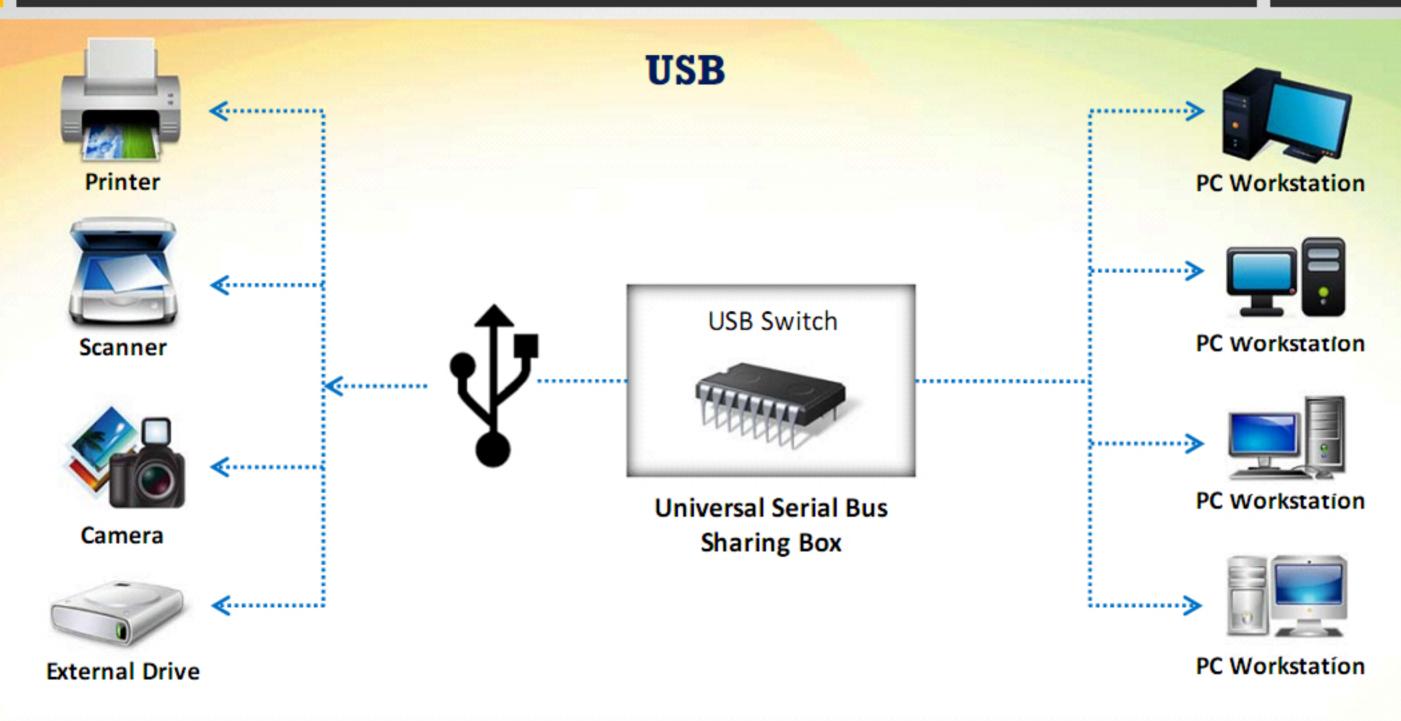


#### **Serial Attached SCSI**

SAS is successor and an advanced alternative to parallel SCSI in enterprise environments

## Hard Disk Interfaces (Cont'd)





#### **Fibre Channel**

- Fibre Channel (FC) is a point-to-point bi-directional serial interface that supports up to 4 Gbps data transfer rates between computer devices
- It is particularly suitable for linking computer system servers to shared storage devices and for interconnecting storage controllers and disk drives

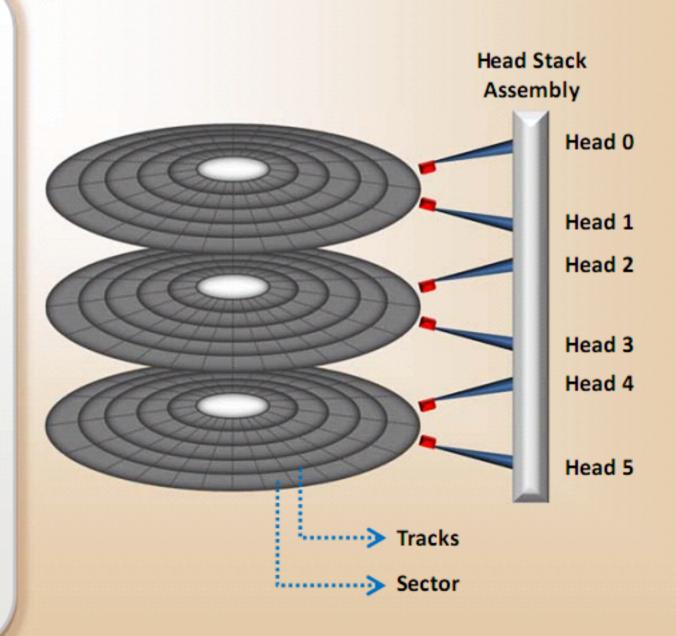
## Tracks



- Tracks are the concentric circles on platters where all the information is stored
- Drive head can access these circular rings in one position at a time
- Tracks are numbered for identification purposes
- Read-write is done by rolling headers from inner to outermost part of the disk

#### **Track Numbering:**

- Track numbering on a hard disk begins at 0 from the outer edge and moves towards the center, typically reaching a value of 1023
- The read/write heads on both surfaces of a platter are tightly packed and locked together on an assembly of head arms
- The arms move in and out together to physically locate all heads at the same track number
- Therefore, a track location is often referred by a cylinder number rather than a track number
- A cylinder is a group of all tracks that start at the same head position on the disk



## Sector



- A sector is the smallest physical storage unit on the disk platter
- It is almost always 512 bytes in size and a few additional bytes for drive control and error correction
- Each disk sector is labelled using the factory track-positioning data
- The optimal method of storing a file on a disk is in a contiguous series
- For example, if the file size is 600 bytes, two 512 bytes sectors are allocated for the file

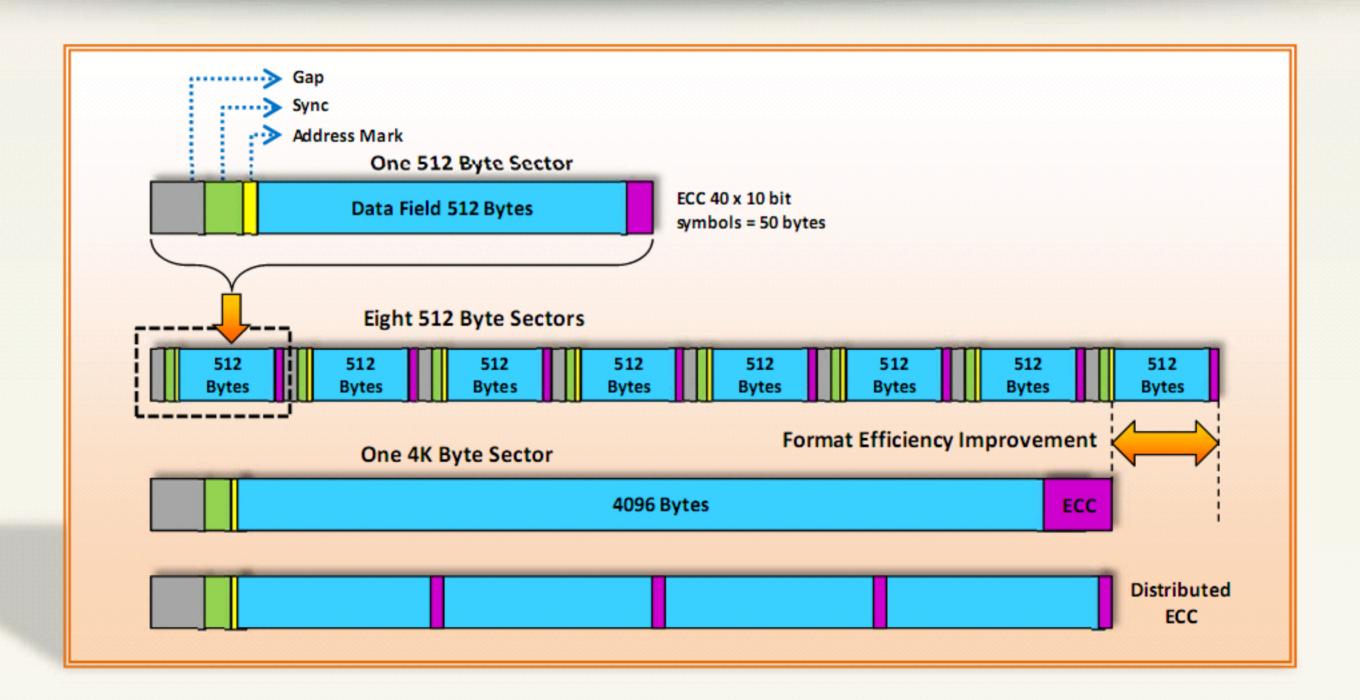
#### **Sector Addressing:**

- Cylinders, heads and sectors (CHS) determine the address of the individual sectors on the disk
- For example, on formatting a disk, 50 tracks are divided into 10 sectors each
- Track and sector numbers are used by the operating system and disk drive to identify the stored information

## Advanced Format: Sectors



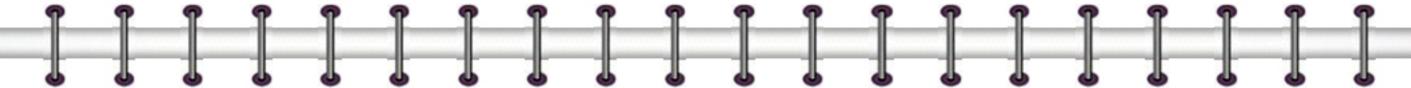
- New hard drives use 4096 byte (4 KB or 4K) advanced format sectors
- Generation-one Advanced Format also called as 4K sector technology, efficiently uses the storage surface media of a disk efficiently by merging eight 512 byte sectors into one single sector (4096 bytes)
- After merging, the structure of 4K sector does not disturb the key design elements of the traditional 512-byte sector



## Clusters



- A cluster is the smallest logical storage unit on a hard disk. It is a set of track sectors, ranging from 2 to 32 or more, depending on the formatting scheme in use
- The file system divides the storage on a disk volume into discreet chunks of data for efficient disk usage and performance. These chunks are called clusters
- The process by which files are allocated to clusters is called allocation, so clusters are also known as allocation units
- In the File Allocation Table (FAT) file system, the clusters linked with a file keep track of file data in the hard disk's file allocation table



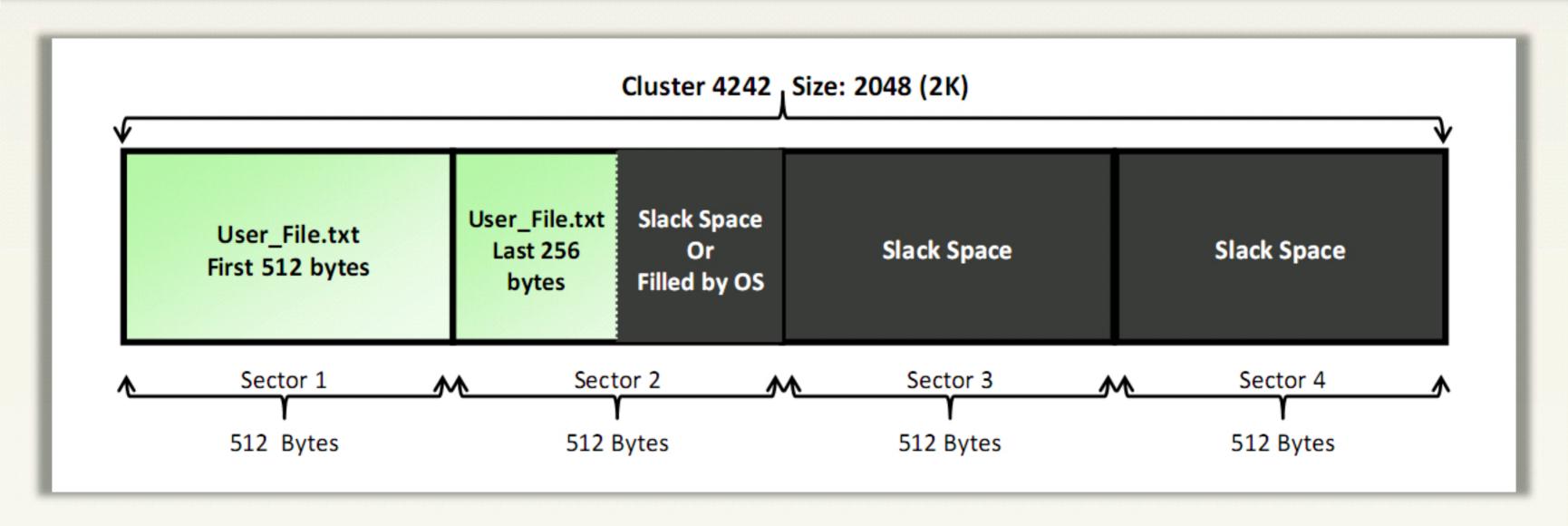
#### **Cluster Size**

- Cluster sizing has a significant impact on the performance of an operating system and disk utilization
- Cluster size can be altered for optimum disk storage
- The size of a cluster depends on the size of the disk partition and type of file system installed on the partition
- Larger cluster size (greater than one sector):
  - Minimizes the fragmentation problem
  - Increases the probability of unused space in the cluster
  - Reduces disk storage area to save information
  - Reduces the unused area on the disk

## Slack Space



- Slack space is the area of a disk cluster between the end of the file and the end of the cluster
- If the file size is less than the cluster size, still a full cluster is assigned to that file. The remaining space remains unused and is called slack space. This remaining unused space is called slack space.
- For example, if the partition size is 4 GB, each cluster will be 32 KB. Even if a file requires only 10 KB, the entire 32 KB will be allocated to that file, resulting in 22 KB of slack space



## Lost Clusters



- When the operating system marks clusters, as used, but does not allocate them to any file, such clusters are known as lost clusters
- A lost cluster is a FAT file system error that results from in what manner the FAT file system allocates space and chains files together
- It is mainly the result of a logical structure error and not a physical disk error
- They usually occur because of interrupted file activities such as, 'the file is not correctly completed and closed' thus, the clusters have involved never linked correctly to a file
- CHKDSK is a system tool in Windows, that authenticates the file system reliability of a volume and repairs logical file system errors

```
Administrator: Command Prompt
           the file system is NTFS.
 ARNING! /F parameter not specified.
  nning CHKDSK in read-only mode.
 tage 1: Examining basic file system structure ...
                                                            84224 file records processed.
 ile verification completed.
                                                            2 large file records processed.
                                                            0 bad file records processed.
Stage 2: Examining file name linkage ...
                                                            91752 index entries processed.
 ndex verification completed.
                                                            @ unindexed files scanned.
                                                            0 unindexed files recovered to lost and found.
Stage 3: Examining security descriptors ...
 ecurity descriptor verification completed.
                                                            3765 data files processed.
CHKDSK is verifying Usn Journal...
                                                            157066512 USN bytes processed.
 sn Journal verification completed.
 indows has scanned the file system and found no problems.
  further action is required.
 204799999 KB total disk space.
 142636808 KB in 77667 files.
    34736 KB in 3766 indexes.
        0 KB in bad sectors.
   310243 KB in use by the system.
    65536 KB occupied by the log file.
  61818212 KB available on disk.
     4096 bytes in each allocation unit.
  51199999 total allocation units on disk.
  15454553 allocation units available on disk.
```

## **Bad Sectors**



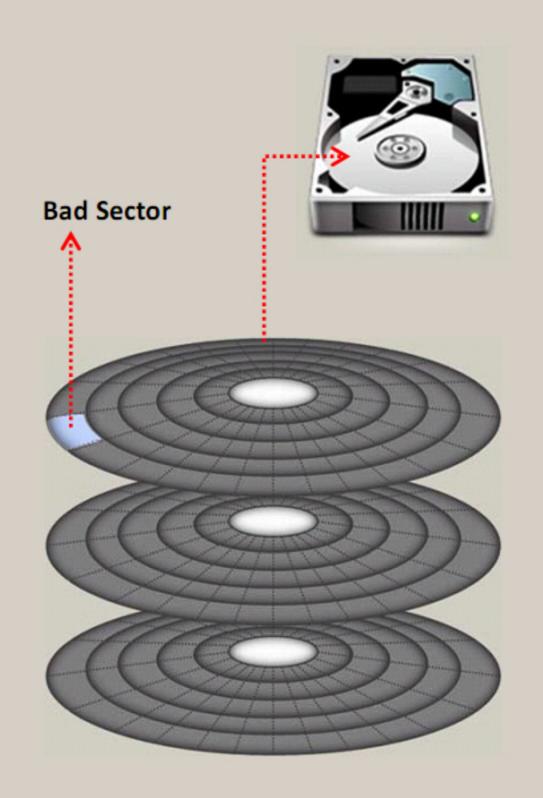
Bad sector is a

damaged portion of
a disk on which no
read/write operation
can be performed

Formatting a disk enables the operating system to identify unusable sectors and mark them as bad

Bad sectors are formed due to configuration problems or any physical disturbances to the disk

If data is in a sector that becomes bad, then it might not be recoverable Data can be recovered using software tools such as Chkdsk



## Understanding Bit, Nibble and Byte



#### Bit:

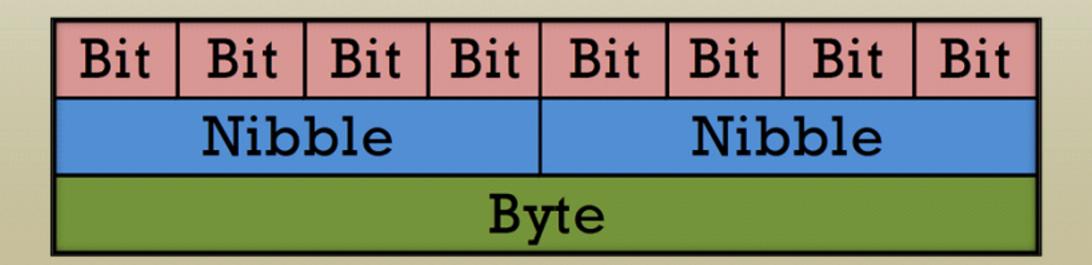
- Short for binary digit
- It is the smallest unit of data stored in a computer and is represented as a binary value, either 1 (true) or 0 (false)

#### **Nibble:**

- It is a group of 4 bits and is half the size of a Byte
- Not a common term as most microprocessors use group of 8 bits or higher to process data

#### **Byte:**

- It is a group of 8 bits and twice the size of a Nibble
- One single character typed from a keyboard takes one byte of storage

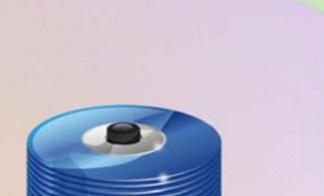


## Hard Disk Data Addressing



Hard disk data addressing is a method of allotting addresses to each physical block of data on a hard disk





Hard Disk Data
Addressing
Methods





CHS (Cylinder-Head-Sector)

- It addresses data by simply specifying the cylinder (radius), head (platter side), and sector (angular position)
- It is used on most IDE drives

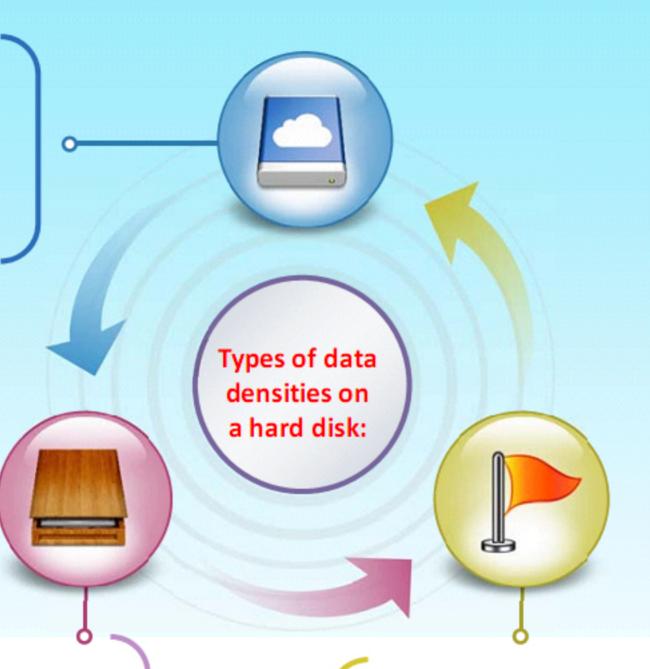
- LBA (Logical Block Address)
- It addresses data by allotting a sequential number to each sector of the hard disk
- It is used on SCSI and enhanced IDE drives

## Data Densities on a Hard Disk



#### **Track Density**

It is defined as the space between tracks on a disk



- Data is recorded onto a hard disk using a method called zoned bit recording (also known as a multiple zone recording)
- In this technique, tracks are combined together into zones depending on their distance from the center of the disk
- Each zone is assigned a number of sectors per track

#### **Areal Density**

It is defined as the number of bits per square inch on a platter

#### **Bit Density**

It is the bits per unit length of track



## Disk Capacity Calculation



## Disk Capacity Calculation Question?



A disk drive has 16,384 cylinders, 80 heads, and 63 sectors per track. Assume - a sector has 512 bytes. What is the capacity of such a disk?

#### **Answer**

The conversion factors appropriate to this hard disk are:

- 16,384 cylinders / disk
- 80 heads / cylinder
- 63 sectors / track
- 512 bytes / sector

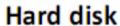
Total bytes = 1 disk \* (16,384 cylinders / disk) \* (80 heads / cylinder) \* (1 track / head) \* (63 sectors / track) \* (512 bytes / sector)

= 42,278,584,320 bytes

## Measuring the Hard Disk Performance









Data stored as files



**CPU** 



Data is stored on the hard disk in the form of files



When running program requests the file, hard disk recovers the byte content of the file and sends them to the CPU one at a time for further processing

Hard disk performance is measured by these factors:

- Data rate: It is a ratio of the number of bytes per second that hard disk sends to the CPU
- Seek time: It is the amount of time required to send the first byte of the file to the CPU, when it requests the file

## Disk Partitions



The HDD partitioning is the creation of logical divisions upon a hard disk that allows user to apply operating system-specific logical formatting

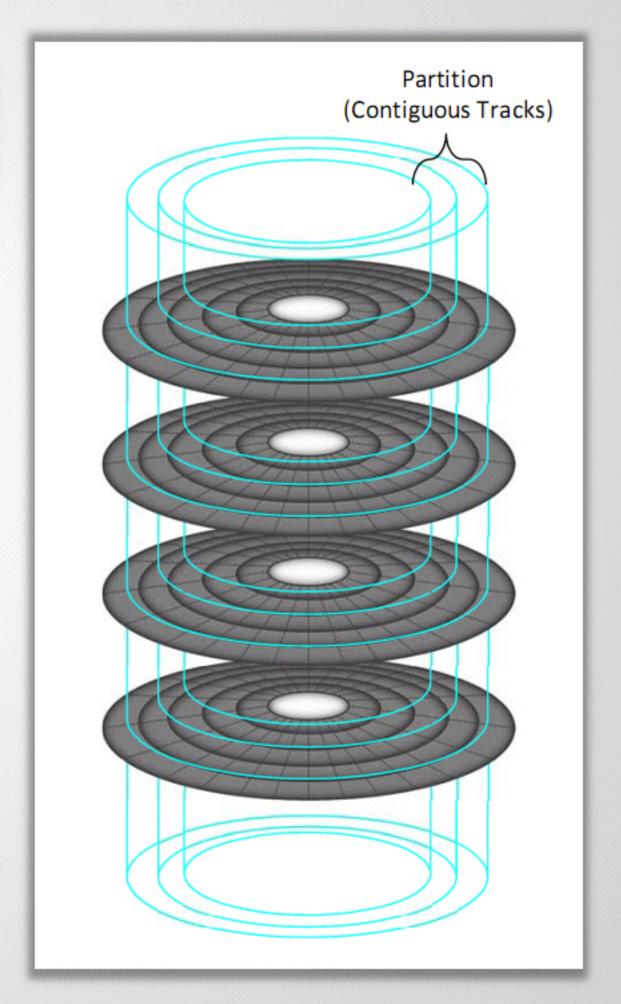
## **Primary Partition**

- It is a drive that holds the information regarding operating system, system area, and other information required for booting
- In MS-DOS and earlier versions of Microsoft Windows systems, the first partition (C:) must be a "primary partition"

## **Extended Partition**

It is the logical drive that holds the information regarding stored data and files in the disk





## BIOS Parameter Block (BPB)



- The BIOS parameter block (BPB) is a data structure in the partition boot sector
- It describes the physical layout of a data storage volume, like the number of heads and the size of the tracks on the drive
- BPB in file systems such as FAT12 (except for in DOS 1.x), FAT16, FAT32, HPFS, and NTFS defines the filesystem structure
- The BPB length varies for FAT16, FAT32, and NTFS boot sectors, due to different types of fields and the amount of data stored in them
- BPB assists investigators to locate the file table on the hard drive

Format of full DOS 7.1 Extended BIOS Parameter Block (79 bytes) for FAT32:			
Sector offset	BPB offset	Field length	Description
0x00B	0x00	25 BYTEs	DOS 3.31 BPB
0x024	0x19	DWORD	Logical sectors per FAT
0x028	0x1D	WORD	Mirroring flags etc.
0x02A	0x1F	WORD	Version
0x02C	0x21	DWORD	Root directory cluster
0x030	0x25	WORD	Location of FS Information Sector
0x032	0x27	WORD	Location of backup sector(s)
0x034	0x29	12 BYTEs	Reserved (Boot file name)
0x040	0x35	BYTE	Physical drive number
0x041	0x36	BYTE	Flags etc.
0x042	0x37	BYTE	Extended boot signature (0x29)
0x043	0x38	DWORD	Volume serial number
0x047	0x3C	11 BYTEs	Volume label
0x052	0x47	8 BYTEs	File-system type

NTFS - Format of Extended BPB for NTFS (73 bytes):			
Sector offset	BPB offset	Field length	Description
0x00B	0x00	25 BYTEs	DOS 3.31 BPB
0x024	0x19	BYTE	Physical drive number (identical to DOS 3.4 EBPB)
0x025	0x1A	BYTE	Flags etc. (identical to DOS 3.4 EBPB)
0x026	0x1B	BYTE	Extended boot signature (0x80 aka "8.0") (similar to DOS 3.4 EBPB and DOS 4.0 EBPB)
0x027	0x1C	BYTE	Reserved
0x028	0x1D	QWORD	Sectors in volume
0x030	0x25	QWORD	MFT first cluster number
0x038	0x2D	QWORD	MFT mirror first cluster number
0x040	0x35	DWORD	MFT record size
0x044	0x39	DWORD	Index block size
0x048	0x3D	QWORD	Volume serial number
0x050	0x45	DWORD	Checksum

## **Master Boot Record (MBR)**



A master boot record (MBR) is the first sector ("sector zero") of a data storage device, such as a hard disk



The information regarding the files on the disk, their location, size, and other important data is stored in the MBR file



Ш

In practice, MBR almost always refers to the 512-byte boot sector or partition sector of a disk



IV

MBR is used for:

- Holding a partition table which refers to the partitions of a hard disk
- Bootstrapping an operating system
- Distinctively recognizing individual hard disk media with a 32-bit disk signature



## Structure of a Master Boot Record



Backing up the MBR

In UNIX/Linux, dd can be used to backup and restore the MBR

Backup the MBR

dd if=/dev/xxx of=mbr.backup bs=512 count=1

Restore the MBR

dd if=mbr.backup of=/dev/xxx bs=512 count=1

Address		Description		Size in bytes	
Hex	Oct	Dec			
0000	0000	0	Code Area		440 (max. 446)
01B8	0670	440	Disk Signatu	re (Optional)	4
01BC	0674	444	Usually Nulls; 0x0000		2
01BE	0676	446	Table of Primary Partitions (Four 16-byte entries, IBM partition table scheme)		64
01FE	0776	510	55h	MBR Signature;	2
01FF	0777	511	AAh	0xAA55	2
MBR, Total Size: 446 + 64 = 2 =			512		



1010010 0010110

# Structure of a Master Boot Record (Cont'd)



#### **Layout of 16-byte Partition Record**

Offset	Description
0x00	Status (0x80 = bootable, 0x00 = non-bootable, other = malformed)
0x01	Cylinder-head-sector address of the first sector in the partition
0x04	Partition type
0x05	Cylinder-head-sector address of the last sector in the partition
0x08	(4 bytes) Logical block address of the first sector in the partition
0x0C	(4 bytes) Length of the partition, in sectors



#### **Layout of IBM Extended Partition Record**

Offset	Description
0x00	Status bits (bit 0 = list on Boot Manager menu, other bits = reserved)
0x01	Space-padded partition name



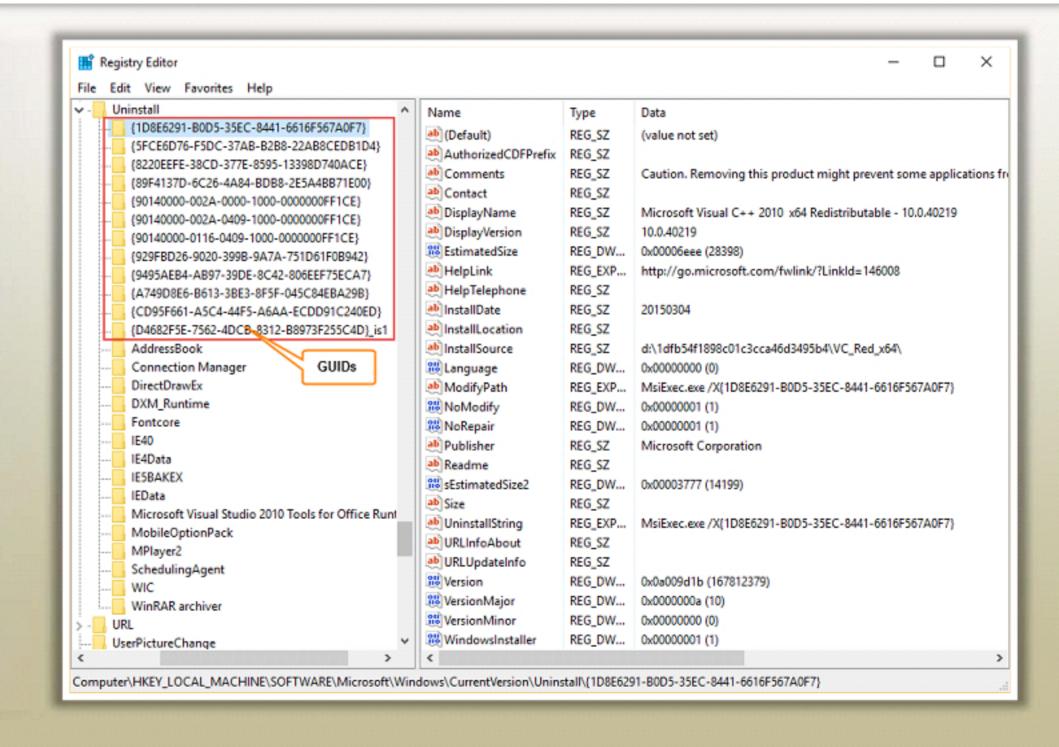
## Globally Unique Identifier (GUID)



- Global Unique Identifier (GUID) is a 128-bit unique reference number used as an identifier in computer software
- In general, GUIDs are displayed as 32 hexadecimal digits with groups separated by hyphens

#### **Common Uses:**

- In Windows registry, GUIDs are used to identify COM DLLs
- In database tables, GUIDs are used as primary key values
- Website assigns GUID to a user's browser to record and track the session
- Windows assigns GUID to a username to identify user accounts



## GUID Partition Table (GPT)



- Unified Extensible Firmware Interface (UEFI) replaces legacy BIOS firmware interfaces
- UEFI is a specification that defines a software interface between an OS and platform firmware
- It uses a partition system known as GUID Partition Table (GPT) that replaces the traditional MBR

#### Advantages of GPT disk layout:

- Supports up to 128 partitions and uses 64-bit Logical Block Addresses (LBAs)
- Supports maximum partition size from 2 Tebibyte (TiB) to 8 Zebibyte (ZiB)
- Provides primary and backup partition tables for redundancy

http://www.invoke-ir.com

### **GUID Partition Table Scheme** LBA 0 Protective MBR LBA 1 Primary GPT Header LBA 2 Entry 2 Entry 1 Entry 3 LBA 3 Entries 5-128 **LBA 34** Partition 1 Partition 2 **Remaining Partitions** LBA -34 Entry 1 Entry 2 Entry 3 Entry 4 LBA -33 Entries 5-128 LBA -2 Secondary GPT Header LBA -1

# GUID Partition Table (GPT) (Cont'd)



#### **Protective MBR:**

- Disk formatted with a GPT disk layout has a Protective MBR located at Logical Block Address (LBA) 0
- Protective MBR provides compatibility with legacy tools that fail to understand the GPT format
- It is alike to the "legacy" MBR in functionality, but has only one partition of type 0xEE (EFI\_GPT\_DISK)
- This partition reserves the entire disk for the formal GUID Partition Table structure

**Note:** The **UEFI Firmware** does not execute the MBR Boot Code (the first 440 bytes)

The Get-MBR cmdlet displays the MBR Partition Table of a GPT formatted disk

http://www.invoke-ir.com

#### **PROTECTIVE MBR**

First sector of drive For breakdown see MBR poster

```
000 33 CO 8E DO BC 00 7C 8E CO 8E D8 BE 00 7C BF 00
010 06 B9 00 02 FC F3 A4 50 68 1C 06 CB FB B9 04 00
020 BD BE 07 80 7E 00 00 7C 0B 0F 85 0E 01 83 C5 10
030 E2 F1 CD 18 88 56 00 55 C6 46 11 05 C6 46 10 00
040 B4 41 BB AA 55 CD 13 5D 72 OF 81 FB 55 AA 75 09
050 F7 C1 01 00 74 03 FE 46 10 66 60 80 7E 10 00 74
   26 66 68 00 00 00 00 66 FF 76 08 68 00 00 68 00
   7C 68 01 00 68 10 00 B4 42 8A 56 00 8B F4
   9F 83 C4 10 9E EB 14 B8 01 02 BB 00 7C 8A 56 00
   8A 76 01 8A 4E 02 8A 6E 03 CD 13 66 61 73 1C FE
   4E 11 75 OC 80 7E 00 80 OF 84 8A 00 B2 80 EB 84
   55 32 E4 8A 56 00 CD 13 5D EB 9E 81 3E FE
   AA 75 6E FF 76 00 E8 8D 00 75 17 FA B0 D1
ODO E8 83 00 BO DF E6 60 E8 7C 00 BO FF E6 64 E8 75
   00 FB B8 00 BB CD 1A 66 23 CO 75 3B 66
   43 50 41 75 32 81 F9 02 01 72 2C 66 68 07
   00 66 68 00 02 00 00 66 68 08 00 00 00 66 53 66
110 53 66 55 66 68 00 00 00 00 66 68 00 7c 00 00 66
   61 68 00 00 07 CD 1A 5A 32 F6 EA 00 7C 00 00 CD
   18 AO B7 O7 EB O8 AO B6 O7 EB O3 AO B5 O7 32 E4
   05 00 07 8B FO AC 3C 00 74 09 BB 07 00 B4 0E CD
   10 EB F2 F4 EB FD 2B C9 E4 64 EB 00 24 02 EO F8
   24 02 C3 49 6E 76 61 6C 69 64 20 70 61
   74 69 6F 6E 20 74 61 62 6C 65 00 45 72 72
   20 6C 6F 61 64 69 6E 67 20 6F 70 65 72 61 74 69
   6E 67 20 73 79 73 74 65 6D 00 4D 69 73 73 69 6E
1AO 67 20 6F 70 65 72 61 74 69 6E 67 20 73 79 73 74
   65 6D 00 00 00 63 7B 9A 00 00 00 00 00 00 00 00
1CO 02 00 EE FF FF FF 01 00 00 00 FF FF FF FF 00 00
```

#### **IMPORTANT PROTECTIVE MBR VALUES**

System id EE – EFI GPT partition GPT header sector offset 1

# GUID Partition Table (GPT) (Cont'd)



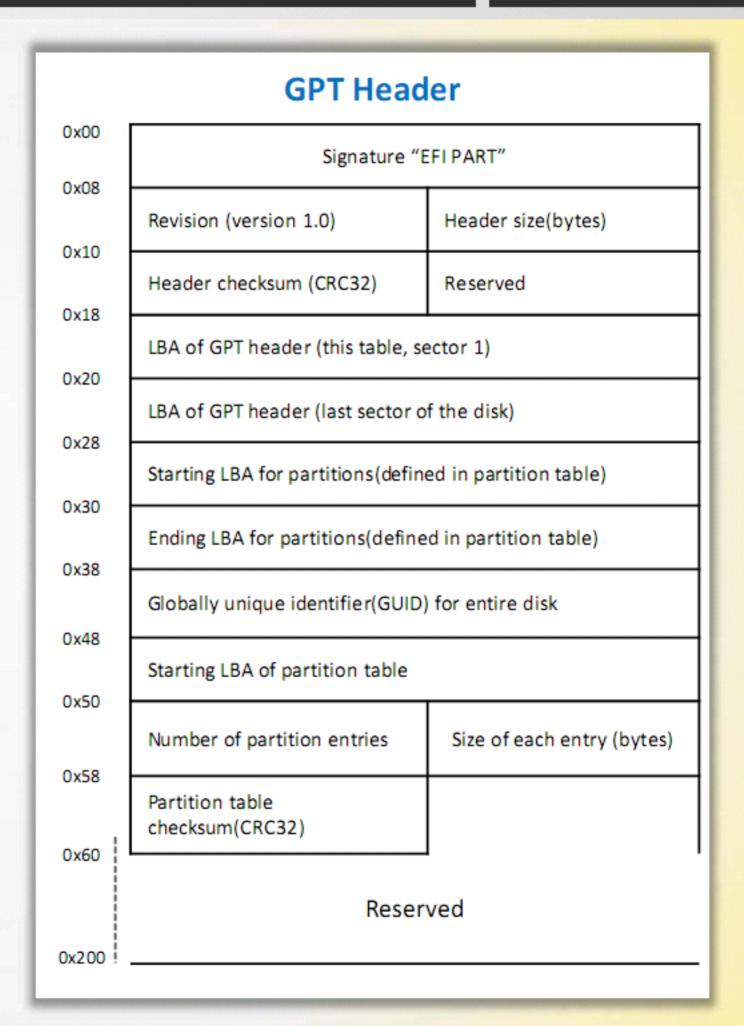
#### **GUID Partition Table:**

The formal GUID Partition Table starts at LBA 1 where the GPT Header is found

#### **GPT Header:**

- It is a pointer to the partition table and defines the complete logical layout
- Contains information such as the "EFI PART" signature and a unique GUID of the disk
- The firmware detects GPT corruption using the CRC32 values
- The MyLBA value, always 1, defines the location of GPT header, while the AlternateLBA value represents the backup GPT and will occupy the last sector on a disk
- The backup GPT replaces the original GPT when it is corrupted
- The FirstUsableLBA and LastUsableLBA values point to the disk portion the partitions can use
- The PartitionEntryLBA value represents start of the array, while the NumberOfPartitionEntries and SizeOfPartitionEntry values denote the overall partition size

http://www.invoke-ir.com



## GUID Partition Table (GPT) (Cont'd)



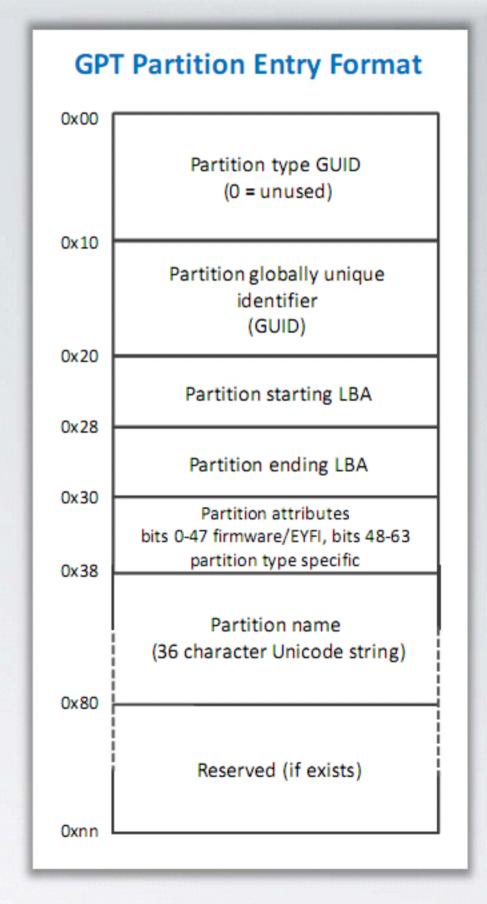
#### **GPT Partition Array:**

- The GPT Header points to the Partition Array via the PartitionEntryLBA value
- The size and number of partitions are defined in the GPT Header

#### **Each partition contains:**

- Two GUIDs, one represents the type of partition, and the second uniquely identifies the partition
- Partition StartingLBA and EndingLBA values describing the location and size of the partition
- Partition type specific attributes
- 36 character user-defined partition name

http://www.invoke-ir.com



## What is the Booting Process?



Booting refers to the process of starting or resetting operating systems when the user turns on a computer system



operating system (stored in the hard disk) to the RAM (working memory)

### **Types of Booting**



- Cold boot (Hard boot)
  - It is the process of starting a computer from a powered-down or off state
- Warm boot (Soft boot)

It is the process of restarting a computer that is already turned on through the operating system



## Essential Windows System Files

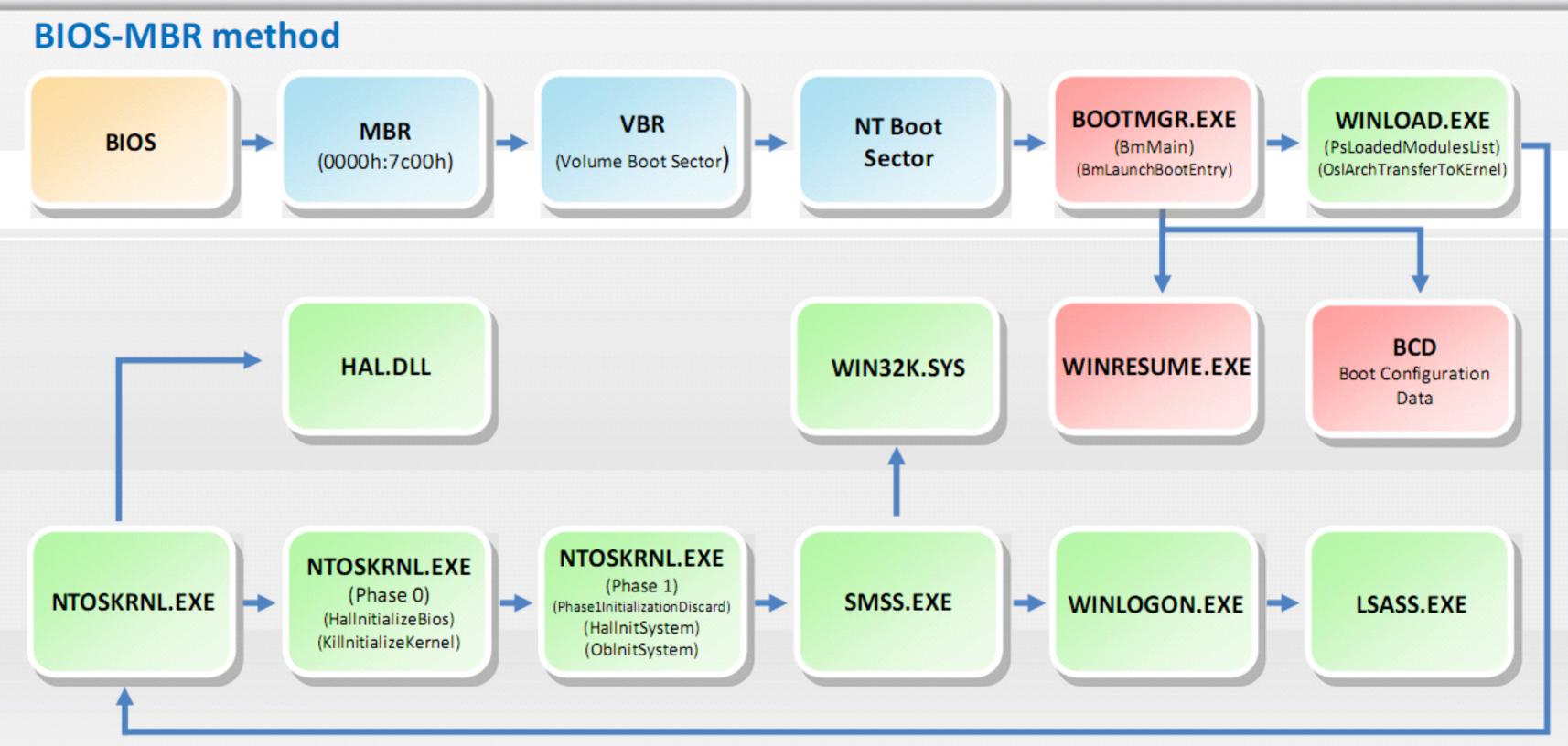


File Names	Description	
Ntoskrnl.exe	Executive and kernel	
Ntkrnlpa.exe	Executive and kernel with support for Physical Address Extension (PAE)	
Hal.dll	Hardware abstraction layer	
Win32k.sys	Kernel-mode part of the Win32 subsystem	
Ntdll.dll	Internal support functions and system service dispatch stubs to executive functions	
Kernel32.dll		
Advapi32.dll	Win22 subsystem DII files	
User32.dll	Win32 subsystem DLL files	
Gdi32.dll		

## Windows Boot Process



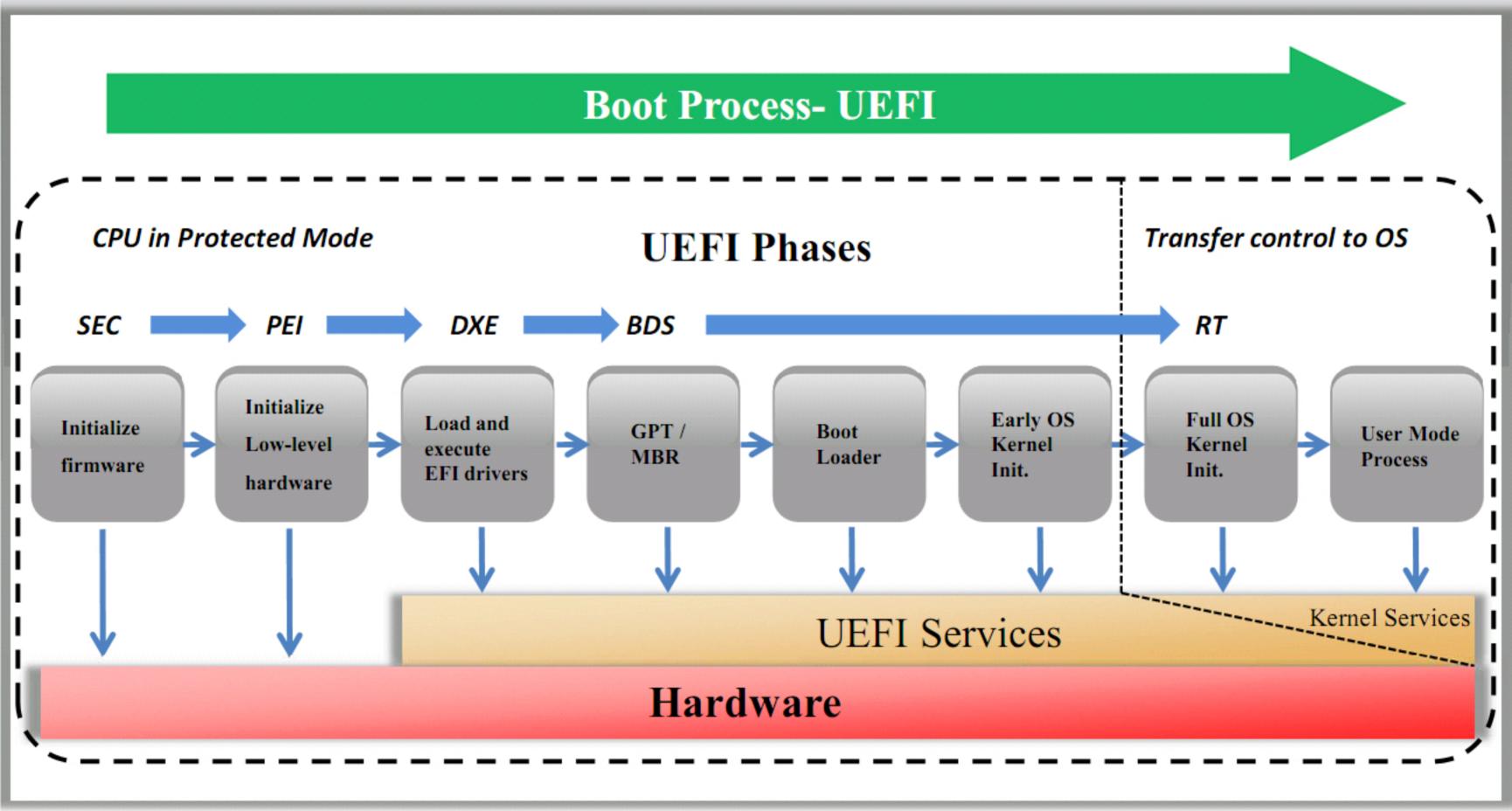
Windows XP, Vista, and 7 OSs power on and start up using the traditional BIOS-MBR method. Whereas, OSs from Windows 8 and above uses either traditional BIOS-MBR method or newer UEFI-GPT method according to the user choice



http://breaking-the-system.blogspot.in

## Windows Boot Process (Cont'd)





http://answers.microsoft.com

# Identifying GUID Partition Table (GPT)



Investigators can use cmdlets given below in Windows PowerShell to identify the presence of GPT:

#### Get-GPT

- It parses the GPT data structure contained within the first few sectors of the device specified
- It requires the use of the -Path parameter which takes the Win32 Device Namespace (ex. \\.\PHYSICALDRIVE1) for the device from which the GPT should be parsed

```
PS C:\Windows\system32> Get-GPT -Path \\.\PHYSICALDRIVE1

Revision : 1.0
HeaderSize : 92
MyLBA : 1
AlternateLBA : 20971519
FirstUsableLBA : 34
LastUsableLBA : 34
LastUsableLBA : 20971486
DiskGUID : f913e110-0835-4cf1-96c7-380b5db4a42d
PartitionEntryLBA : 2
NumberOfPartitionEntries : 128
SizeOfPartitionEntry : 128
PartitionTable : {Microsoft reserved partition, Basic data partition}
```

If Get-GPT is run against a disk formatted with a MBR, it will throw an error prompting to use Get-MBR instead

```
PS C:\Windows\system32> Get-GPT -Path \\.\PHYSICALDRIVEO

Get-GPT : No GPT found. Please use Get-MER cmdlet

At line:1 char:1
+ Get-GPT -Path \\.\PHYSICALDRIVEO
+ CategoryInfo : NotSpecified: (:) [Get-GPT], Exception
+ FullyQualifiedErrorId : System.Exception,InvokeIR.PowerForensics.Cmdlets.GetGPTCommand
```

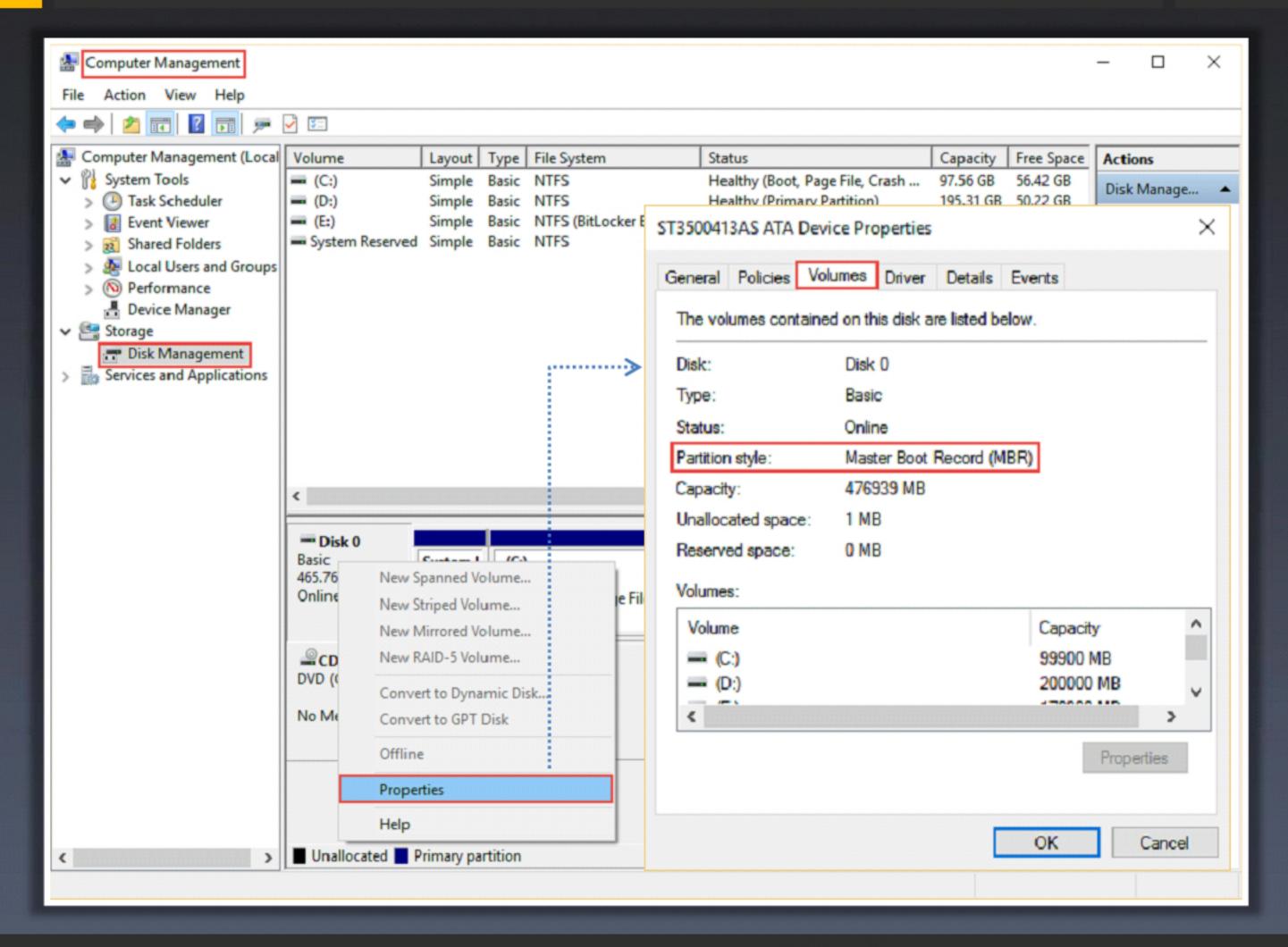
#### Alternate Method:

- Open "Computer Management" application and click "Disk Management" on the left pane. Right-click on the primary disk (here, Disk 0) and then click Properties.
- In the Device Properties window, click "Volumes" tab to see the Partition style

http://www.invoke-ir.com

# Identifying GUID Partition Table (GPT) (Cont'd)





# Identifying GUID Partition Table (GPT) (Cont'd)



#### **Get-BootSector**

It reviews the hard drive's first sector and determines if the disk is formatted using the MBR or GPT partitioning scheme. Once done, it acts just as Get-MBR or Get-GPT would, respectively.

#### Get-BootSector run against a disk formatted using the GPT partitioning scheme:

```
PS C:\Windows\system32> Get-BootSector -Path \\.\PHYSICALDRIVE1
Revision
                         : 1.0
                         : 92
HeaderSize
AlternateLBA
                         : 20971519
FirstUsableLBA
                         : 20971486
LastUsableLBA
                         : f913e110-0835-4cf1-96c7-380b5db4a42d
DiskGUID
PartitionEntryLBA
NumberOfPartitionEntries : 128
SizeOfPartitionEntry
                         : 128
                         : {Microsoft reserved partition, Basic data partition, Basic data partition}
PartitionTable
```

#### Get-BootSector run against a disk formatted using the MBR partitioning scheme:

# Identifying GUID Partition Table (GPT) (Cont'd)



#### Get-PartitionTable

It determines the type of boot sector (MBR or GPT) and returns the correct partition object (PartitionEntry or GuidPartitionTableEntry)

Get-PartitionTable run against an MBR formatted disk, returning an PartitionEntry object:

```
PS C:\Windows\system32> Get-PartitionTable -Path \\.\PHYSICALDRIVEO

Bootable SystemID StartSector EndSector

------

True NTFS 2048 125827071
```

Get-PartitionTable run against a GPT formatted disk, returning an array of GuidPartitionTableEntry Objects:

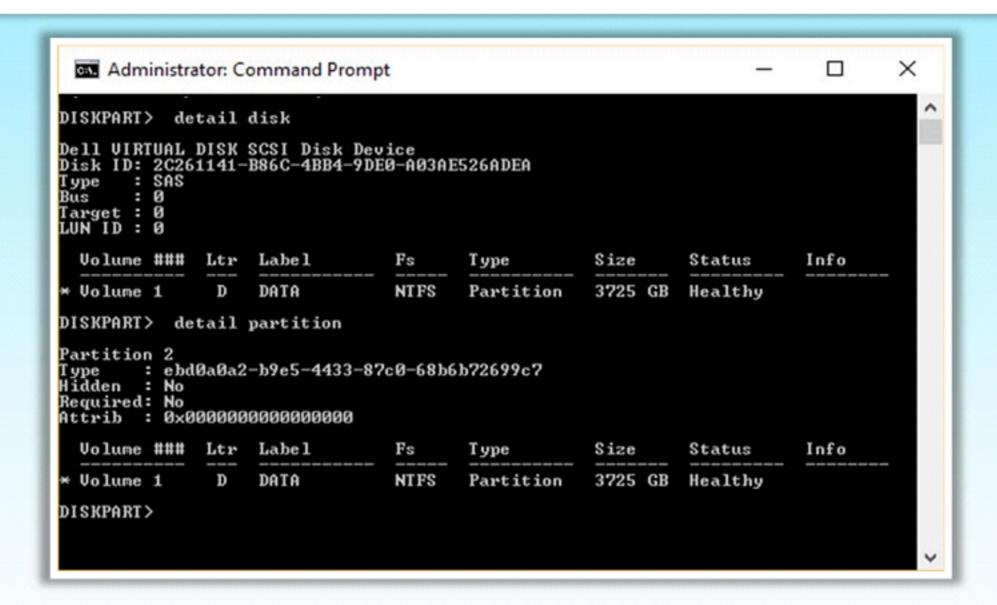
```
PS C:\Windows\system32> Get-PartitionTable -Path \\.\PHYSICALDRIVE1
PartitionTypeGUID : e3c9e316-0b5c-4db8-817d-f92df00215ae
UniquePartitionGUID: ff1a8a47-08f8-43ab-b410-53697f0b2323
StartingLBA
                   : 34
                   : 65569
EndingLBA
Attributes
                   : Microsoft reserved partition
PartitionName
PartitionTypeGUID : ebd0a0a2-b9e5-4433-87c0-68b6b72699c7
UniquePartitionGUID: 6d76ae42-b6c1-4fbe-8d42-20cd366026b4
StartingLBA
                 : 67584
EndingLBA
                   : 2164735
Attributes
                   : 0
PartitionName
                   : Basic data partition
PartitionTypeGUID : ebd0a0a2-b9e5-4433-87c0-68b6b72699c7
UniquePartitionGUID : d6795c3a-8a4d-4fb4-91a0-488812cce027
StartingLBA
                   : 2164736
EndingLBA
                   : 4261887
Attributes
                   : 0
PartitionName
                   : Basic data partition
```

http://www.invoke-ir.com

# Analyzing the GPT Header and Entries



- Most of the operating systems that support GPT disk access come up with a basic partitioning tool, which displays details about GPTs
  - E.g.: DiskPart tool (Windows), OS X Disk utility (Mac), GNU parted tool (Linux)



- Sleuthkit (mmls command) can be used to view detailed partition layout for GPT disk
- Alternatively, details about GPT header and partition entries can be obtained via manual analysis using a hex editor

## **GPT Artifacts**



#### **Deleted and Overwritten GUID Partitions**

#### Case 1:

- If the MBR disk is repartitioned or converted to GPT, then the sector zero will be generally overwritten with a protective MBR
- To recover data from the previous MBR partitioned volumes, investigators can use standard forensic methods used to perform an extensive search for file systems

#### Case 2:

- If the GPT disk is repartitioned or converted to MBR, then the GPT header and tables may remain intact based on the tool used
- Implementation of general partition deletion tools on a GPT disk might only delete the protective MBR, which can be recreated by simply reconstructing the disk

As per UEFI specification, if all the fields in a partition entry are zeroed, it implies that the entry is not in use. In this case, data recovery about deleted GUID partition entries is not possible

#### **GUID Identifiers**

- The GPT scheme provides GUIDs which are of investigative value as they are unique and hold potential information within them
- GUIDs possess unique identifying information for both disks and individual partitions
- Investigators can use tools such as uuid to decode various versions of GUID/UUID

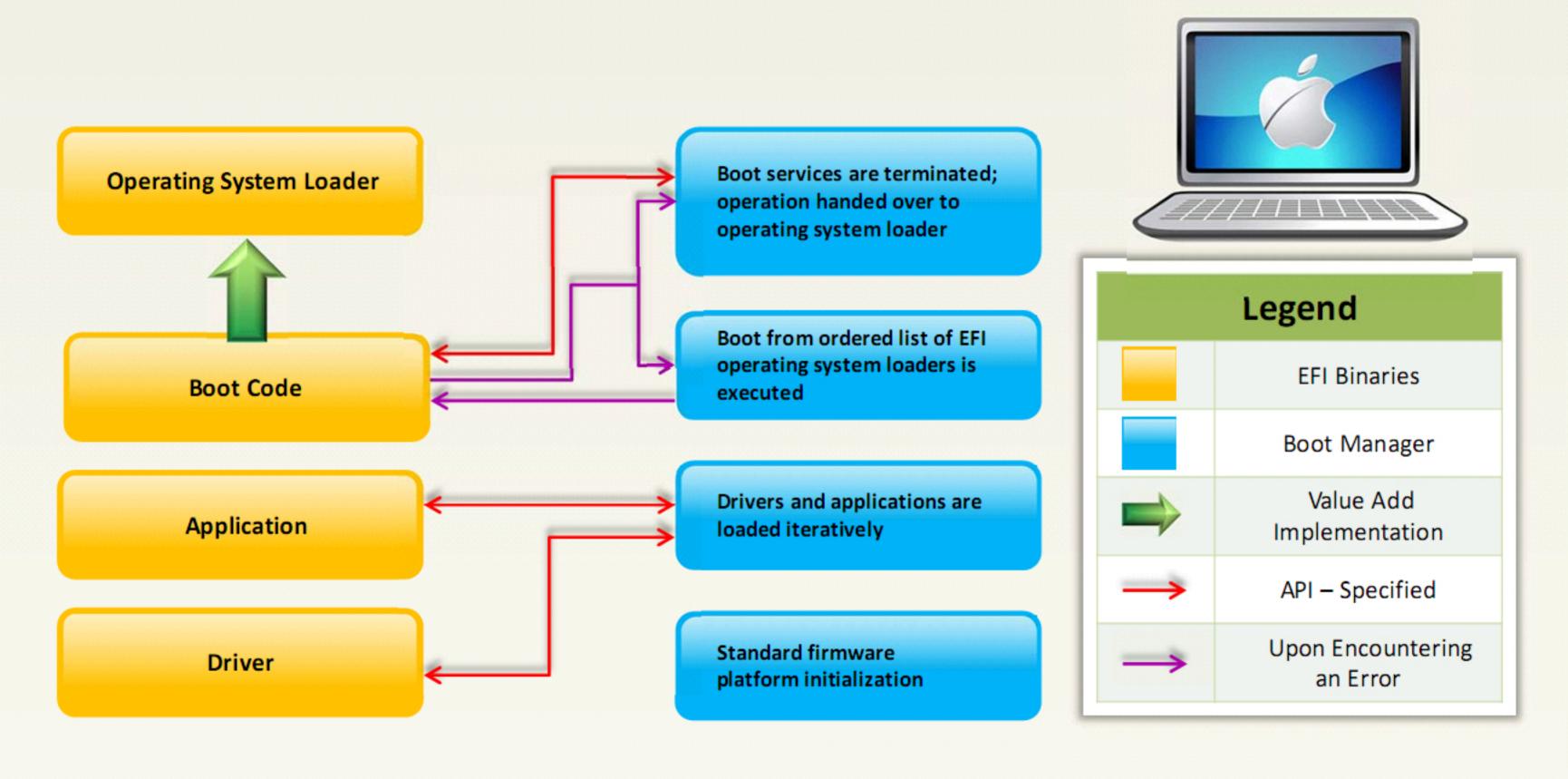
#### **Hidden Information on GPT Disks**

- Intruders may hide data on GPT disks as they do it on traditional MBR disks
- Data hiding places on GPT disks may be inter-partition gaps, unpartitioned space towards the end of the disk, GPT header, and reserved areas

Current forensic methods and tools to perform GPT analysis are not satisfactory

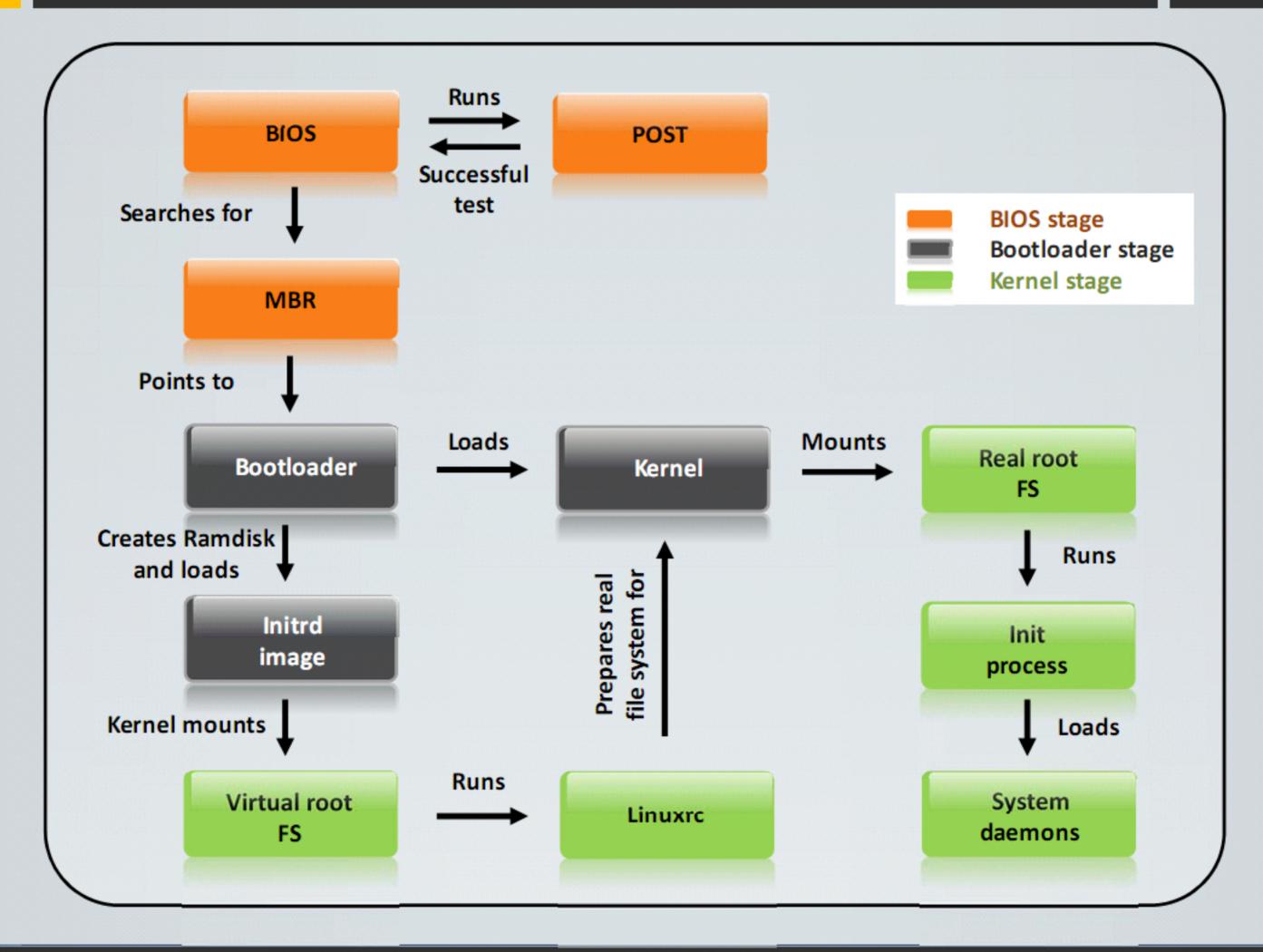
## **Macintosh Boot Process**





## Linux Boot Process





## Understanding File Systems





The file system is a set of data types, which is employed for storage, hierarchical categorization, management, navigation, access, and recovering the data



It provides a mechanism for users to store data logically in a hierarchy of files and directories



It also includes a **format** for specifying the path to a file through the structure of directories



They are organized in the form of tree-structured directories, and directories require access authorization



Major file systems include FAT, NTFS, HFS, HFS+, Ext2, Ext3, Ext4, etc.

## Types of File Systems



#### **Shared Disk File Systems**

In this file system, a number of systems (servers) can access same external disk subsystem

#### **Special Purpose File Systems**

In this file system, files are arranged dynamically by software, intended for such purposes as communication between computer processes or temporary file space

#### **Tape File Systems**

This file system is designed for storing and recovering the file on the tape in a self-describing form

#### **Disk File Systems**

This file system is designed for storing and recovering the file on the storage devices, usually a hard disk

#### Network File Systems

This file system is created to access the files on other computers that are connected by a network

#### Database File Systems

File management, where, instead of or in addition to hierarchically structured management, files are identified by their characteristics, such as the type of file, topic, author, or similar metadata

#### Flash File Systems

This file system is designed for storing and recovering the file on the **flash memory devices** 



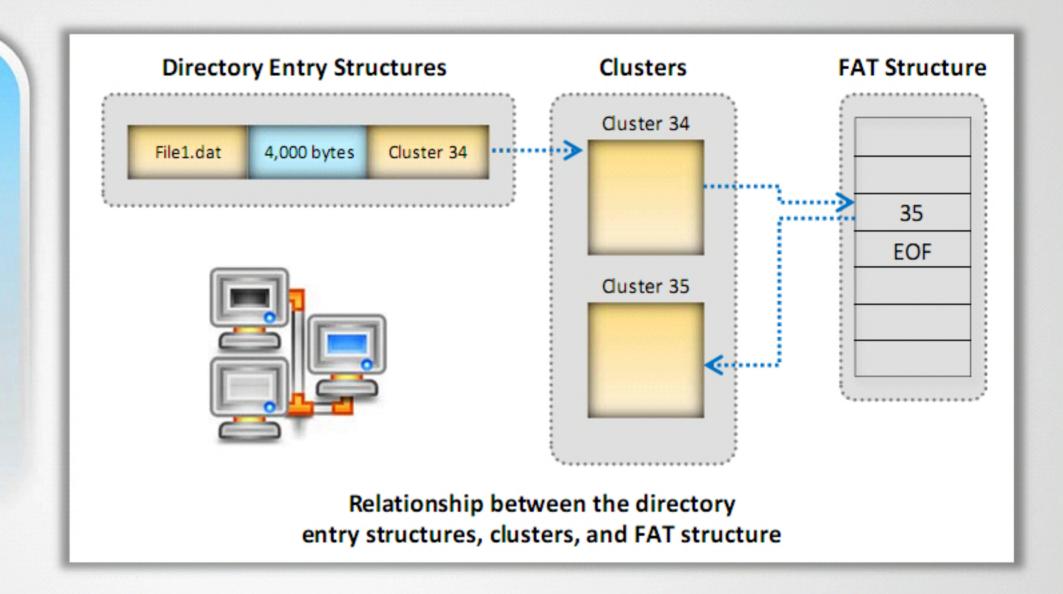


## Windows File Systems

## File Allocation Table (FAT)



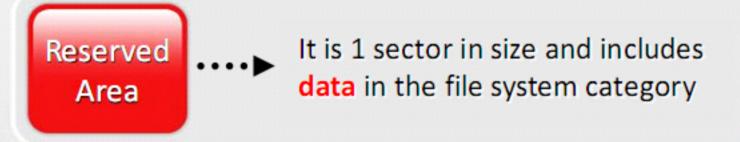
- The FAT file system is used with DOS, and it was the first file system used with the Windows OS
- It is named for its method of organization, the file allocation table, which resides at the beginning of the volume
- FAT contains three different versions (FAT12, FAT16, and FAT32) and differs due to the size of the entries in the FAT structure



System	Bytes Per Cluster within File Allocation Table	Cluster Limit
FAT12	1.5	Fewer than 4087 clusters
FAT16	2	Between 4,087 and 65,526 clusters, inclusive
FAT32	4	Between 65,526 and 268,435,456 clusters, inclusive

## FAT File System Layout





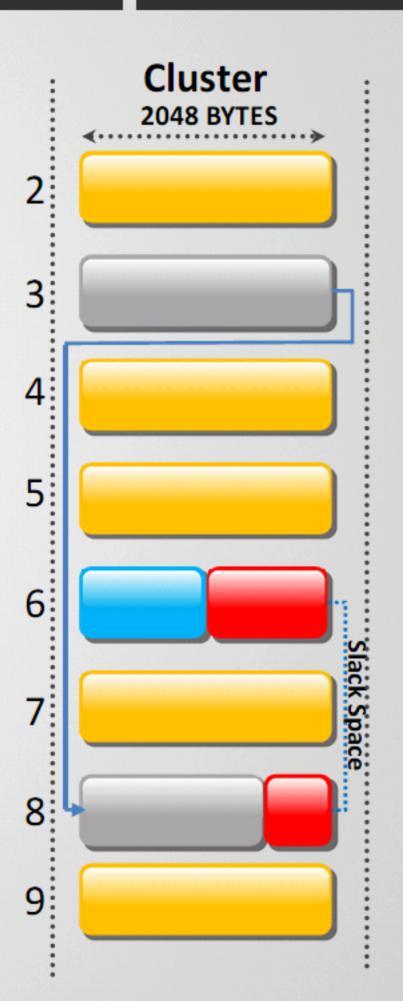
It contains the FAT structures and its size is calculated based on their number and size

Data
Area

It contains the clusters allocated
to store the content of the file and
directory



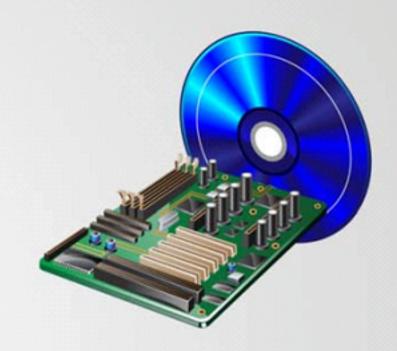
Root Directory					
File Name	Size	Cluster			
hello.jpg	3973	3			
gary.txt	1034	6			
File Allo	cation	Table			
Cluster	ı	Next			
2	0:	k0000			
3		8			
4		XFFF7 Cluster)			
5	0:	k0000			
6	0:	xFFFF			
7	0:	k0000			
8	0:	xFFFF			
9	0:	k0000			
Cluster = 2	048 = 4	Sectors			



### FAT Partition Boot Sector



- Boot Sector is the first sector (512 bytes) of a FAT file system
- FAT partition boot sector holds data that the file system uses to access the partition or volume
- MBR of x86-based computer systems uses this boot sector on the system partition to load the system kernel files



Byte Offset (in Hex)	Field Length	Sample Value	Meaning
00	3 bytes	EB 3C 90	Jump instruction
03	8 bytes	MSDOS5.0	OEM name in text
ОВ	25 bytes		BIOS Parameter Block (BPB)
24	26 bytes		Extended BIOS parameter block
3E	448 bytes		Bootstrap code
1FE	2 bytes	0x55AA	End of the sector marker

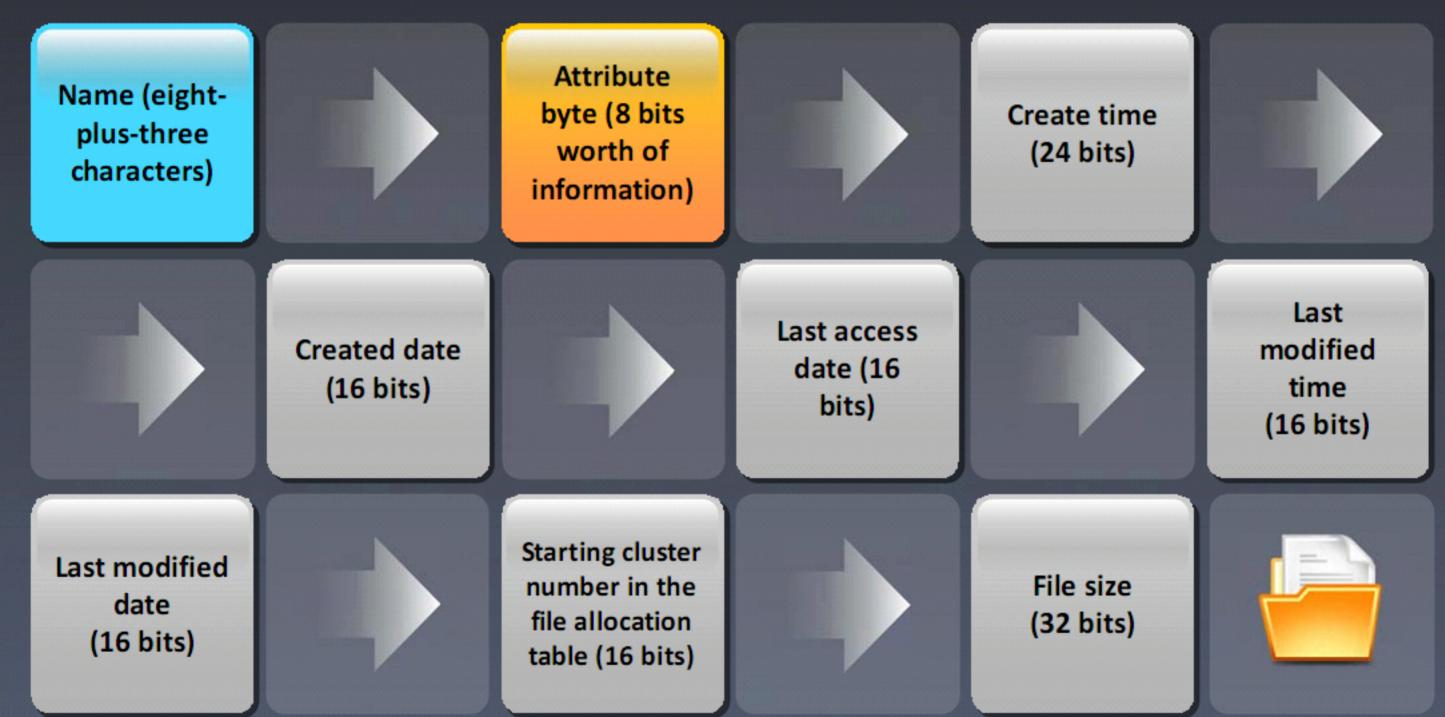
			-	0.000	1000										00000	00000		
Physical :	Sect	tor	C	v1 (	0, 5	Side	1	. Se	ct	or I	1							
000000000:	EB	30	90	4D	53	44	4F	53	-	35	2E	30	00	02	40	01	00	.<.MSD0S5.0@
00000010:	02	00	02	00	00	F8	FC	00	-	3F	00	40	00	3F	00	00	00	
00000020:	01	FO	3E	00	80	00	29	A8	-	8B	36	52	4E	4F	20	4E	41	>)6RNO NA
00000030:	4D	45	20	20	20	20	46	41	-	54	31	36	20	20	20	33	CO	ME FAT16 3.
00000040:	8E	DO	BC	00	70	68	CO	07	-	1F	AO	10	00	F7	26	16	00	h
00000050:	03	06	0E	00	50	91	B8	20	-	00	F7	26	11	00	88	1E	OB	P 6
00000060:	00	03	C3	48	F7	F3	03	C8	-	89	0E	08	02	68	00	10	07	Hh
00000070:	33	DB	8F	06	13	02	89	1E	-	15	02	0E	EΒ	90	00	72	57	3rW
00000080:	33	DB	88	OΕ	11	00	8B	FB	-	51	В9	OB	00	BE	DC	01	F3	3Q
00000090:	A6	59	74	05	83	C3	20	E2	-	ED	E3	37	26	8B	57	1A	52	.Yt 74.W.R
000000A0:	B8	01	00	68	00	20	07	33	-	DB	0E	E8	48	00	72	28	5B	h3H.r([
000000B0:	8D	36	0B	00	8D	3E	OB	02	-	1E	8F	45	02	C7	05	F5	00	.6>E
000000000:	1E	8F	45	06	C7	45	04	0E	-	01	8A	16	24	00	EA	03	00	EE
000000D0:	00	20	BE	86	01	EB	03	BE	-	AZ	01	E8	09	00	BE	Cl	01	
000000E0:	E8	03	00	FB	EB	FE	AC	OA.	-	CO	74	09	B4	0E	BB	07	00	t
000000F0:	CD	10	EB	F2	C3	50	4A	4A	-	AO	OD	00	32	E4	F7	EZ	03	PJJ2
00000100:	06	08	02	83	DZ	00	A3	13	-	02	89	16	15	02	58	AZ	07	X
00000110:	02	Al	13	02	8B	16	15	02	-	03	06	10	00	13	16	1E	00	
00000120:	F7	36	18	00	FE	CZ	88	16	-	06	02	33	DZ	F7	36	1A	00	.636
00000130:	88	16	25	00	A3	04	02	Al	-	18	00	ZA	06	06	02	40	AE	
00000140:	06	07	02	76	05	0.4	07	02	-	32	E4	50	B4	02	88	0E	04	v2.P
00000150:								02	-	86	E9	88	16	24	00	CD	13	
00000160:								F9	-	CB	58	28	06	07	02	76	11	X(v.
00000170:	01	06	13	02	83	16	15	02	-	00	F7	26	OB	00	03	DB	EB	
00000180:	90	A2	07	02	F8	CB	42	4F	-	4F	54	ЗA	20	43	6F	75	6C	BOOT: Coul
00000190:	64	6E	27	74	20	66	69	6E	-	64	20	4E	54	4C	44	52	OD	dn't find MTLDR.
000001A0:	OA	00	42	4F	4F	54	3A	20	-	49	2F	4F	20	65	72	72	6F	BOOT: I/O erro
000001B0:	72	20	72	65	61	64	69	6E	-	67	20	64	69	73	6B	OD	AO	r reading disk
000001CO:		50							-	69	6E	73	65	72	74	20	61	.Please insert a
000001D0:	6E	6F	74	68	65	72	20	64	-	69	73	6B	00	4E	54	40	44	nother disk.NTLD
000001E0:	52	20	20	20	20	20	20	00	-	00	00	00	00	00	00	00	00	R
000001F0:	00	00	00	00	00	00	00	00	-	00	00	00	00	00	00	55	AA	U.

### FAT Folder Structure



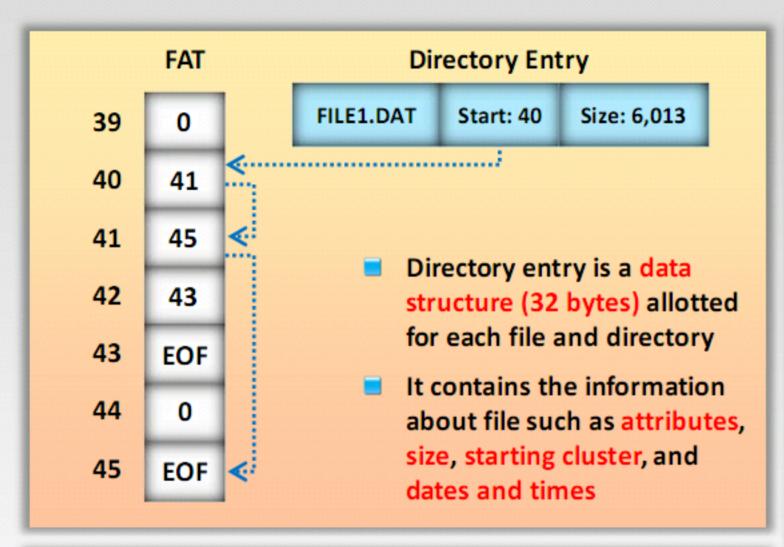
- FAT file systems have a set of 32-byte folder entries for every file
- Folder entries in FAT system:

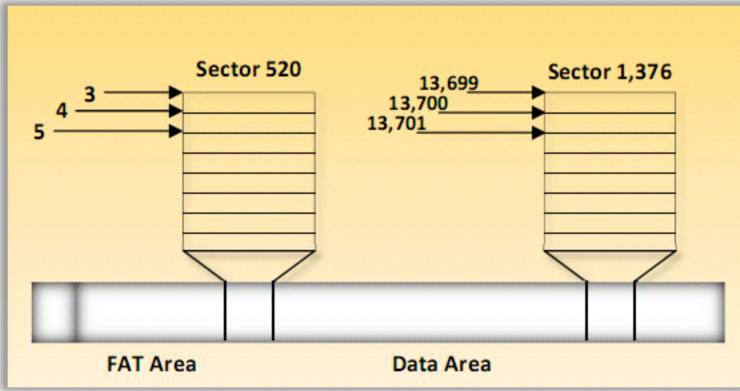




# **Directory Entries and Cluster Chains**







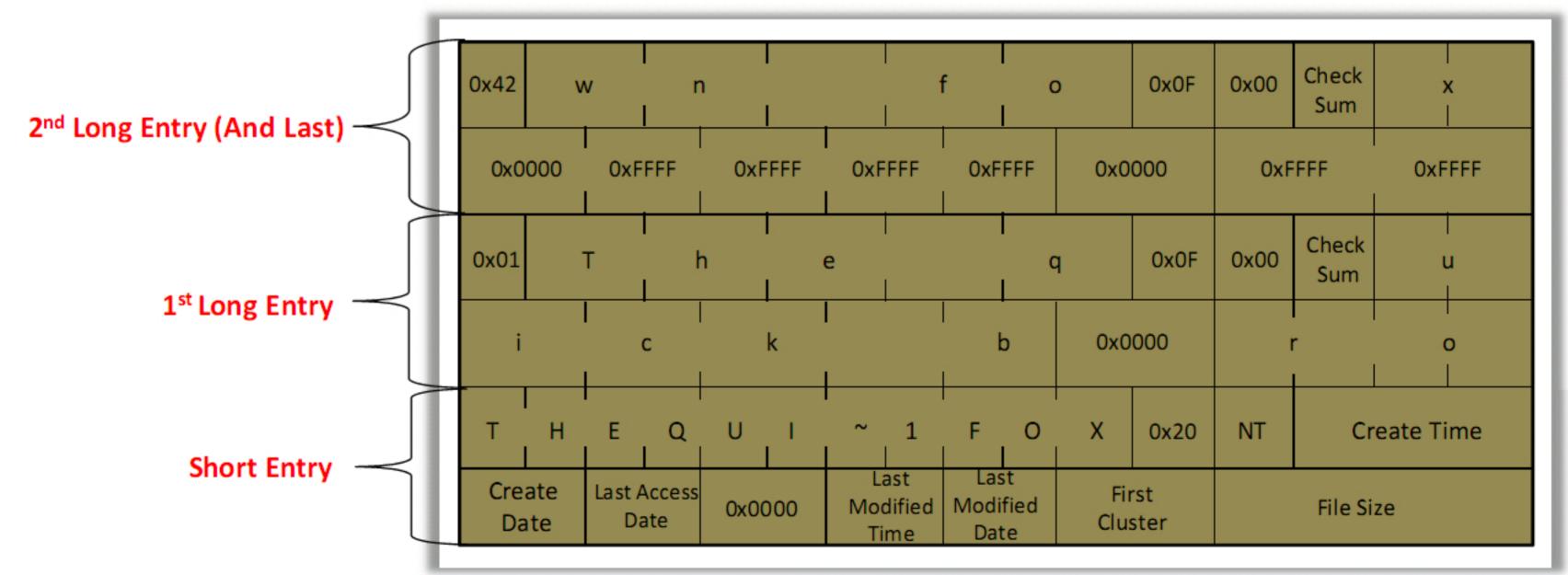
Byte Range	Description
0-0	First character of file name in ASCII and allocation status (0xe5 or 0x00 if unallocated)
1 – 10	Characters 2 to 11 of file name in ASCII
11 – 11	File Attributes
12 – 12	Reserved
13 – 13	Created time (tenths of second)
14 – 15	Created time (hours, minutes, seconds)
16 – 17	Created day
18 – 19	Accessed day
20 – 21	High 2 bytes of first cluster address (0 for FAT12 and FAT16)
22 – 23	Written time (hours, minutes, seconds)
24 – 25	Written day
26 – 27	Low 2 bytes of first cluster address
28 – 31	Size of file (0 for directories)

### Filenames on FAT Volumes



- Whenever users create or rename a file on FAT volume, Windows uses attribute bits to support long file names and creates an eight-plus-three name for the file
- Windows also creates many secondary folder entries for the file
- Below diagram shows all of the folder entries for the file Thequi~1.fox, which has a long name of The quick brown.fox





http://technet.microsoft.com

## FAT32

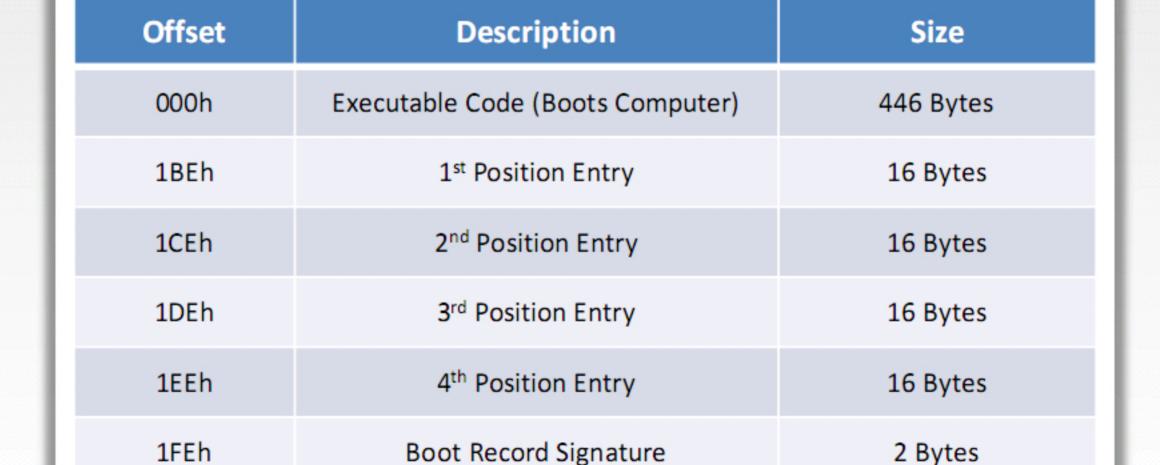




- FAT32 file system is derived from a FAT file system and supports drives up to 2 terabytes in size
- It uses drive space efficiently and uses small clusters
- It takes backup of the file allocation table instead of the default copy











MBR table of FAT32

# New Technology File System (NTFS)





NTFS is the **standard file system** of Windows NT and its descendants Windows XP, Vista, 7, 8.1,10, server 2003, server 2008, and server 2012



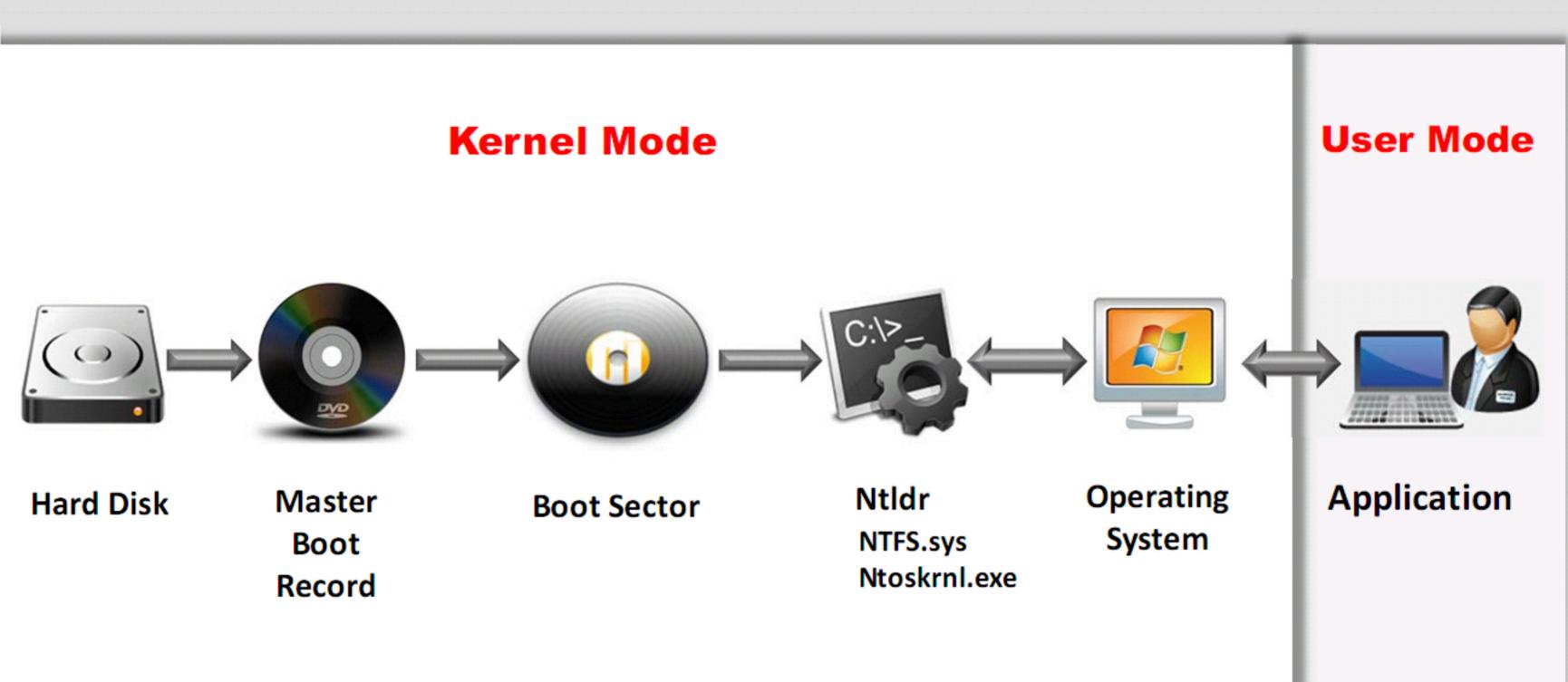
From Windows NT 3.1, it is the default file system of Windows NT family



It has several improvements over FAT such as improved support for metadata and the use of advanced data structures to improve performance, reliability, and disk space utilization plus additional extensions such as security access control lists and file system journaling

## NTFS Architecture





## NTFS System Files



File Name	Description
\$attrdef	Contains definitions of all system-and user-defined attributes of the volume
\$badclus	Contains all the bad clusters
\$bitmap	Contains bitmap for the entire volume
\$boot	Contains the volume's bootstrap
\$logfile	Used for recovery purposes
\$mft	Contains a record for every file
\$mftmirr	Mirror of the MFT used for recovering files
\$quota	Indicates disk quota for each user
\$upcase	Converts characters into uppercase Unicode
\$volume	Contains volume name and version number

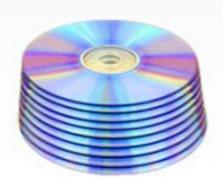
### NTFS Partition Boot Sector



- When a volume is formatted as an NTFS volume, the format program allocates the first 16 sectors for the boot sector and the bootstrap code
- Partition identifier: 0x07 (MBR) EBD0A0A2-B9E5-4433-87C0-68B6B72699C7 (GPT)

#### **Boot Sector of an NTFS Volume:**

Byte Offset	Field Length	Field Name
0x00	3 bytes	Jump Instruction
0x03	LONGLONG	OEM ID
0x0B	25 bytes	BIOS Parameter Block (BPB)
0x24	48 bytes	Extended BPB
0x54	426 bytes	Bootstrap Code
0x01FE	WORD	End of Sector Marker





```
Physical Sector: Cyl 0, Side 1, Sector 1
     00000000:EB 52 90 4E 54 46 53 20 -20 20 20 00 02 08 00 00 .R.NTFS ......
     00000020:00 00 00 00 80 00 80 00 -4A F5 7F 00 00 00 00 ......J.....
     00000030:04 00 00 00 00 00 00 -54 FF 07 00 00 00 00 ......T.....
     00000090: OF B6 D1 80 E2 3F F7 E2 -86 CD C0 ED 06 41 66 OF .....?......Af.
     000000B0:16 24 00 CD 13 72 OF 81 -FB 55 AA 75 09 F6 C1 01 .$...r...U.u....
     000000C0:74 04 FE 06 14 00 C3 66 -60 1E 06 66 R1 10 00 66 t.....f`..f...f
            :B4 42 8A 16 24 00 16 1F -8B F4 CD 13 66 58 5B 07 .B..$.....fX [..
     00000110:66 58 66 58 1F EB 2D 66 -33 D2 66 0F B7 0E 18 00 fXfX.-f3.f.....
     00000120:66 F7 F1 FE C2 8A CA 66 -8B D0 66 C1 EA 10 F7 36 f.....f...f....6
     000001A0:0D 0A 4E 54 4C 44 52 20 -69 73 20 6D 69 73 73 69 ..NTLDR is missi
     000001B0:6E 67 00 0D 0A 4E 54 4C -44 52 20 69 73 20 63 6F ng...NTLDR is co
     000001C0:6D 70 72 65 73 73 65 64 -00 0D 0A 50 72 65 73 73 mpressed...Press
     000001D0:20 43 74 72 6C 2B 41 6C -74 2B 44 65 6C 20 74 6F Ctrl+Alt+Del to
     000001E0:20 72 65 73 74 61 72 74 -OD OA 00 00 00 00 00 00 restart......
```

## Cluster Sizes of NTFS Volume



- A cluster is the smallest allocation unit on the hard disk that is used to hold a file
- NTFS uses clusters of different sizes to hold the files, depending on the size of the NTFS volume
- List of the default cluster sizes for NTFS volume:

Sectors Per Cluster	Default Cluster Size
1	512 bytes
2	1024 bytes (1 GB)
4	2048 bytes (2 GB)
8	4 KB
	Per Cluster 1



## NTFS Master File Table (MFT)



Each file on an NTFS volume is represented by a record in a special file called the master file table (MFT)

[ash] Loggohost:~root\$\_

It reserves the first 16 records of the table for special information

2 5

If the first MFT record is corrupted,
NTFS reads the second record to find
the MFT mirror file, whose first record
is identical to the first record of the
MFT

The locations of the data segments for both the MFT and MFT mirror file are recorded in the boot sector, a duplicate of the boot sector is located at the logical center of the disk

The first record of this table describes the master file table itself, followed by an MFT mirror record



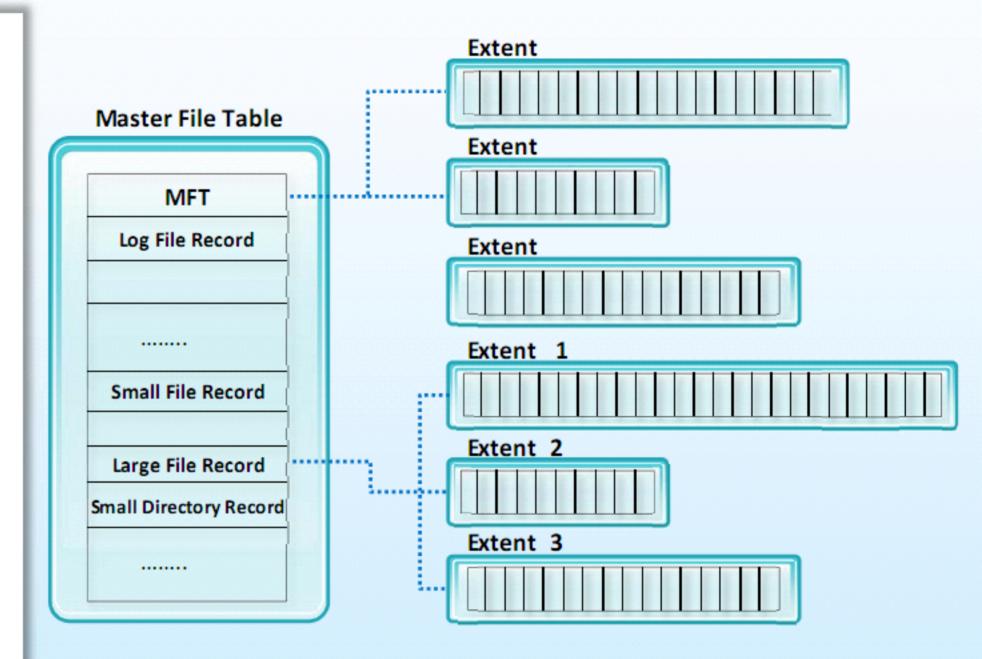


The third record of the MFT is the log file, used for file recovery. The seventeenth and following records of the master file table are for each file and directory (also viewed as a file by NTFS) on the volume

# NTFS Master File Table (MFT) (Cont'd)



- MFT is a relational database, which consists of information related to the files and the file attributes
- The rows consist of file records and the columns consist of file attributes
- It has information of every file on the NTFS volume including its own information
- It has 16 records reserved for system files
- For small folder, MFT is represented as follows:



Structure of a Master File Table on an NTFS volume

Standard Directory Information Name Data or Unused Space

# Metadata Files Stored in the MFT



System File	File Name	MFT Record	Purpose of the File
Master file table	\$Mft	0	It contains one base file record for each file and folder on an NTFS volume
Master file table mirror	\$MftMirr	1	It guarantees access to the MFT in case of a single-sector failure
Log file	\$LogFile	2	It contains information used by NTFS for faster recoverability
Volume	\$Volume	3	It contains information about the volume
Attribute definitions	\$AttrDef	4	It lists attribute names, numbers, and descriptions
Root file name index		5	The root folder
Cluster bitmap	\$Bitmap	6	It represents the volume by showing free and unused clusters
Boot sector	\$Boot	7	It includes the BPB used to mount the volume
Bad cluster file	\$BadClus	8	It contains bad clusters for a volume
Security file	\$Secure	9	It contains unique security descriptors for all files within a volume
Upcase table	\$Upcase	10	It converts lowercase characters to matching Unicode uppercase characters
NTFS extension file	\$Extend	11	It is used for various optional extensions

## NTFS Attributes





Every file has unique identities such as name, security information, and metadata of file system in the file



Every attribute is identified by an attribute type code and attribute name



There are two categories of attributes:

- Resident attributes: These are the attributes that are contained in the MFT
- Non-resident attributes: These are the attributes that are allocated with one or more clusters of disk space

Attribute Type	Purpose of the Attribute
Standard information	Lists the information regarding the time stamp data and link count information
Attribute list	List of attributes that are not in the MFT
File name	The file name is stored here and can be a long or short name
Security descriptor	Ownership and access rights to the file are listed here
Data	Stores file data
Object id	File identifier volume –unique identifier
Logged tool stream	Used by the encrypted file system service
Reparse point	Volume mount point used for installable file system filter drivers
Index root	Employed for use of folders and files
Index allocation	Employed for use of folders and files
Bitmap	Employed for use of folders and files
Volume information	Version number of the volume is listed
Volume name	The volume label is listed here

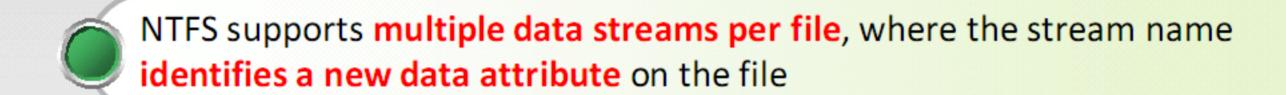
## NTFS Data Stream





NTFS data stream is a unique set of file attributes







C:\>ECHO text\_message > myfile.txt:stream1



Command to display the contents of the data stream: C:\>MORE < myfile.txt:stream1



A data stream does not appear when a file is opened in a text editor. The only way to see if a data stream is attached to a file is by examining the MFT entry for the file



When you copy an NTFS file to a FAT volume, such as a floppy disk, data streams and other attributes not supported by FAT are lost

## NTFS Data Stream (Cont'd)



```
×
 Administrator: Command Prompt
C:\promgry>dir
 Volume in drive C has no label.
 Volume Serial Number is 746E-941F
Directory of C:\promqry
04/04/2016 02:07 PM
04/04/2016 02:07 PM
                       <DIR>
                                7,487 EULA.txt
12/07/2004 02:25 PM
                              32,768 promgry.exe
01/31/2005 01:07 PM
                               40,255 bytes
              2 File(s)
              2 Dir(s) 65,523,535,872 bytes free
C:\promqry>_
```

```
Administrator: Command Prompt
                                                                  ×
C:\promary>dir
Volume in drive C has no label.
Volume Serial Number is 746E-941F
Directory of C:\promgry
04/04/2016 02:07 PM
                       <DIR>
04/04/2016 02:07 PM
                       <DIR>
12/07/2004 02:25 PM
                               7,487 EULA.txt
01/31/2005 01:07 PM
                               32,768 promqry.exe
              2 File(s)
                               40,255 bytes
              2 Dir(s) 65,523,535,872 bytes free
C:\promqry>ECHO text message > myfile.txt:stream1
C:\promqry>
```

```
×
 Administrator: Command Prompt
Directory of C:\promqry
04/04/2016 02:07 PM
                       <DIR>
04/04/2016 02:07 PM
                       <DIR>
12/07/2004 02:25 PM
                                7,487 EULA.txt
01/31/2005 01:07 PM
                               32,768 promgry.exe
              2 File(s)
                                40,255 bytes
              2 Dir(s) 65,523,535,872 bytes free
C:\promqry>ECHO text_message > myfile.txt:stream1
C:\promqry>more < myfile.txt:stream1
text message
C:\promqry>_
```

```
Administrator: Command Prompt
C:\promqry>dir
Volume in drive C has no label.
Volume Serial Number is 746E-941F
Directory of C:\promqry
04/04/2016 02:09 PM
                       <DIR>
34/94/2016 02:09 PM
12/97/2004 02:25 PM
                                7,487 EULA.txt
04/04/2016 02:09 PM
                                    @ myfile.txt
01/31/2005 01:07 PM
                               32,768 promgry.exe
              3 File(s)
                               40,255 bytes
              2 Dir(s) 65,522,970,624 bytes free
C:\promqry>_
```

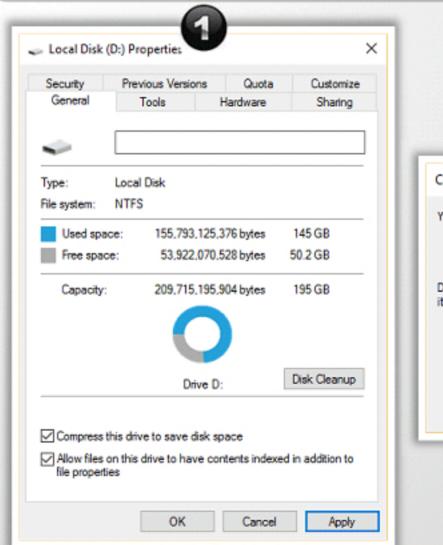
### NTFS Compressed Files



- The compressed files present on an NTFS volume can be read and written by any Windows-based application without first being decompressed by another program
- NTFS promotes compression of individual files, all the files within a folder, and all the files/folders within an NTFS volume
- The file is automatically decompressed by filter driver when Windows applications requests the access
- NTFS compression algorithms support cluster sizes of up to 4 KB

#### **Setting the Compression State of a Volume:**

- Right-click on the drive that is to be compressed and click Properties
- On the General tab, choose "Compress this drive to save disk space" check box and click Apply
- In the Confirm Attribute Changes dialog box, choose an option and click OK

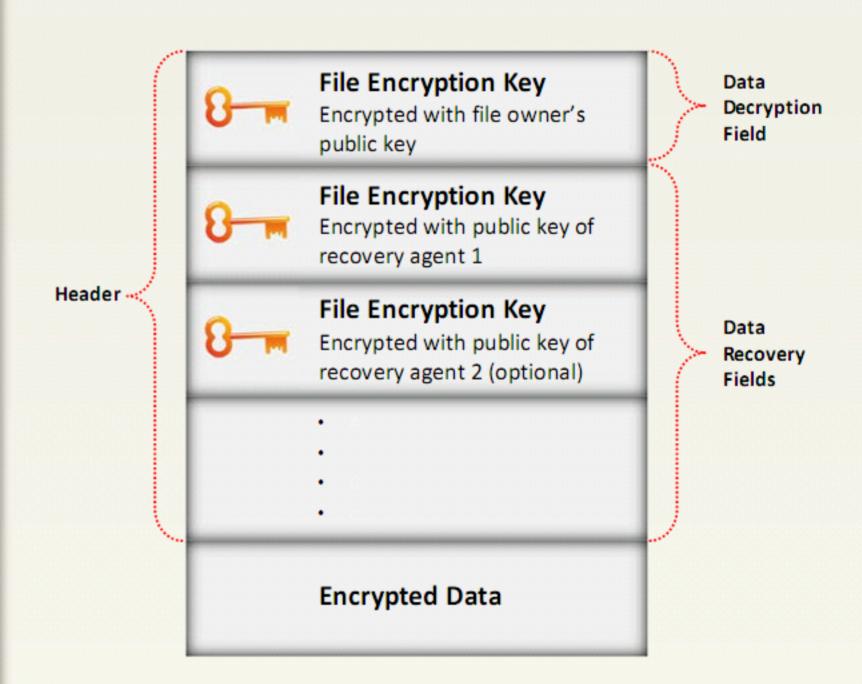




### **Encrypting File Systems (EFS)**

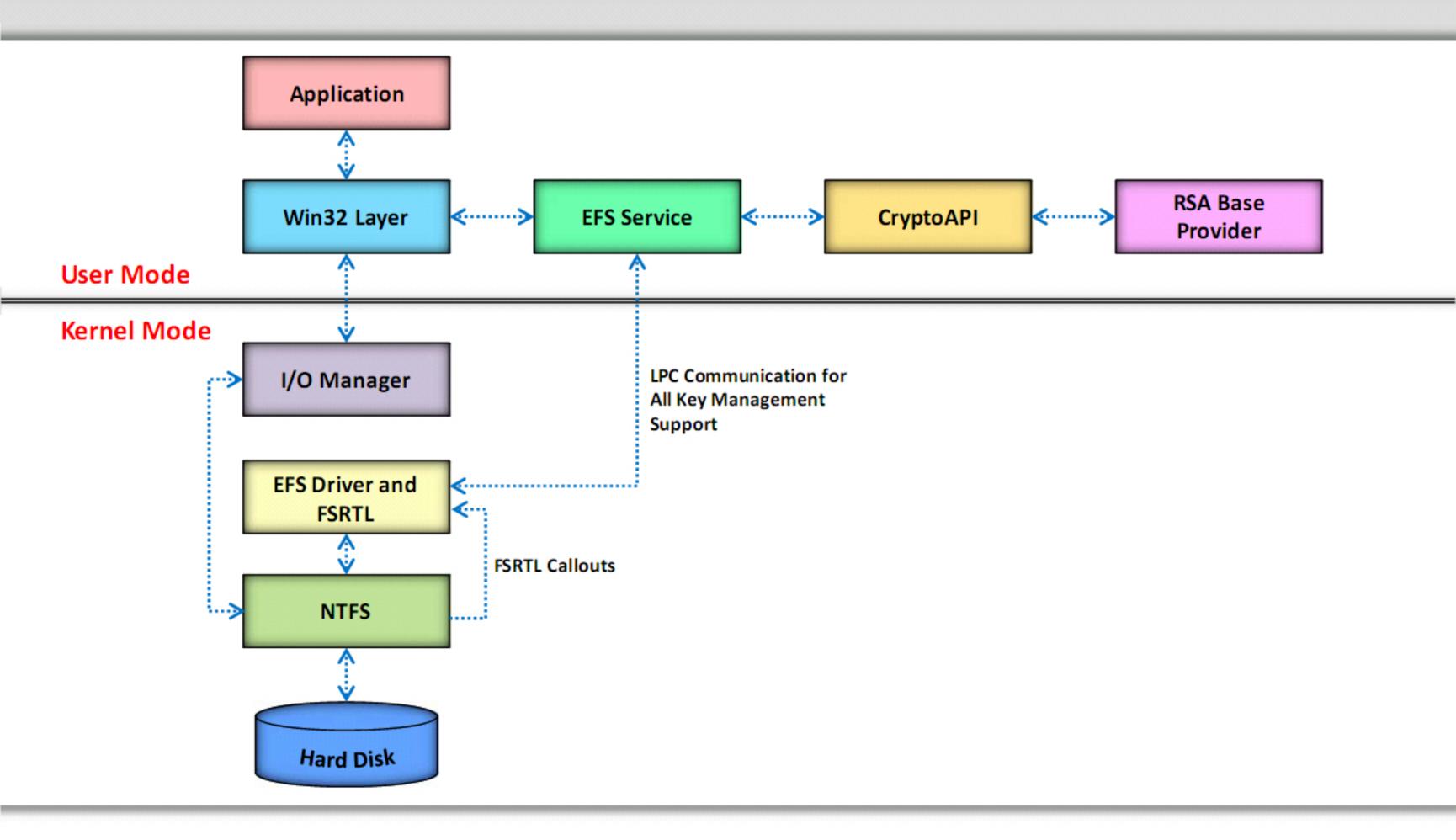


- Encrypting File System (EFS) was first introduced in version 3.0 of NTFS, that offers filesystem-level encryption
- This encryption technology maintains a level of transparency to the user who encrypted the file, which means there is no need for users to decrypt the file to access it to make changes
- After a user is done with the file, the encryption policy is automatically restored
- When any unauthorized user tries to access an encrypted file, he or she is denied access
- To enable the encryption and decryption facilities, a user has to set the encryption attributes of the files and folders that the user wants to encrypt or decrypt



## Components of **EFS**



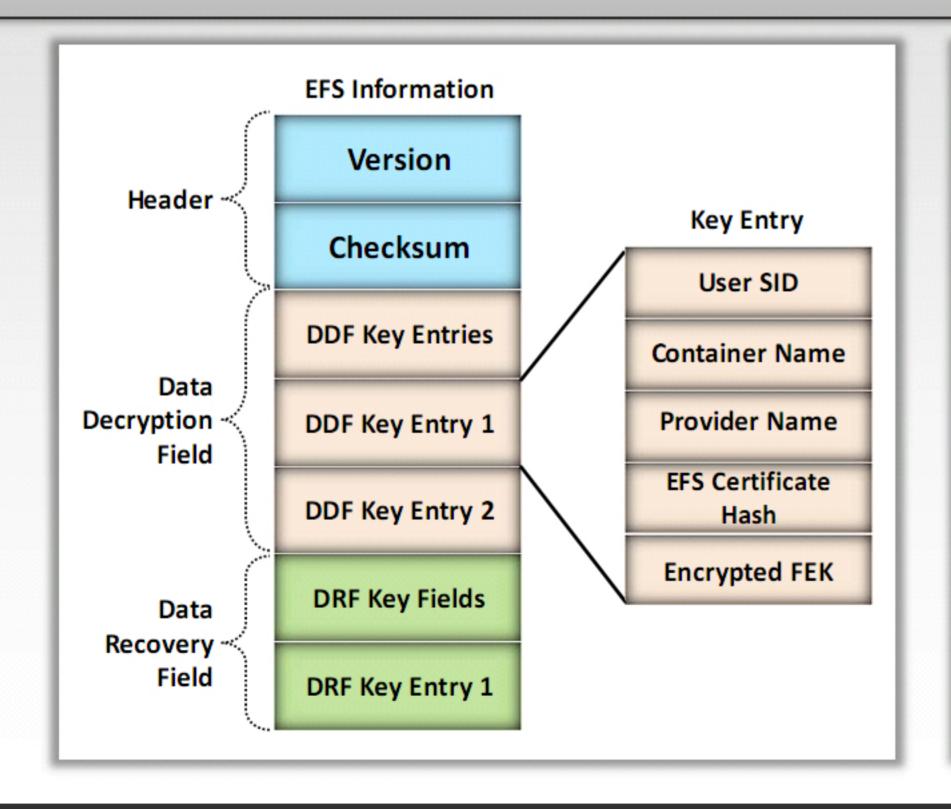


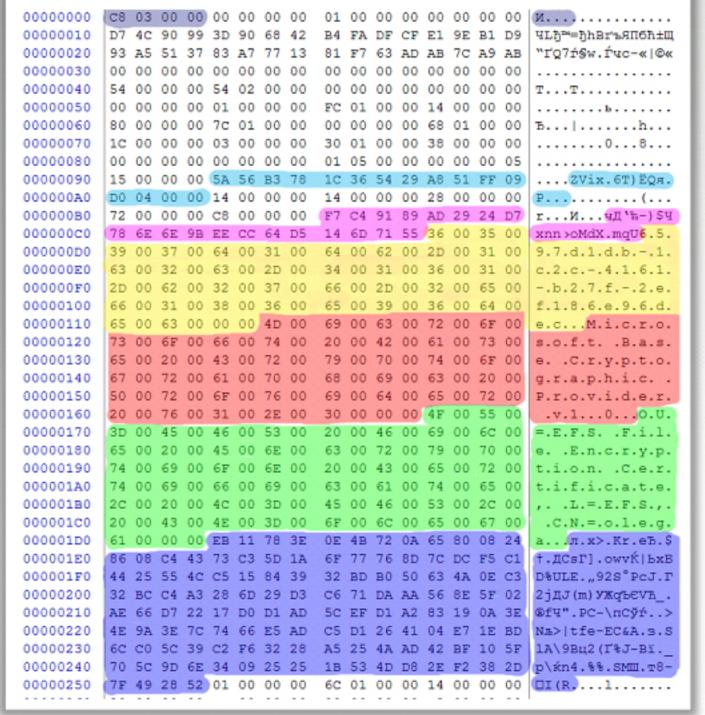
### EFS Attribute



- NTFS sets a flag for the file once you encrypt it and creates an EFS attribute where it stores Data Decryption Field (DDF) and Data Recovery Field (DDR)
- This attribute has Attribute ID = 0x100 in NTFS







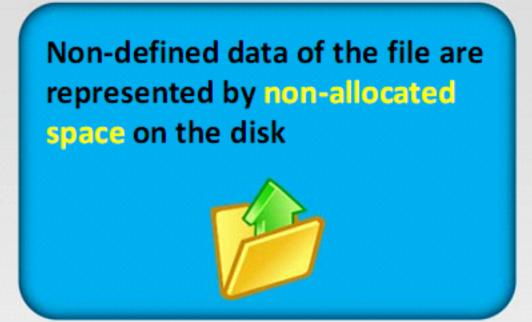
## Sparse Files

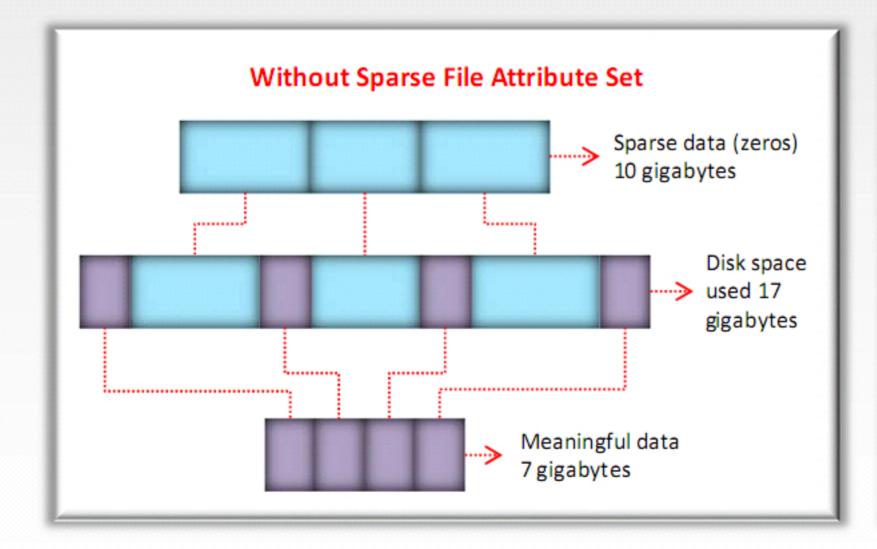


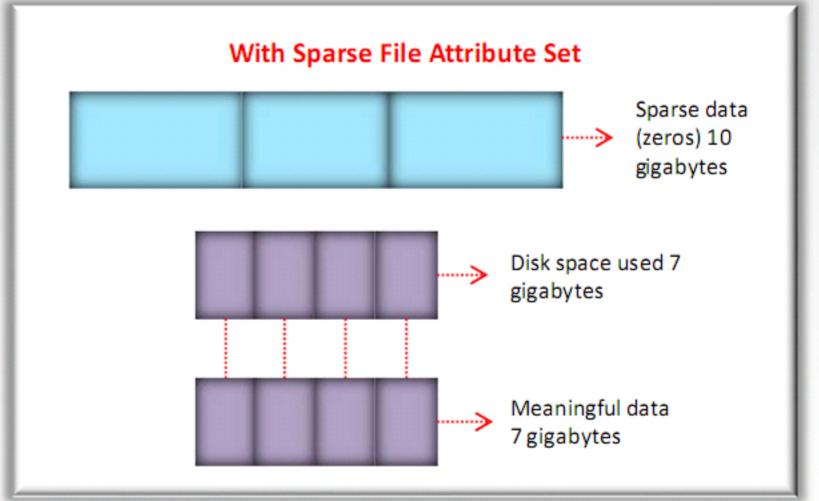
Sparse files provide a method of saving disk space for files by allowing I/O subsystem to allocate only meaningful (nonzero) data



If NTFS file is marked as sparse, it assigns hard disk cluster only for the data defined by the application







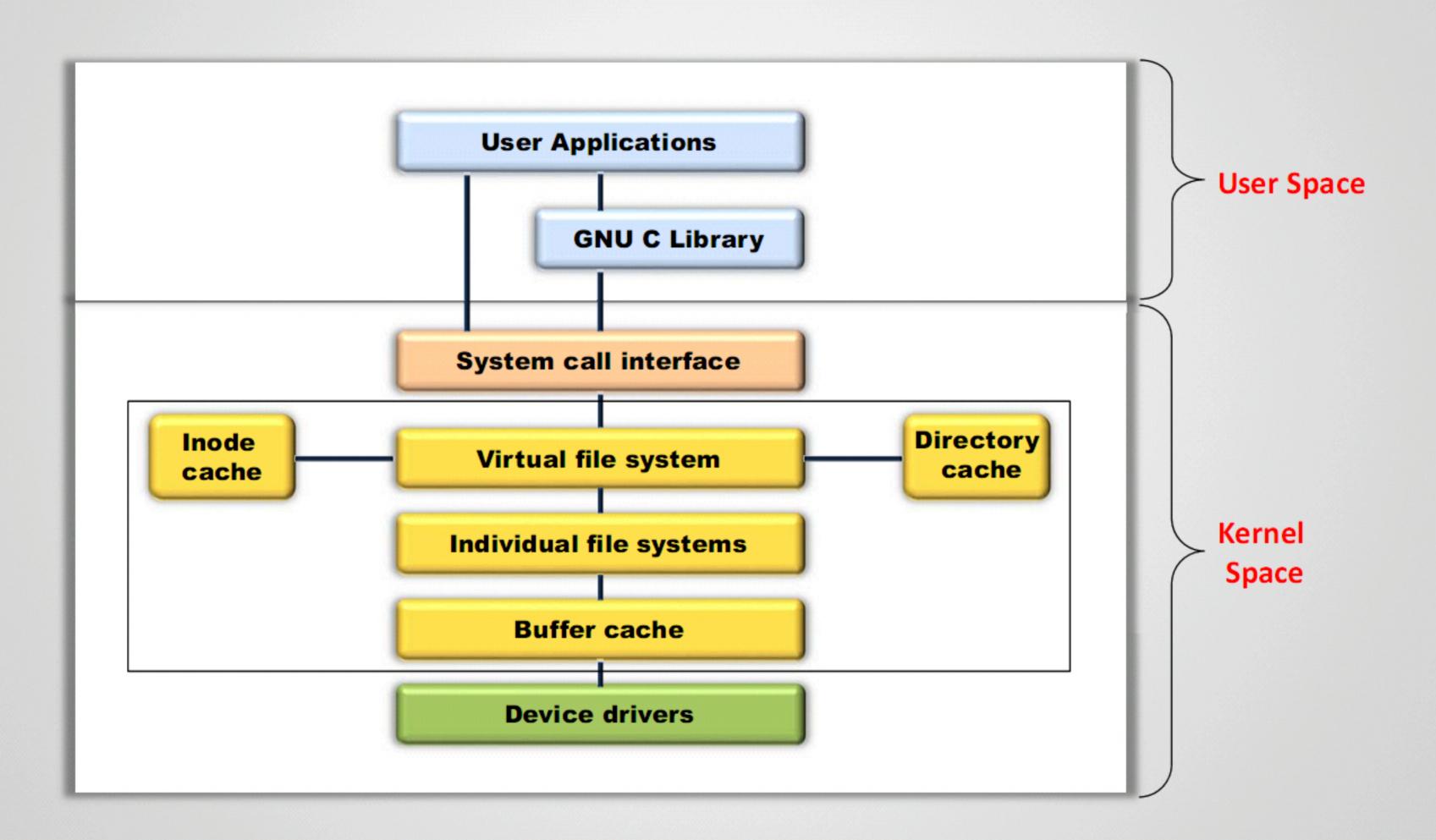




# Linux File Systems

### Linux File System Architecture





# Filesystem Hierarchy Standard (FHS)



#### Table displaying directories and their description specific to the FHS

- The Filesystem Hierarchy Standard (FHS) defines the directory structure and its contents in Linux and Unix-like operating systems
- In the FHS, all files and directories are present under the root directory (represented by /)

Directory	Description
/bin	Essential command binaries. Ex: cat, ls, cp.
/boot	Static files of the boot loader. Ex: Kernels, Initrd
/dev	Essential device files. Ex: /dev/null
/etc	Host-specific system configuration files
/home	Users' home directories, holding saved files, personal settings, etc.
/lib	Essential libraries for the binaries in /bin/ and /sbin/
/media	Mount points for removable media
/mnt	Temporarily mounted filesystems
/opt	Add-on application software packages
/root	Home directory for the root user
/proc	Virtual file system providing process and kernel information as files
/run	Information about running processes. Ex: running daemons, currently logged- In users
/sbin	Contains the binary files required for working
/srv	Site-specific data for services provided by the system
/tmp	Temporary files
/usr	Secondary hierarchy for read-only user data
/var	Variable data. Ex: logs, spool files, etc.

## Extended File System (EXT)



- First file system for the Linux operating system to overcome certain limitations of the Minix file system
- It has a maximum partition size of 2 GB and a maximum file name size of 255 characters
- It removes the two major Minix file system limitations of a 64 MB partition size and short file names
- The major limitation of this file system is that it doesn't support separate access, inode modification, and data modification time stamps
- It is replaced by the second extended file system

# Second Extended File System (EXT2)





EXT2 is a standard file system that uses improved algorithms, which greatly enhances its speed significantly, and it maintains additional time stamps



It maintains a special field in the superblock that keeps track of the file system status and identifies it as either clean or dirty

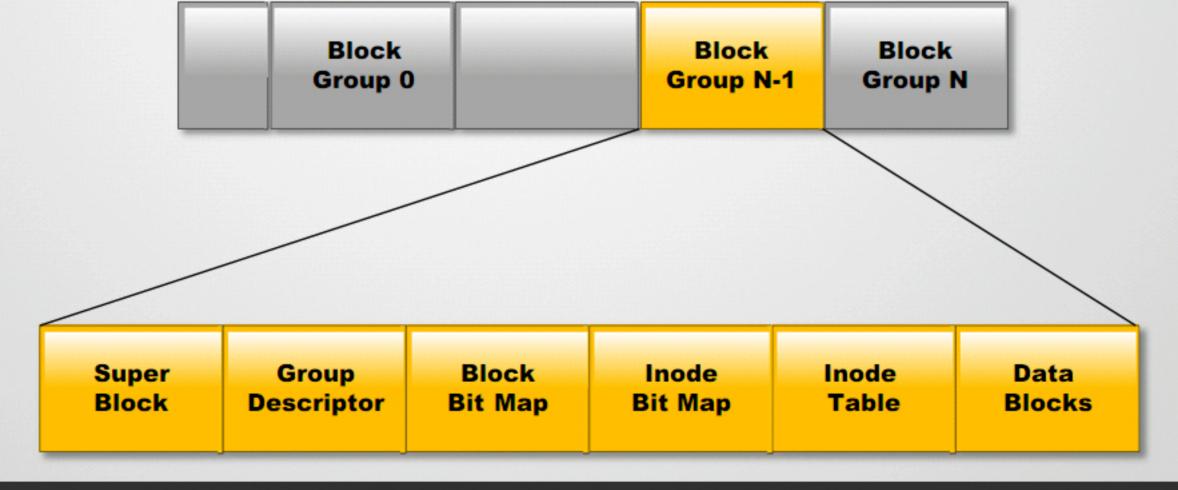


Its major shortcomings are the risk of file system corruption when writing to EXT2, and it is not a journaling file system



Physical layout of the EXT2 File system:





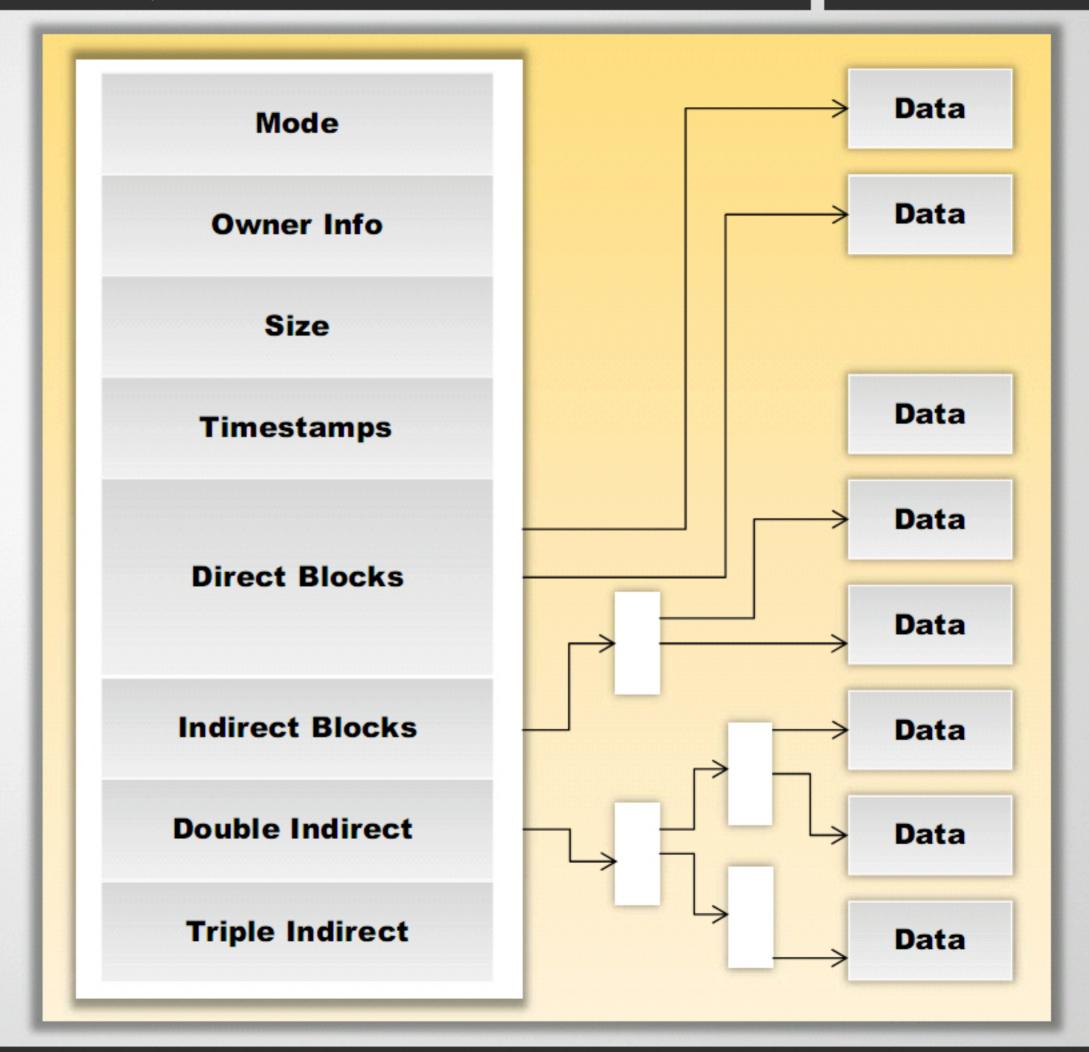
# Second Extended File System (EXT2) (Cont'd)



### **EXT2 Inode**

- The inode is a basic building block of the EXT2 file system
- Each file and directory are described by a single inode
- The inodes for each file system block are placed together in an inode table





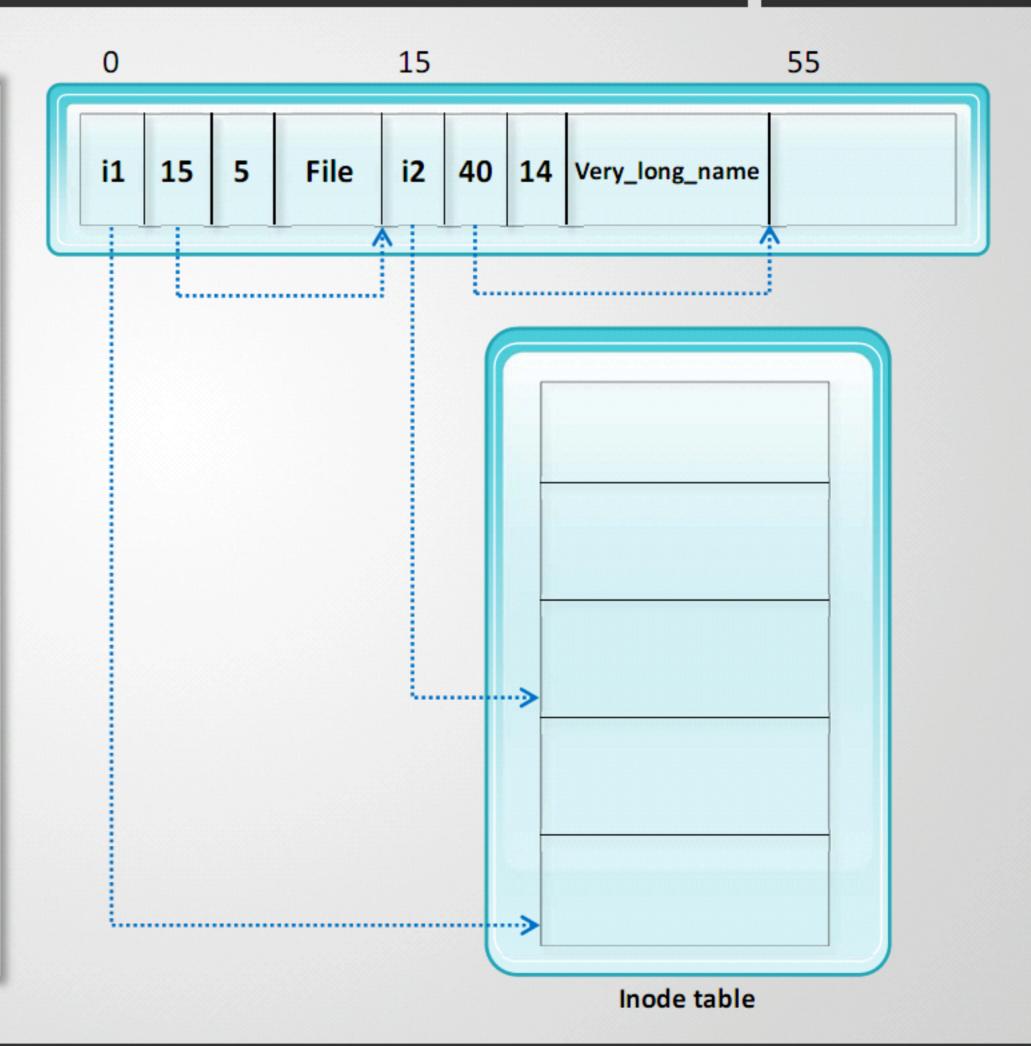
# Second Extended File System (EXT2) (Cont'd)



### **EXT2 Directories**

- EXT2 directories are particular files that create and hold the access path of the files in the file system
- These files contain the list of directory entries with the following information:
  - Directory inode
  - Length of the file name
  - Name of the directory





# Third Extended File System (Ext3)



- Ext3 is a journaling version of the EXT2 file system and is greatly used with the Linux operating system
- It is the enhanced version of the EXT2 file system
- It uses file system maintenance utilities (like fsck) for maintaining and repairing alike EXT2 file system
- Command to convert EXT2 to EXT3 file system:
  - # /sbin/tune2fs -j <partition-name>

### **Ext3 Features**

### **Data Integrity**

It provides stronger data integrity for events that occur due to computer system shutdowns



### Speed

As the EXT3 file system is journaling the file system, it has higher throughput in most cases than EXT2



### **Easy Transition**

The user can easily change the file system from EXT2 to EXT3 and increase the performance of the system



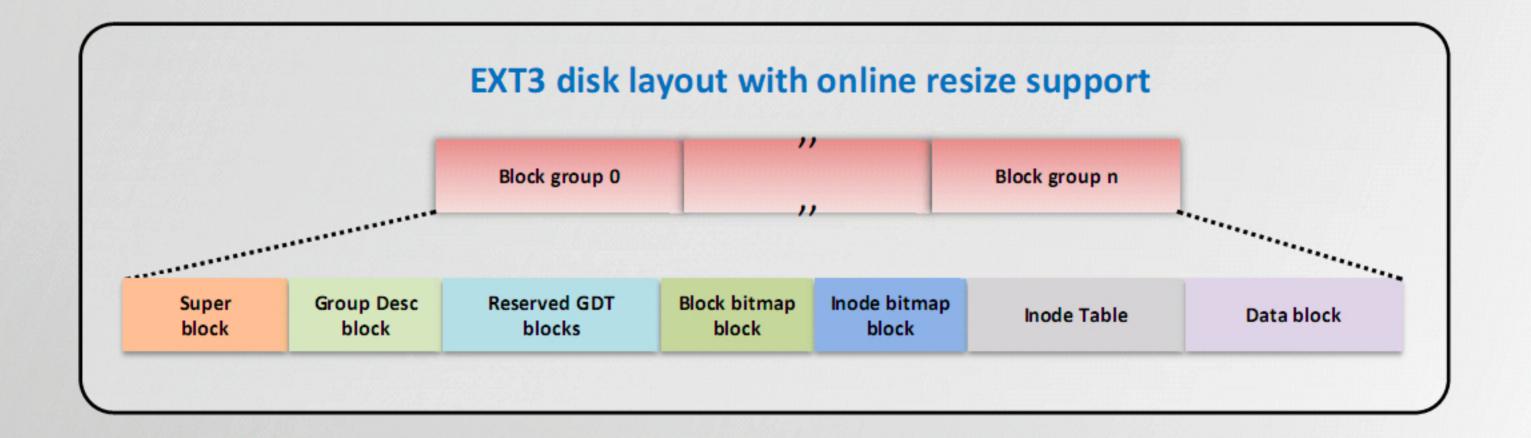
# Third Extended File System (EXT3) (Cont'd)



### File System Journaling:

- It records file system updates which help in quick file system recovering in case of a system crash
- The EXT3 journal uses inode 8 and its location is stated in the superblock
- The first block in the EXT3 journal is for superblock and contains general information





# Fourth Extended File System (EXT4)



- EXT4 is a journaling file system, developed as the replacement of commonly used EXT3 file system
- With incorporation of new features, EXT4 has significant advantages over EXT3 and EXT2 file systems particularly in terms of performance, scalability, and reliability
- Supports Linux Kernel v2.6.19 onwards

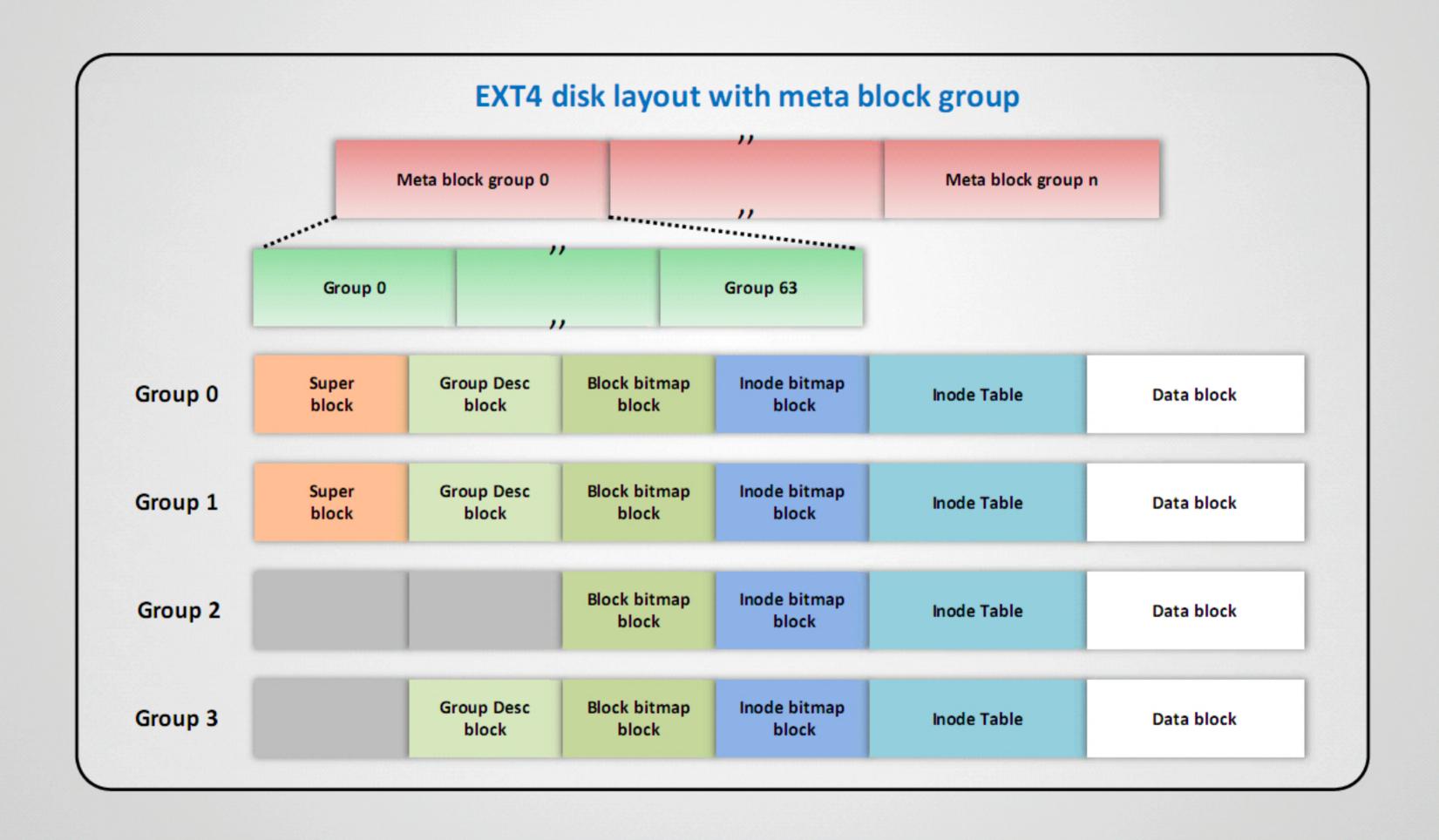
### Key Features



- File System Size supports maximum individual file size 16TB and overall maximum EXT4 file system size 1EB (exabyte)
- Extents replaces block mapping scheme used by EXT2 and EXT3, improving large file performance and reducing fragmentation
- Delayed allocation improves performance and reduces fragmentation by effectively allocating larger amounts of data at a time
- Multi-block allocation allocates files contiguously on disk
- fsck speed supports faster file system checking
- Journal checksumming uses checksums in the journal to improve reliability
- Persistent preallocation preallocates on-disk space for a file
- Improved Timestamps provides timestamps measured in nanoseconds
- Backward compatibility makes it possible to mount EXT3 and EXT2 as EXT4

# Fourth Extended File System (EXT4) (Cont'd)









## Mac OS X File Systems

## Mac OS X File Systems





### Hierarchical File System (HFS)

Developed by Apple Computer to support Mac operating system

### **HFS Plus**

HFS Plus (HFS+) is a successor of HFS and is used as a primary file system in Macintosh

### **UNIX File System (UFS)**

- Derived from the Berkeley Fast File System (FFS) that was originally developed at Bell Laboratories from the first version of UNIX FS
- All BSD UNIX derivatives including FreeBSD, NetBSD, OpenBSD, NeXTStep, and Solaris use a variant of UFS
- Acts as a substitute for HFS in Mac OS X

## HFS vs. HFS Plus



Feature	HFS	HFS Plus	Benefit/Comment
User visible name	Mac OS Standard		
Number of allocation blocks	16 bits worth	32 bits worth	Radical decrease in disk space used on large volumes, and a large number of files per volume
Long file names	31 characters	255 characters	Obvious user benefit; also improves cross-platform compatibility
File name encoding	MacRoman	Unicode	Allows for international-friendly file names, including mixed script names
File/folder attributes	Support for fixed size attributes (FileInfo and ExtendedFileInfo)	Allows for future metadata extensions	Future systems may use metadata for a richer Finder experience
OS startup support	System Folder ID	Also supports a dedicated startup file	May help non-Mac OS systems to boot from HFS Plus Volumes
Catalog node size	512 bytes	4 KB	Maintains efficiency in the face of the other changes. (This larger catalog node size is due to the much longer file names [512 bytes as opposed to 32 bytes], and larger catalog records (because of more/larger fields)).
Maximum file size	2 <sup>31</sup> bytes	2 <sup>63</sup> bytes	Obvious user benefits, especially for multimedia content creators.

## Hierarchical File System (HFS)

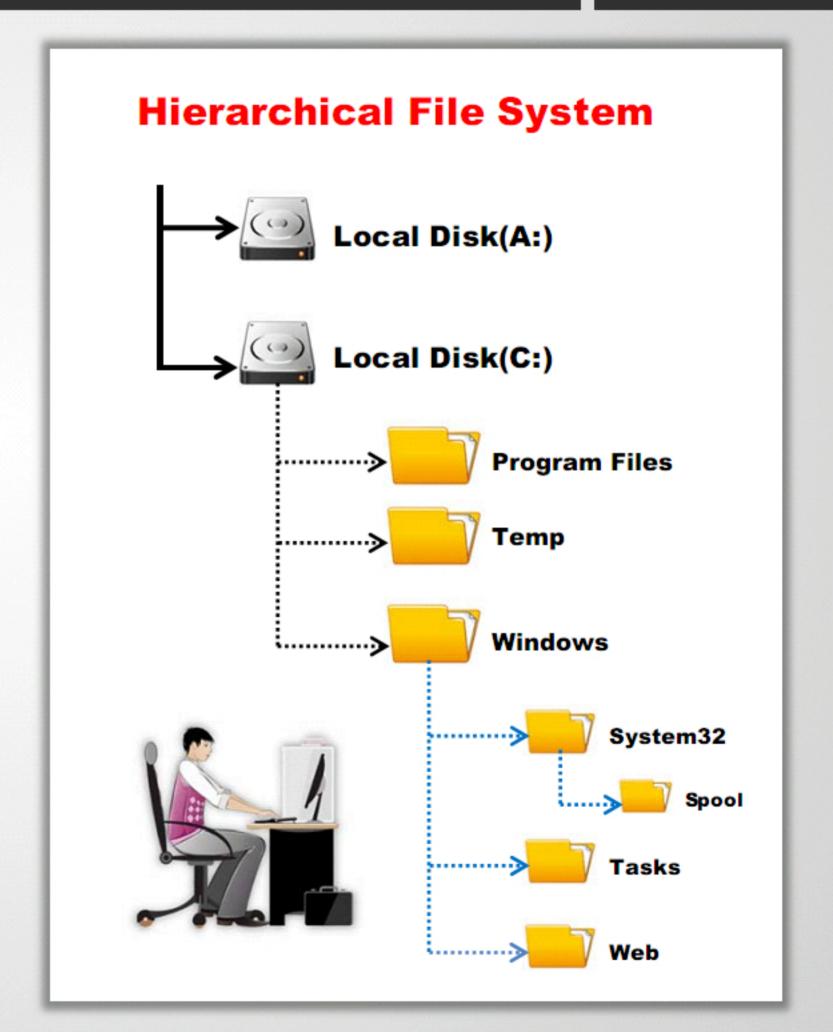


Hierarchical File System (HFS, also referred as Mac OS Standard) is a file system designed by **Apple** in 1985 for MAC operating system

It groups files into directories and each directory also groups with other directories

It displays drives, directories, and files in groups

It divides a logical volume into logical blocks of 512 bytes



# Hierarchical File System Plus (HFS+)



HFS+ is a successor of HFS and used as a primary file system in Macintosh



It supports large files and uses

Unicode for naming the items

(files and folders)



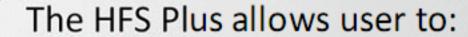








It is also called Mac OS Extended (HFS Extended) and is one of the formats used in the Apple iPod



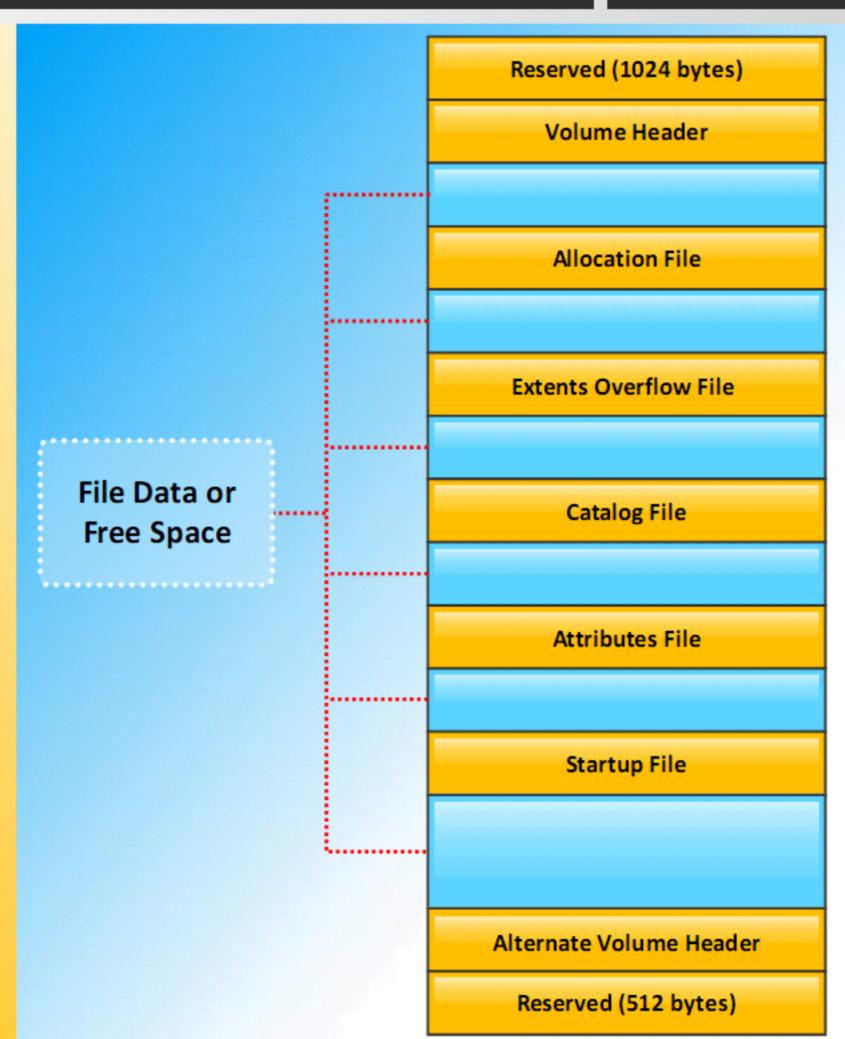
- Efficiently use hard disk space
- Use only international-friendly file names
- Easily boot on non-Mac OS operating systems

### HFS+ Volumes



- HFS+ volumes are divided into logical blocks (sectors) of 512 bytes
- These sectors are clustered into allocation blocks
- The total number of allocation blocks depends on the volume size
- The bulk of an HFS+ volume contains seven types of sectors:
  - User file fork
  - Allocation file
  - Catalog file
  - Extent overflow file
  - Attribute file
  - Startup file
  - Unused space





## HFS+ Journal

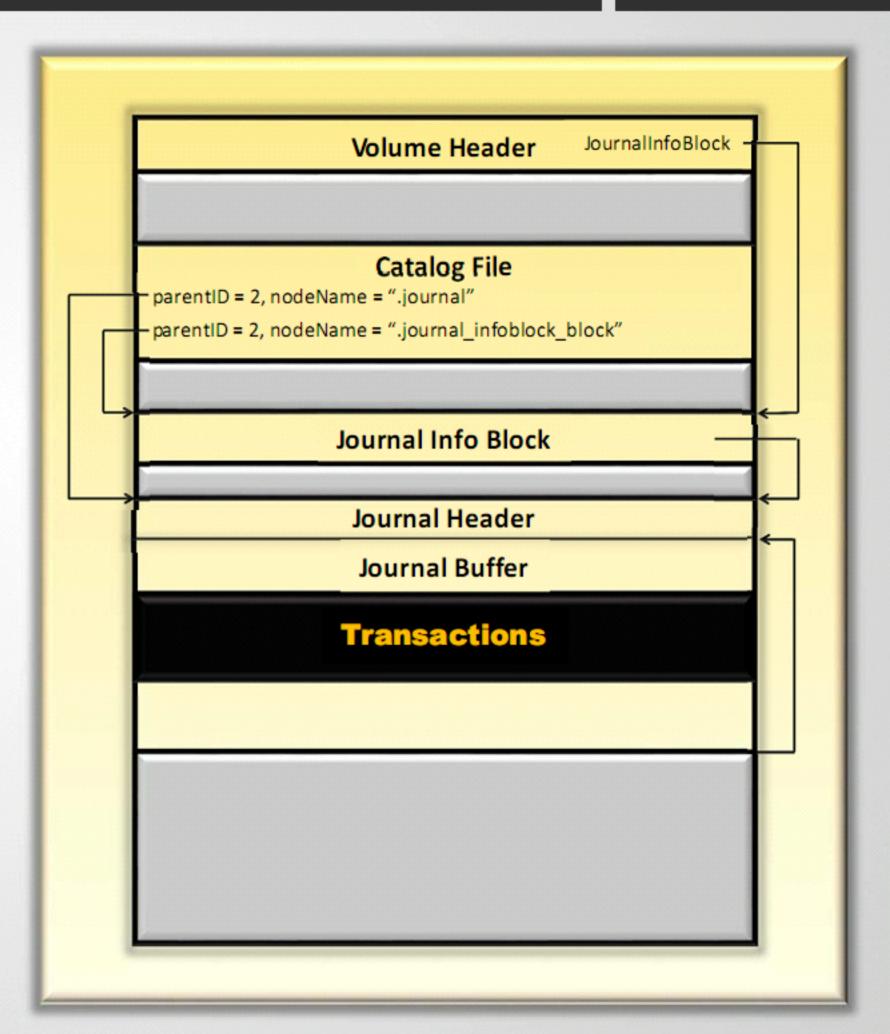


HFS+ volume has an optional journal which helps in mounting an unmounted volume in the case of a system crash

Journal **restores** the volume structures to a trustworthy state without scanning all of the structures

The journal info block (.journal\_info\_block) is stored as a file on the HFS+ volume's root directory





## Oracle Solaris 11 File System:



- ZFS is the default disk-based and root file system used in the Oracle Solaris 11
- It provides a simple management interface, which is robust, scalable, and easy to administer

### **Features:**

- ZFS Pooled Storage Model
- Data integrity Model
- Simplified Administration
- Copy-on-Write transactional model
- End-to-End Checksums
- Self-Healing Data
- Unparalled Scalability
- ZFS and Solid-State Storage
- Snapshots and Clones
- Encryption
- Deduplication
- Compression

### **Benefits:**

- Simplifies and reduces storage management tasks
- Increases storage agility and data protection
- Delivers superior performance and availability

http://www.oracle.com

## CD-ROM/DVD File System



The ISO (International Organization for Standardization) 9660 defines a file system for CD-ROM and DVD-ROM media



Windows supports two types of file systems on CD-ROM and Digital Versatile Disk (DVD):

- Compact Disc File System (CDFS)
- Universal Disk Format (UDF)

To exchange data, it supports various computer operating systems such as Microsoft Windows, Mac OS, and UNIX-based systems



ISO 13490 is a combination of ISO 9660 with multisession support



Common extensions to ISO 9660 were to deal with the limitations:

- Longer ASCII coded names and UNIX permissions are facilitated by Rock Ridge
- Unicode naming (like non-Roman scripts) are also supported by Joliet
- Bootable CDs are facilitated by El Torito



# Compact Disc File System (CDFS)



CD File System (CDFS) is a file system for Linux operating system



It transfers all **tracks** and **boot images** on a CD as normal files



It unlocks the information in old ISO images



For e.g., suppose multisession CD contains two ISO images, mounting the CD with CDFS file system, results in two sessions as files:

- [root@k6 /root] # mount -t cdfs -o ro /dev/cdrom /mnt/cdfs
- [root@k6 /root] # ls -l /mnt/cdfs
  total 33389
  -r--r-- 1 ronsse ronsse 33503232 Aug 8 19:36 sessions\_1-1.iso
  -r--r-- 1 ronsse ronsse 34121728 Aug 8 19:99 sessions 1-2.iso

## Virtual File System (VFS) and Universal Disk Format File System (UDF)



### Virtual File System (VFS)

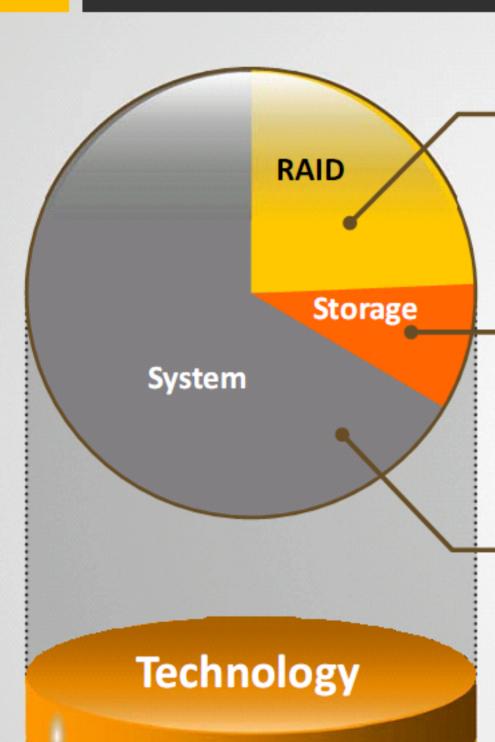
- A VFS is programming that specifies an interface between the OS's kernel and the different file systems
- VFS gives client applications access to the various concrete file system in a systematic manner
  E.g.: provision of transparent access to local and network storage devices without noticeable difference by the client application
- VMware Virtual Machine File System (VMFS), New Technology File System (NTFS), Global File System (GFS) and the Oracle Clustered File System (OCFS) are some of the VFS examples

### Universal Disk Format File System (UDF)

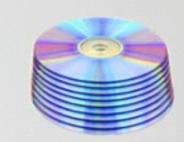
- UDF is a file system specification defined by the Optical Storage Technology Association (OSTA), aimed to replace the ISO9660 file system on optical media and also FAT on removable media
- It is an open source file system based on ISO/IEC 13346 and ECMA-167 standards that defines how data is stored and interchanged on a wide variety of optical media

## RAID Storage System





Redundant Array of Independent Disks (RAID) is a technology that uses multiple smaller disks simultaneously which function as a single large volume



It provides a particular method of accessing one or many separate hard disks, thereby decreasing the risk of losing all data in case of hard disk failures or damage, and improve access time



- This technology is developed to:
  - Maintain a large amount of data storage
  - Achieve a greater level of input/output performance
  - Achieve a greater reliability through data redundancy



## Levels of RAID Storage System



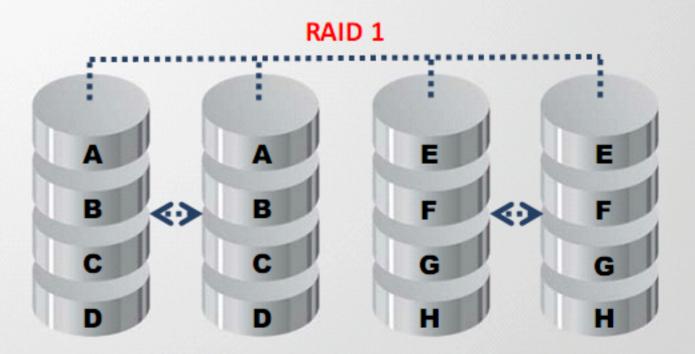
### RAID 0

- Data is split into blocks and written equally across multiple hard drives
- It improves I/O performance by spreading the I/O load across many channels and disk drives
- If any drive fails, data recovery is not possible
- It does not provide data redundancy
- It requires minimum two drives for setting up

# 

### RAID 1

- It is made up of two disks for each volume and is designed for data recovery in the event of disk failure
- The contents of the two disks are identical
- It ensures that data is not lost and assists in preventing computer downtime



# Levels of RAID Storage System (Cont'd)



### RAID 2

- It provides rapid access and increased storage by configuring two or more disks as one large volume, similar to RAID 1
- Data is written to a disk on a bit level

C0 to C3 = Word C; D0 to D3 = Word D

- Error correcting code (ECC) is used to verify whether the write is successful
- It has better data integrity checking and is slower than RAID 0

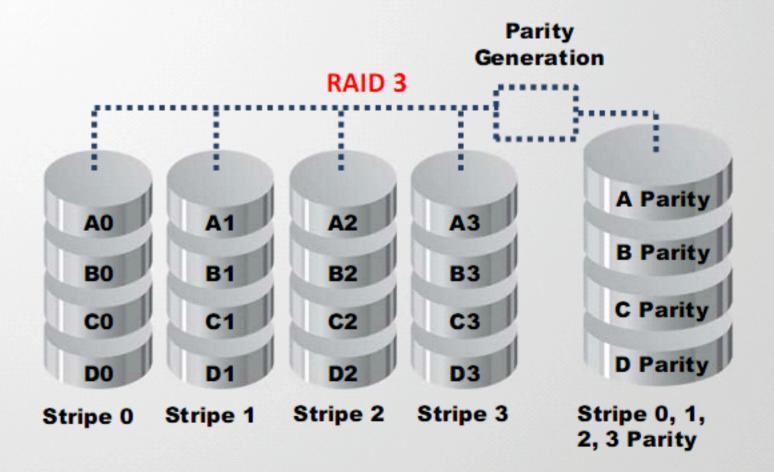
#### RAID 2 ECC/Ax ECC/Ay ECC/Az ECC/By ECC/Bz ECC/Bx **B3** BO **B1** CO ECC/Cx ECC/Cy C1 ECC/Cz ECC/Dy ECC/Dx ECC/Dz A0 to A3 = Word A; B0 to B3 = Word B;

ECC/Ax to Az = Word A ECC;

ECC/Bx to Bz = Word B ECC; ECC/Cx to Cz = Word C ECC; ECC/Dx to Dz = Word D ECC;

### RAID 3

- It uses data stripping and dedicated parity, and requires at least three disks
- Data is striped at a byte level across multiple drives and one drive is set to store parity information
- If any drive fails, data recovery and error correction is possible through the parity drive



# Levels of RAID Storage System (Cont'd)



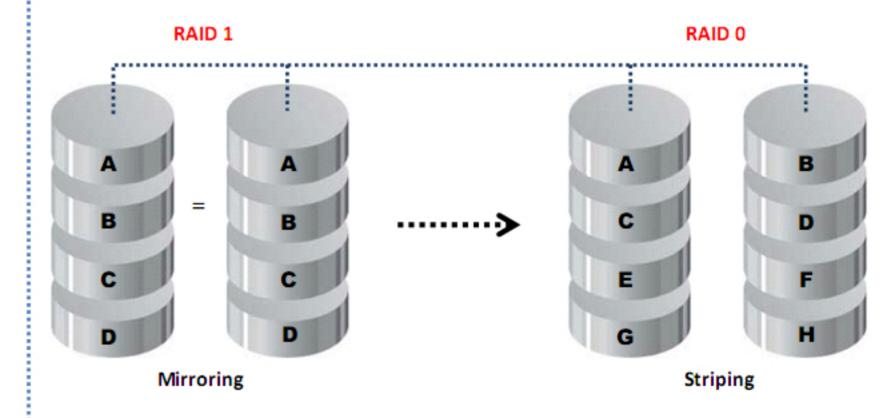
### RAID 5

- Data is striped at a byte level across multiple drives and parity information is distributed among all member drives
- Data writing process is slow
- It requires a minimum of three drives for setup

#### Generation B0 CO DO 0 Parity C1 1 Parity 1 Parity **B1** A1 2 Parity 2 Parity D2 **B2** A2 3 Parity 3 Parity A3 C3 D3 **D4** 4 Parity **B4 Blocks** Blocks **Blocks Blocks Blocks**

### **RAID 10 or Mirrored Striping**

- It is a combination of RAID 0 (Striping Volume Data) and RAID 1 (Disk Mirroring) and requires at least four drives to implement
- It has same fault tolerance as RAID level 1 and the same overheads as mirroring alone
- It allows mirroring of disks in pairs for redundancy and improved performance, and then data is striped across multiple disks for maximum performance



## Host Protected Areas (HPA) and Device Configuration Overlays (DCO)



Host Protected Areas (HPA) and Device Configuration Overlays (DCO) are the hidden areas of a hard disk

#### HPA:

- HPA is the reserved area on a HDD, meant to store data in a way that the user, BIOS, or OS cannot modify, change, or access it
- Information about HDD utilities, diagnostic tools, boot sector code, etc. is available in this area

### DCO:

- DCO is an additional hidden area available on modern hard disks, which enables system vendors to buy HDDs of varying sizes from different manufacturers and configure all of them to have equal number of sectors
- It can also be used to enable/disable features on the HDD
- With an intent to hide information, intruders use certain tools to modify and write to the HPA and DCO areas on the HDD
- HPA and DCO areas are of concern during the investigation as many tools fail to detect their presence
- Investigators can use tools such as EnCase, TAFT (an ATA (IDE) forensics tool), Sleuth Kit, etc. to detect and image HPA and/or DCO areas

## File System Analysis



### Understanding American Standard Code for Information Interchange (ASCII), Unicode, and Offset

- ASCII and Unicode are the two most common techniques used to encode the characters
- File contents can be viewed in ASCII or by extracting the ASCII strings from binary files for analysis

### **ASCII:**

- It is a character encoding scheme developed from telegraphic codes
- ASCII encodes 128 specified characters into 7-bit integers. The encoded characters are:
  - Numbers 0 to 9
  - Lowercase letters a to z
  - Uppercase letters A to Z
  - Basic punctuation symbols
  - Control codes that originated with Teletype machines
  - A space
- The ASCII table has 3 divisions namely, non printable (system codes between 0 and 31), lower ASCII (codes between 32 and 127) and higher ASCII (codes between 128 and 255)

## File System Analysis (Cont'd)



### UNICODE:

- 0
- It is an international encoding standard which supports consistent encoding, representation, and management of text expressed in many writing systems
- lt provides a unique number for every character, irrespective of the platform, program, and language
- 0
- UTF-8 and UTF-16 are the most widely used UTF character encodings

UTF Encoding	Description	
UTF-8	a 8-bit, variable-width encoding. Maximizes compatibility with ASCII	
UTF-16	a 16-bit, variable-width encoding	
UTF-32	a 32-bit, fixed-width encoding	

### OFFSET:



- In computing, an offset usually refers to either the start of a file or the start of a memory address
- Its value is added to a base address to derive the actual address
  - Example: If "A" denotes address 80, then the expression A+20 implies the address 100, where 20 in the expression is the offset

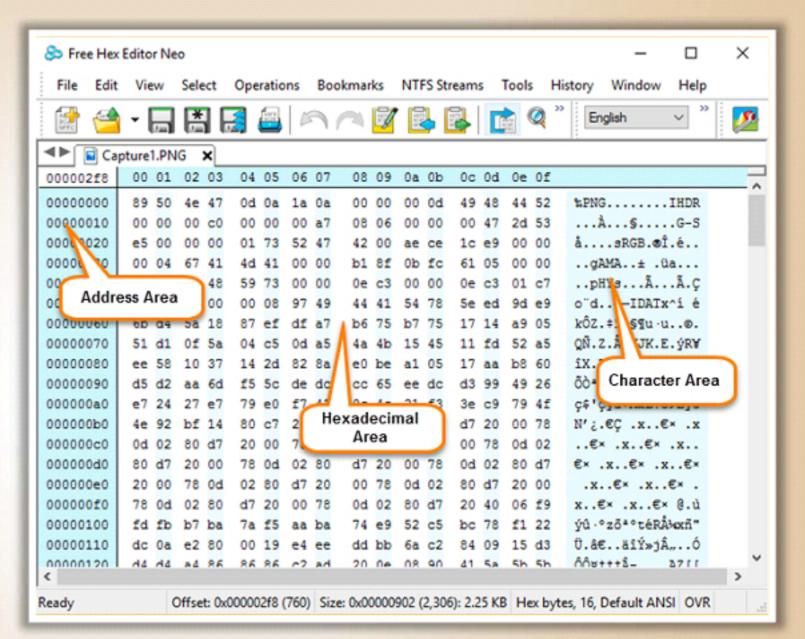


## File System Analysis (Cont'd)



### **Understanding Hex Editor:**

- Hex editors are the programs meant to examine or modify the physical (i.e. byte per byte) structure of a binary file
- Usually, hex-editors have three areas:
  - Address area: located on the left and exhibits the address of the first byte of each line usually in hexadecimal format
  - Hexadecimal area: located in the center, lists each byte of the file in a table, usually 16 bytes per line
  - Character area: located on the right, exhibits the ASCII representation of each of the bytes in the hexadecimal area
- In forensics, hex editors are of use to view stored or deleted data from both files and disk sectors
- In general, investigators use hex editors to examine evidence in particular parts of a disk
- Apart from the hexadecimal view of the data, many hex editors also display data both in binary and ASCII forms



http://www.hhdsoftware.com

## File System Analysis (Cont'd)



### **Understanding Hexadecimal Notation:**

- Hexadecimal numeral system, also known as hex, is a numeral system with base 16
- In hexadecimal notation 0–9 represent zero to nine values and English alphabets A, B, C, D, E, and F represent ten to fifteen values
  - E.g.: 2BA in hexadecimal is the same as 0010 1011 1010 binary
- Hexadecimal notation allows using powers of 2 easily instead of writing the whole thing in binary

HEX	Binary	Base 10
0	0000	0
1	0001	1
2	0010	2
3	0011	3
4	0100	4
5	0101	5
6	0110	6
7	0111	7
8	1000	8
9	1001	9
Α	1010	10
В	1011	11
С	1100	12
D	1101	13
Е	1110	14
F	1111	15

## File Carving



- It is a technique to recover files and fragments of files from unallocated space of the hard disk in the absence of file metadata
- In this technique, file identification and extraction is based on certain characteristics like file header or footer rather than the file extension or metadata
- A file header is a signature (also known as a magic number) which is a constant numeric or text value that determines a file format

**Example:** A suspect may try to hide an image from being detected by investigators by changing the file extension from .jpg to .dll

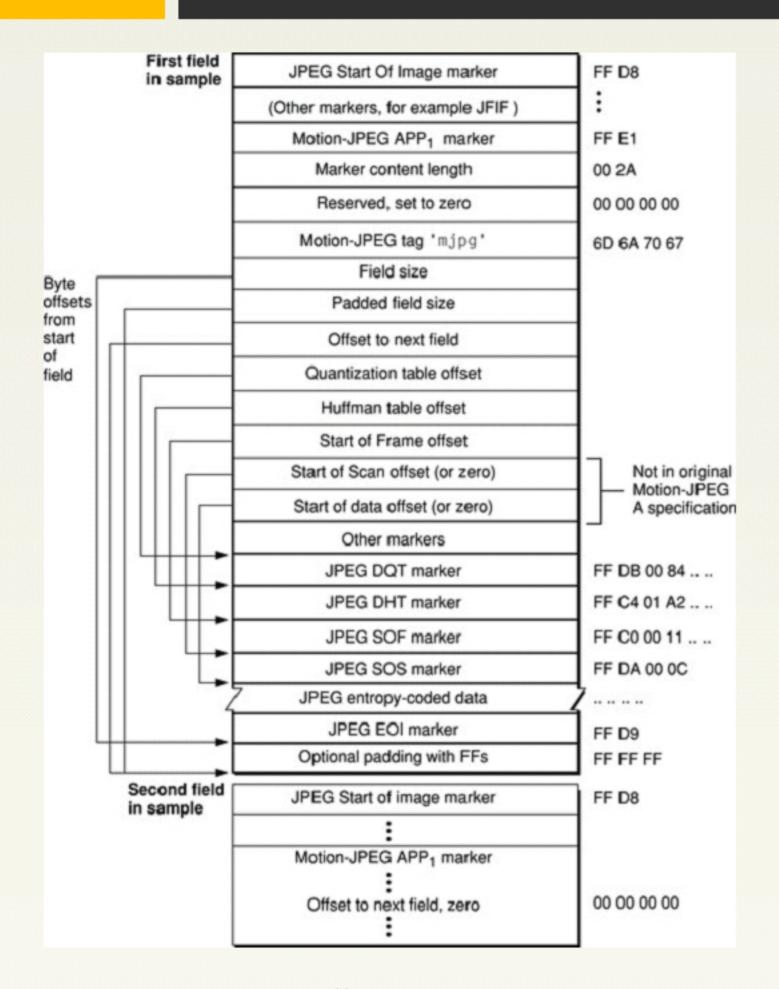
However, changing the file extension does not change to the file header, and analysis tells the actual file format

**Example:** A file format is confirmed as .jpg if it shows "JFIF" in the file header and hex signature as "4A 46 49 46"

Investigators can take a look at file headers to verify the file format using tools such as 010 Editor, CI Hex Viewer, Hexinator, Hex Editor Neo, Qiew, WinHex, etc.

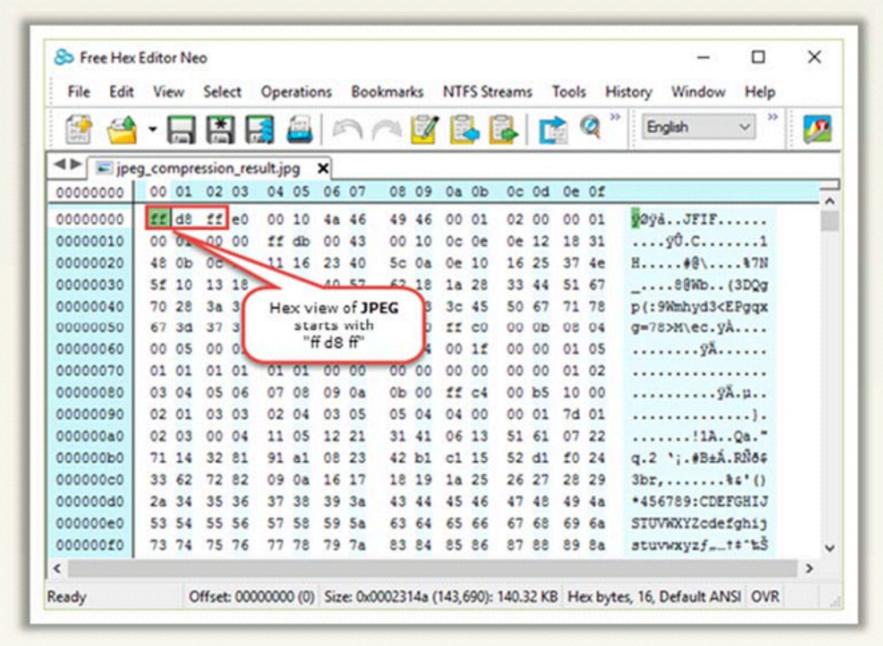
## Image File Formats: JPEG





https://developer.apple.com

- The JPEG is a commonly used method to compress photographic images
- It uses a compression algorithm to minimize the size of the file without affecting the quality of the image



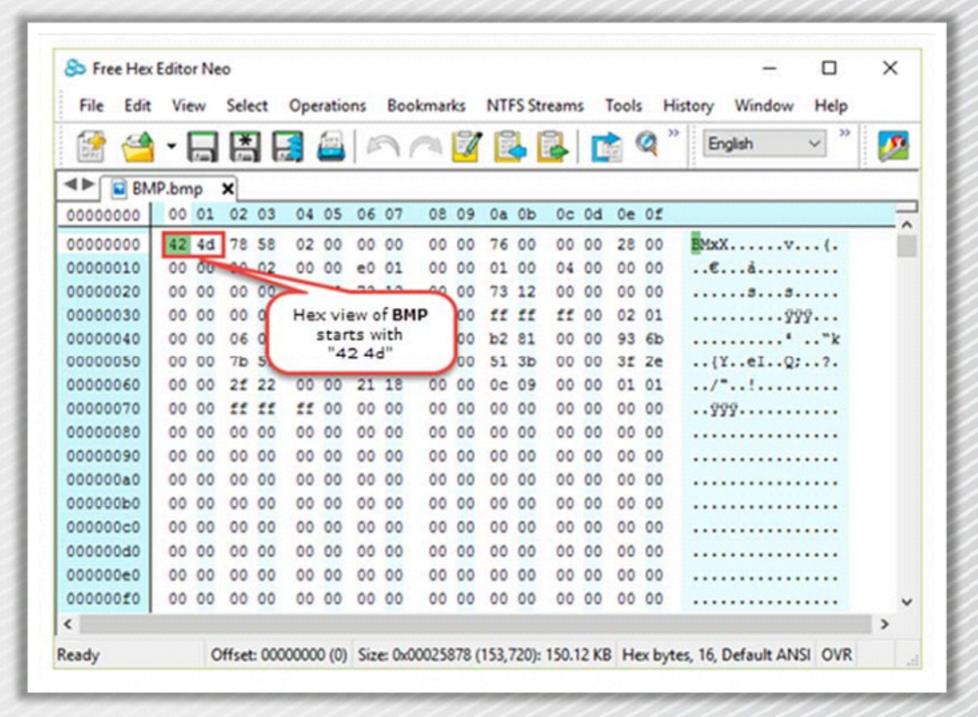
**Hex View of JPEG File Format** 

### Image File Formats: BMP



Basic BMP File Format					
N	ame	Size			
Header		14 bytes			
Signa	iture	2 bytes			
File Size		4 bytes			
reserved		4 bytes			
DataOffset		4 bytes			
Color Table		4 * NumColors bytes			
	Red	1 byte			
	Green	1 byte			
	Blue	1 byte			
	reserved	1 byte			
Repe	ated Num(	Colors times			
Raster Da	ta	Info ImageSize bytes			

- The Bitmap (BMP) is a standard file format for a Windows Device Independent Bitmap (DIB) file
- The size and color of these images can vary from 1 bit per pixel (black and white) to 24-bit color (16.7 million colors)



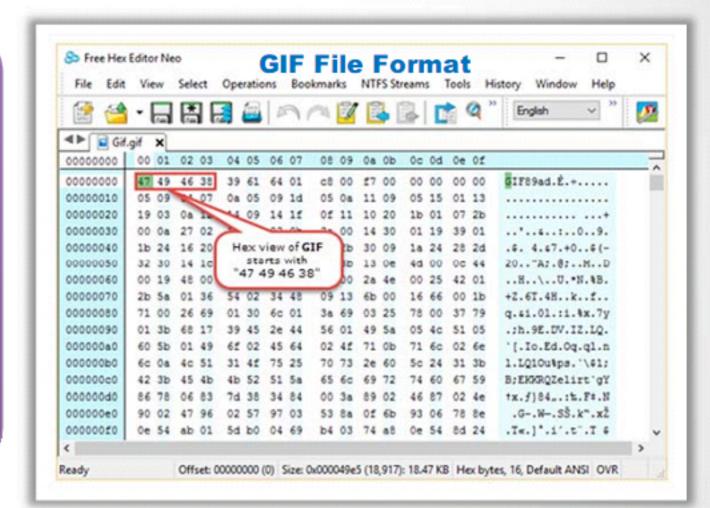
**Hex View of BMP File Format** 

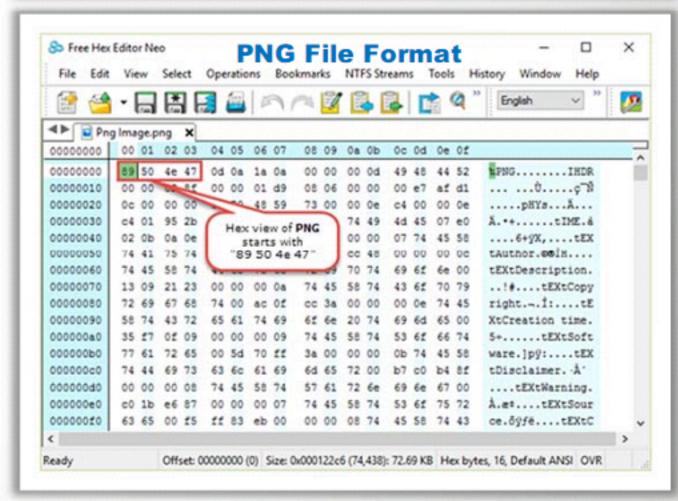
# Hex View of Popular Image File Formats



- The GIF is a file format that contains 8 bits per pixel and displays 256 colors per frame
- It uses lossless data compression techniques, which maintain the visual quality of the image

- The PNG is a lossless data compression image format, intended to replace the GIF and TIFF formats
- It supports:
  - Indexed / Palette-based images (24-bit RGB or 32-bit RGBA colors)
  - Grayscale images (with or without alpha channel)
  - Transparency (both normal and alpha channel)





### PDF File Format



- Attackers use PDF and Microsoft Office (Word, PowerPoint, and Excel) documents as attack vectors because of their wide usage by individuals and organizations
- Thus, it is essential for an investigator to understand the PDF and Microsoft Office file formats and structures, which may assist during malicious document analysis

#### PDF File Structure:

#### Header:

The first line of the PDF file specifies the version number of a PDF file format

Ex: %PDF-1.3

### Body:

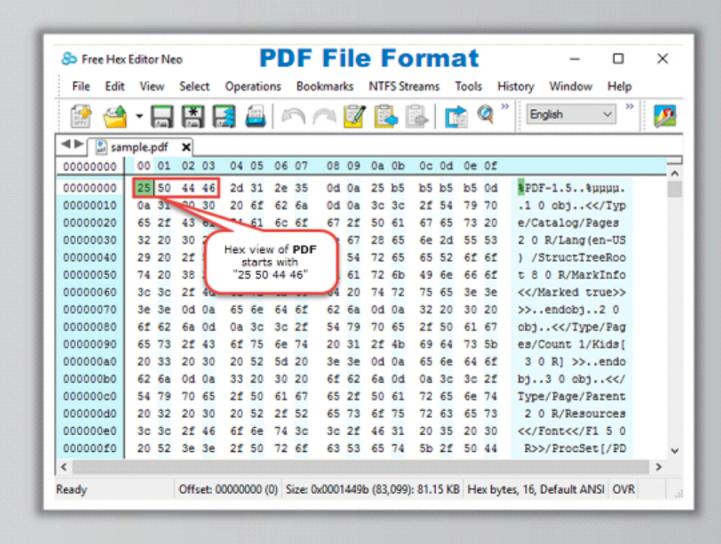
Consists of objects (images, fonts, form fields, text streams, annotations, bookmarks, etc.) that constitute the contents of the document

### The Cross-reference table (xref table):

- Contains links to all the objects in a file
- Allows random access to objects
- Allows to trace updated changes made to the PDF file

#### The Trailer:

- Contains links to the xref table and to main objects in the trailer dictionary
- Ends with %%EOF to recognize end of file



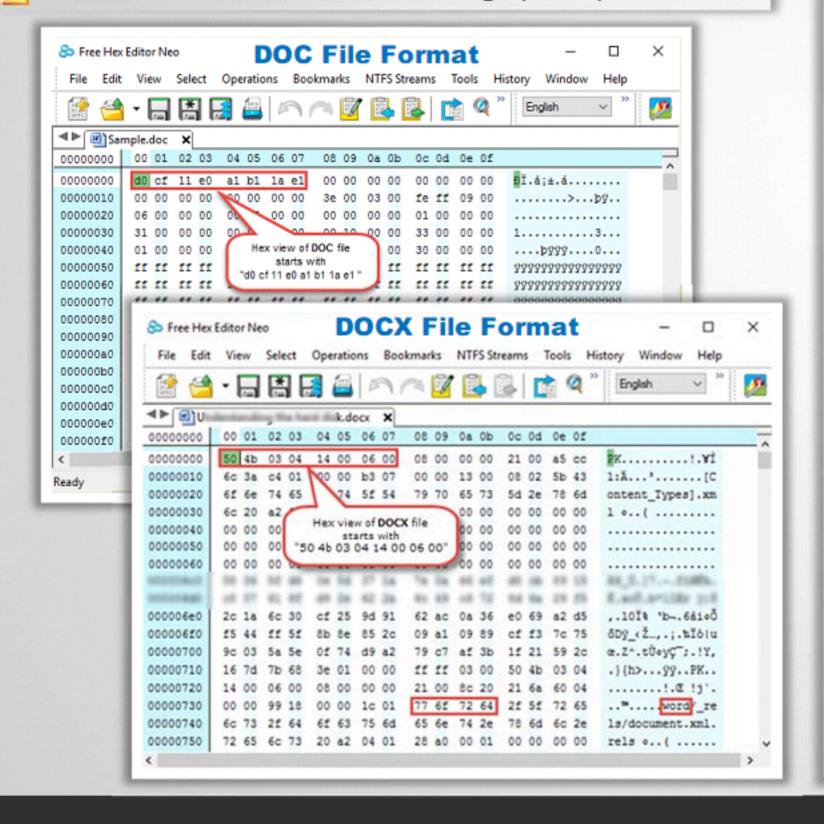
### Word File Formats



### MS Office Documents - File Format:

### **Binary File Format**

Usually, attackers target MS Office documents of the binary format because they are still in use and the file structures are highly complex



### Microsoft Word File Structure (.doc/.docx):

### WordDocument Stream/Main Stream

- Contains binary data of the Word document and Word file header (known as the File Information Block (FIB)) located at the offset 0
- FIB contains information about the document, file length, and specifies pointers to elements in the document file

### **Summary Information Streams**

The summary information is stored in two storage streams: Summary Information and DocumentSummaryInformation

### Table Stream (OTable or 1Table)

- Contains data referenced from the FIB and other parts of the file
- stores various plex of character positions (PLCs) and tables defining document's structure
- Has predefined structure only for encrypted files

### Data Stream (Optional)

- No predefined structure
- Contains data referenced from the FIB in the mainstream or other parts of the file

### **Object Streams**

Holds binary data for embedded OLE objects within the .doc file

### **PowerPoint File Formats**



## Microsoft PowerPoint Presentation File Structure (.ppt/.pptx):

### **Current User Stream**

Maintains CurrentUserAtom record, that identifies the last user's name to open/modify a target PPT and location of the most recent user edit

### **PowerPoint Document Stream**

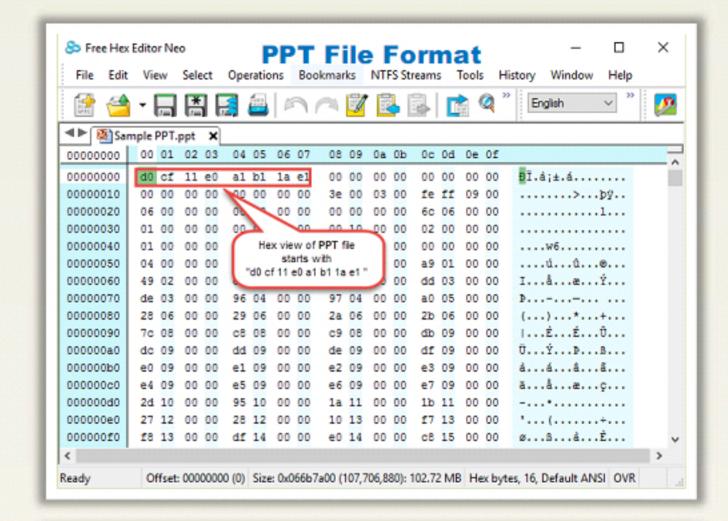
Contains information about the presentation layout and its contents

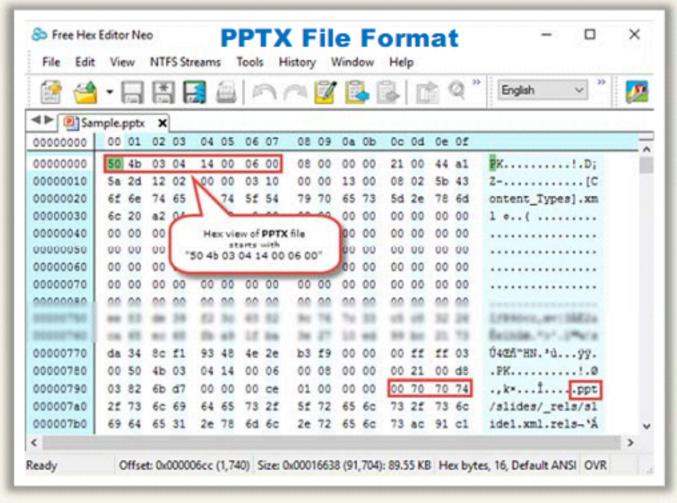
### **Pictures Stream (Optional)**

 Contains information about embedded image files within the presentation

### **Summary Information Streams (Optional)**

The summary information is stored in two storage streams: SummaryInformation and DocumentSummaryInformation





### **Excel File Formats**



## Microsoft Excel File Structure (.xls/.xlsx):

An OLE compound file saved in Binary Interchange File Format (BIFF)

#### Streams

Workbook stream is the primary stream in .xls file, which contains many substreams

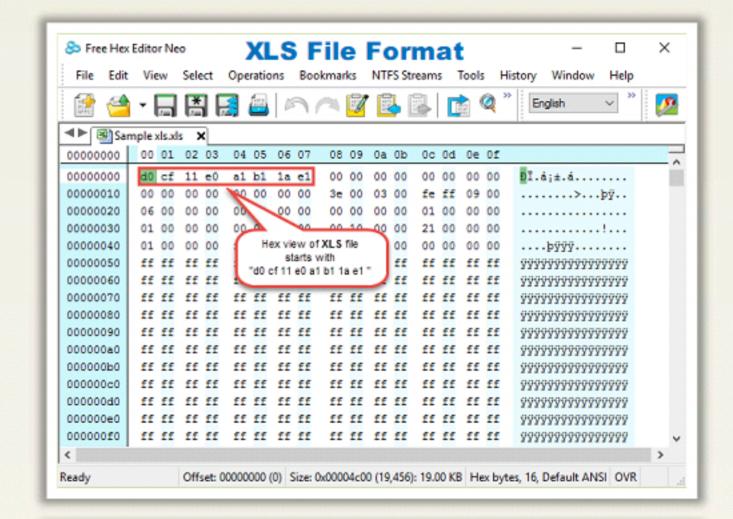
#### Substreams

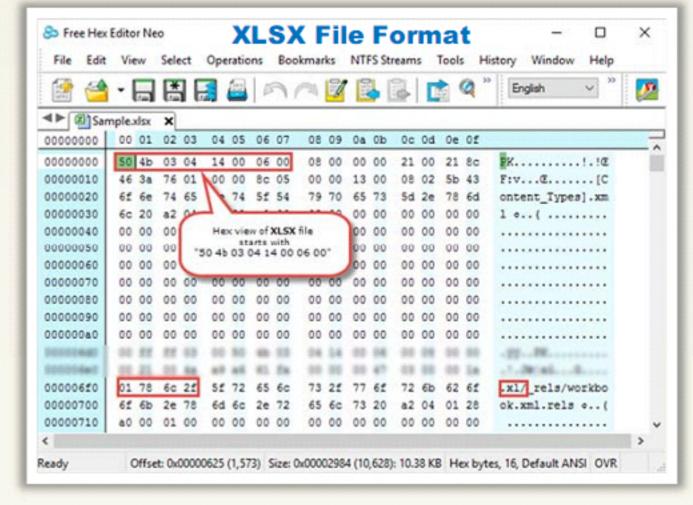
- Global substream specifies global properties and data in a workbook
- Worksheet substream specifies a sheet in a workbook

### Records

- Holds information about each workbook's features
- Components include record size, record type, and record data

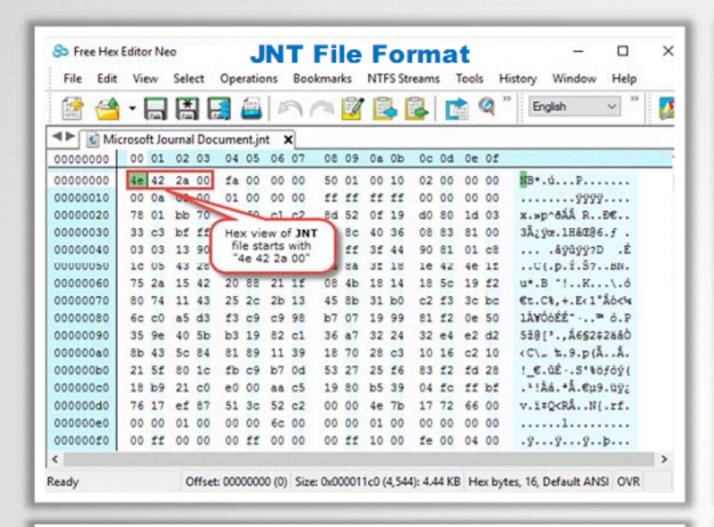
**Note:** Office Open XML Format (MS Office 2007 and above) is less vulnerable compared to the binary format and is therefore not widely used by attackers as a vector of attack

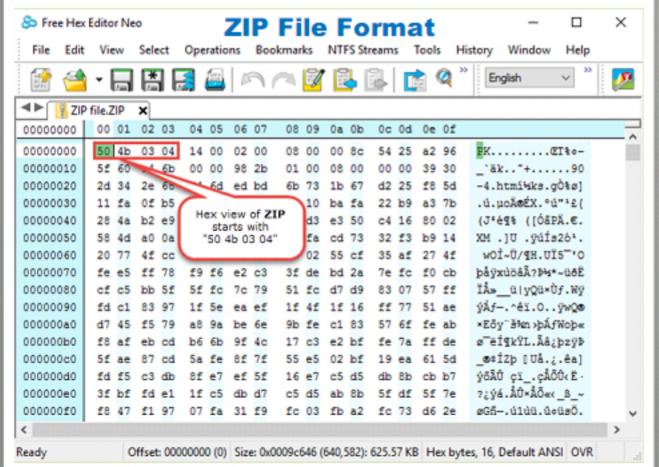


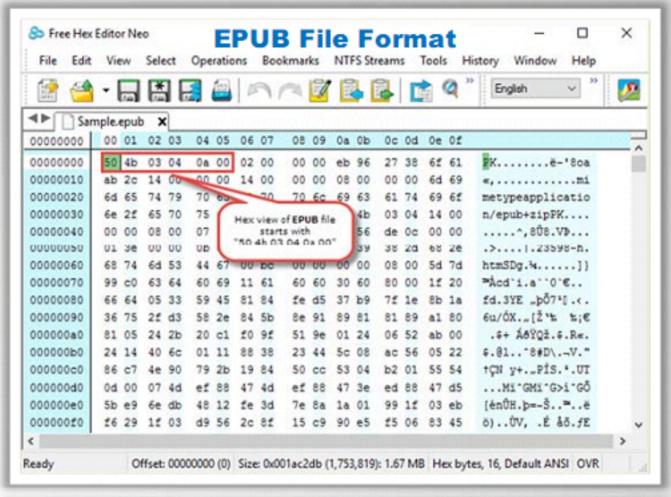


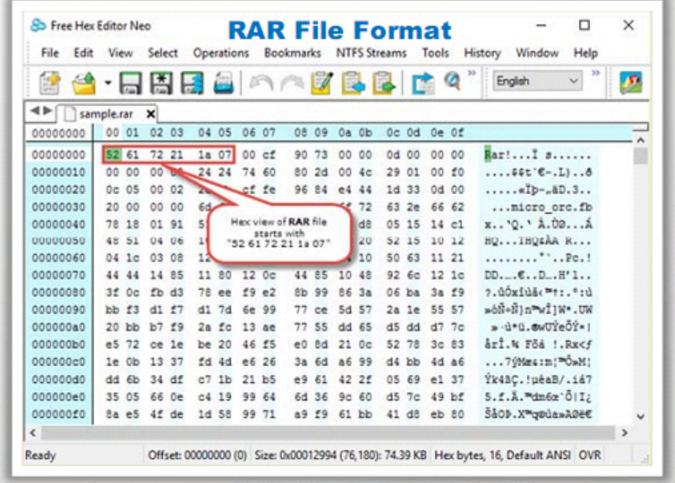
# Hex View of Other Popular File Formats





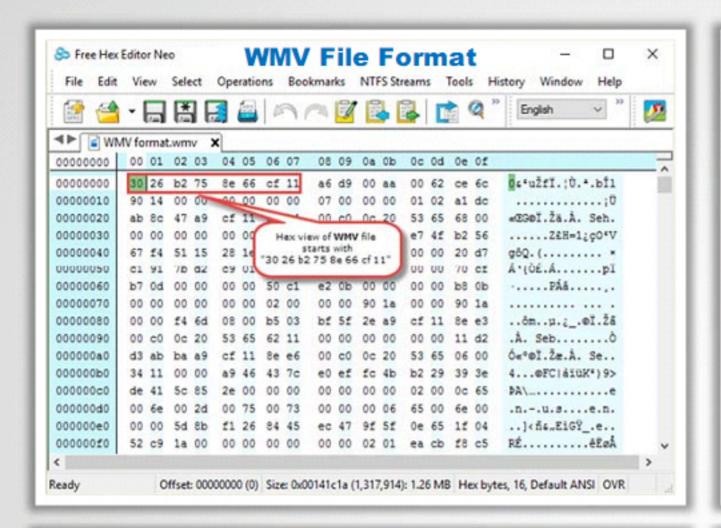


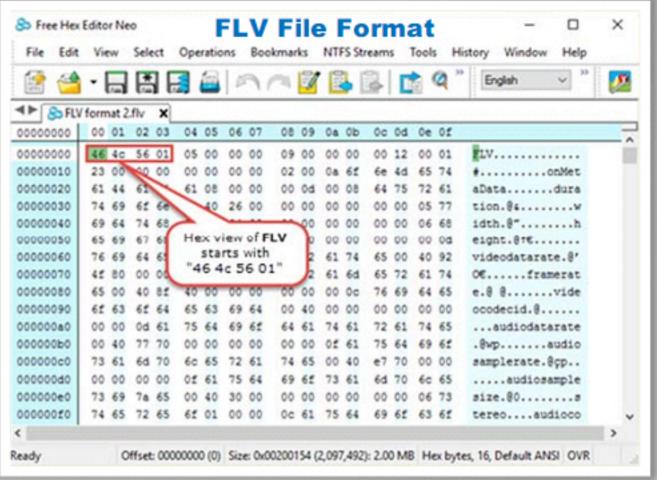


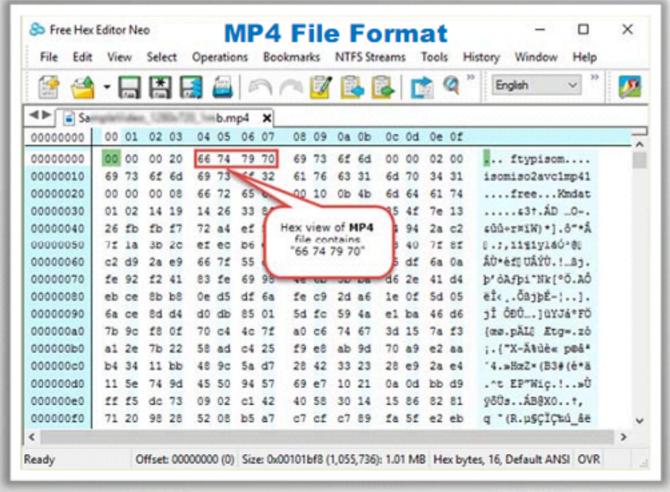


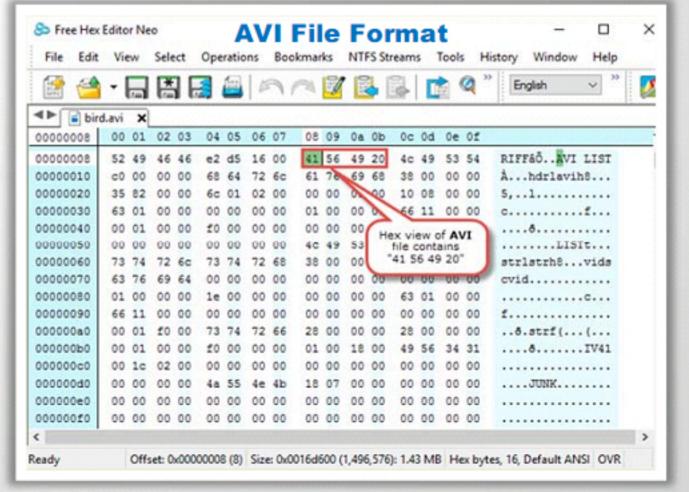
# Hex View of Popular Video File Formats





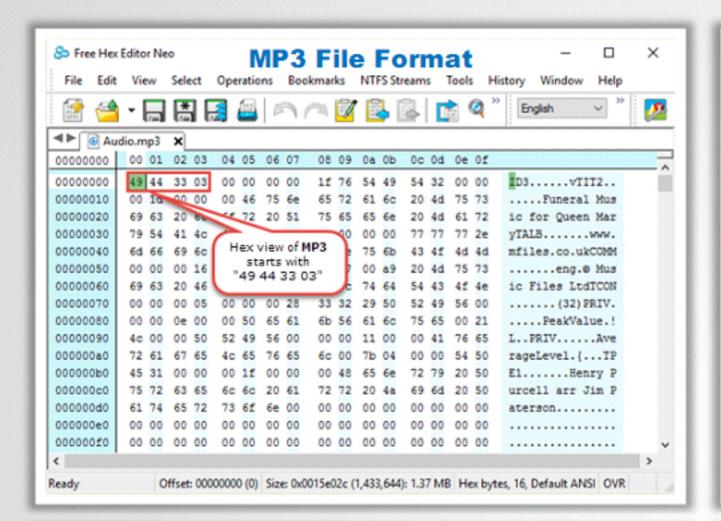


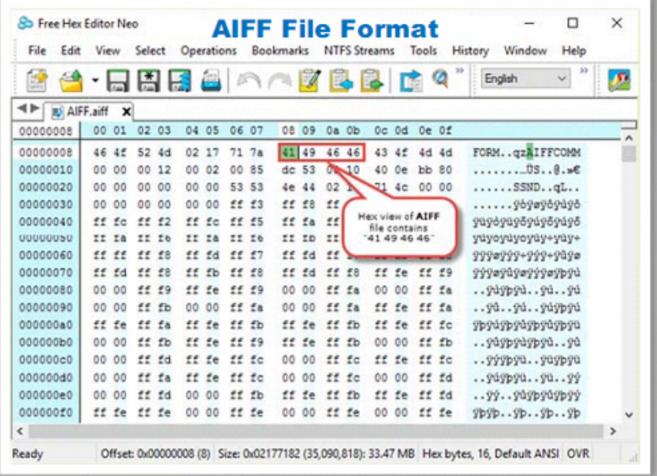


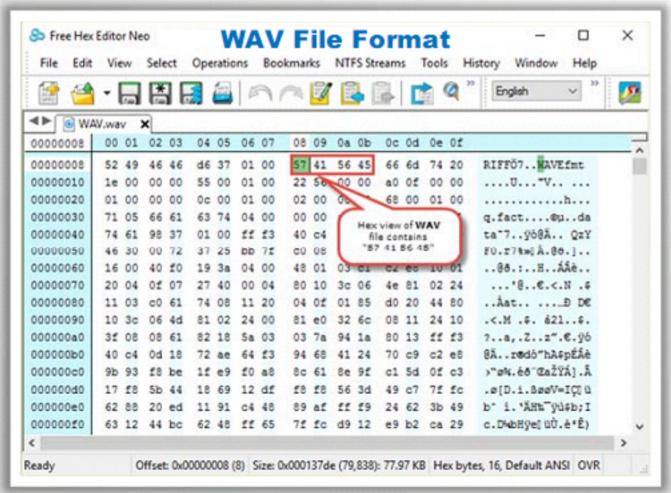


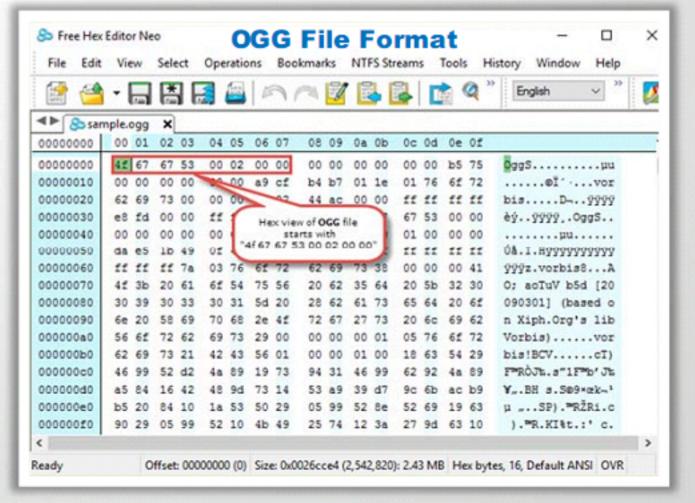
# Hex View of Popular Audio File Formats







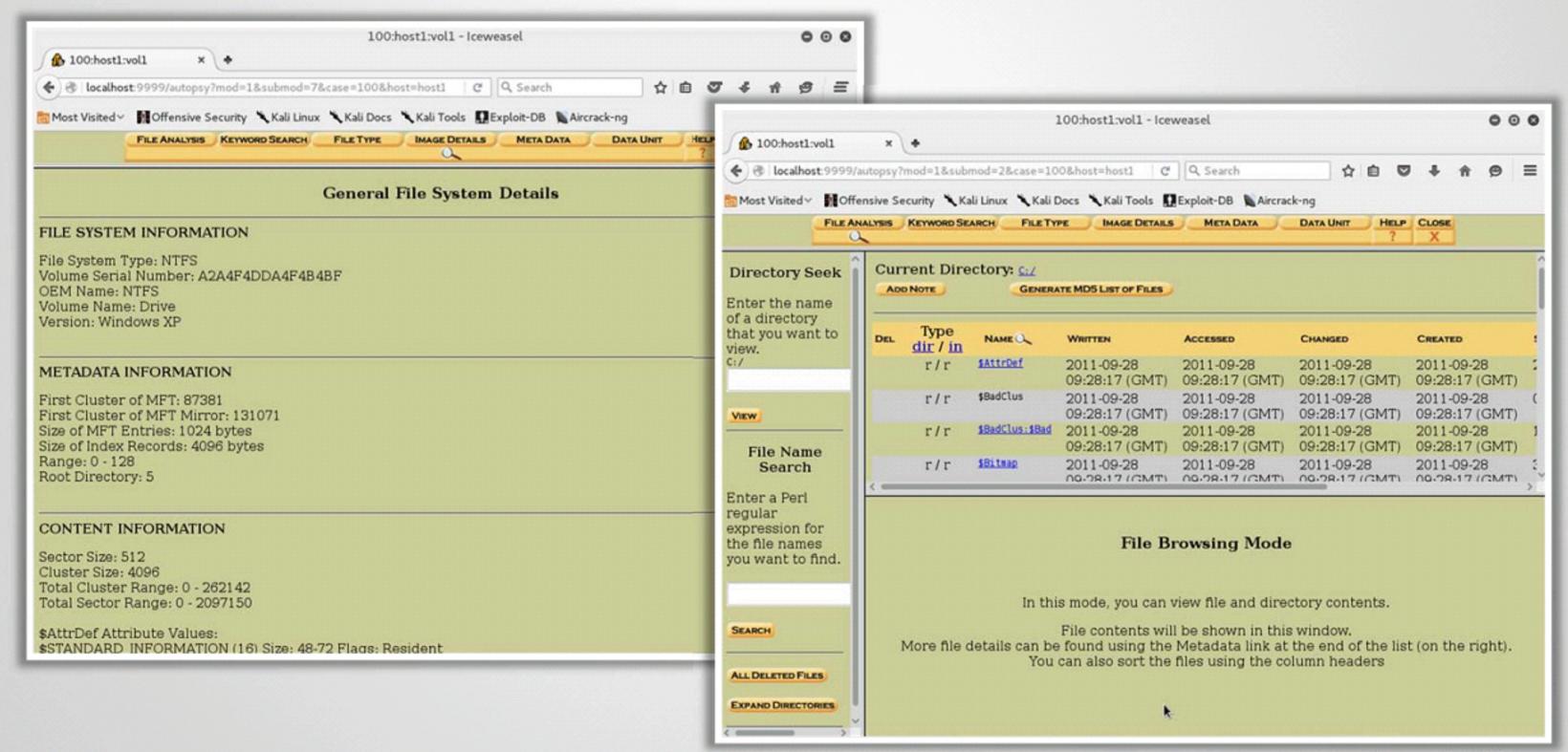




### File System Analysis Using Autopsy



Autopsy is a digital forensics platform and graphical interface to The Sleuth Kit and other digital forensics tools that can be used to investigate activities which happened on a computer



http://www.sleuthkit.org

# File System Analysis Using The Sleuth Kit (TSK)



The Sleuth Kit (TSK) is a library and a collection of command line tools that allow to investigate volume and file system data



The file system tools allow to examine file systems of a suspect computer in a non-intrusive fashion

5



The volume system (media management) tools allow to examine the layout of disks and other media



It supports DOS partitions, BSD partitions (disk labels), Mac partitions, Sun slices (Volume Table of Contents), and GPT disks



It analyzes raw (i.e. dd), Expert Witness (i.e. EnCase) and AFF file systems and disk images



It supports the NTFS, FAT, ExFAT, UFS 1, UFS 2, EXT2FS, EXT3FS, EXT4, HFS, ISO 9660, and YAFFS2 file systems



http://www.sleuthkit.org

**Note:** To perform analysis, create a forensics image .dd or .E01 of hard disk or pen drive using disk imaging tools. Here, we have created forensics image of a pen drive (.E01 format) using AccessData FTK Imager.

## The Sleuth Kit (TSK): fsstat







```
C:\WINDOWS\system32\cmd.exe
C:\Users\user\Downloads\sleuthkit-4.1.3-win32\sleuthkit-4.1.3-win32\bin>
fsstat -f ntfs "C:\Users\user\image.E01"
FILE SYSTEM INFORMATION
File System Type: NTFS
Volume Serial Number: B034FDE334FDAD0A
OEM Name: NTFS
Version: Windows XP
METADATA INFORMATION
First Cluster of MFT: 786432
First Cluster of MFT Mirror: 2
Size of MFT Entries: 1024 bytes
Size of Index Records: 4096 bytes
Range: 0 - 256
Root Directory: 5
CONTENT INFORMATION
Sector Size: 512
Cluster Size: 4096
Total Cluster Range: 0 - 3796986
Total Sector Range: 0 - 30375902
SAttrDef Attribute Values:
$STANDARD_INFORMATION (16) Size: 48-72 Flags: Resident
$ATTRIBUTE_LIST (32) Size: No Limit Flags: Non-resident
$FILE_NAME (48) Size: 68-578 Flags: Resident, Index
$OBJECT_ID (64) Size: 0-256 Flags: Resident
$SECURITY_DESCRIPTOR (80) Size: No Limit Flags: Non-resident
$VOLUME_NAME (96) Size: 2-256 Flags: Resident
$VOLUME_INFORMATION (112) Size: 12-12 Flags: Resident
SDATA (128) Size: No Limit Flags:
$INDEX_ROOT (144) Size: No Limit Flags: Resident
```

### The Sleuth Kit (TSK): istat (1 of 4)

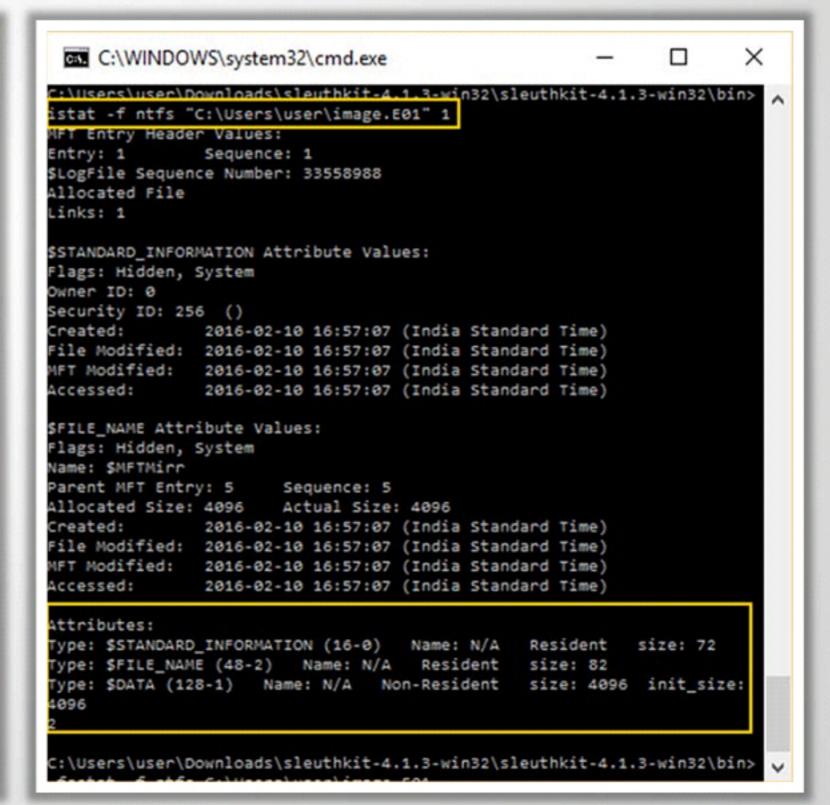


■ The istat tool in TSK shows the details of a directory entry and its output for a given entry

### **MFT File Overview**

#### C:\WINDOWS\system32\cmd.exe X C:\Users\user\Downloads\sleuthkit-4.1.3-win32\sleuthkit-4.1.3-win32\bin> istat -f ntfs "C:\Users\user\image.E01" 0 MFT Entry Header Values: Sequence: 1 \$LogFile Sequence Number: 33558918 Allocated File Links: 1 SSTANDARD\_INFORMATION Attribute Values: Flags: Hidden, System Owner ID: 0 Security ID: 256 () Created: 2016-02-10 16:57:07 (India Standard Time) File Modified: 2016-02-10 16:57:07 (India Standard Time) MFT Modified: 2016-02-10 16:57:07 (India Standard Time) Accessed: 2016-02-10 16:57:07 (India Standard Time) SFILE\_NAME Attribute Values: Flags: Hidden, System Name: SMFT Parent MFT Entry: 5 Allocated Size: 16384 Actual Size: 16384 2016-02-10 16:57:07 (India Standard Time) File Modified: 2016-02-10 16:57:07 (India Standard Time) NFT Modified: 2016-02-10 16:57:07 (India Standard Time) Accessed: 2016-02-10 16:57:07 (India Standard Time) Attributes: Type: \$STANDARD\_INFORMATION (16-0) Name: N/A Resident size: 72 Type: \$FILE\_NAME (48-3) Name: N/A Resident size: 262144 init\_size Type: \$DATA (128-6) Name: N/A Non-Resident 262144 786432 786433 786434 786435 786436 786437 786438 786439 786440 786441 786442 786443 786444 786445 786446 786447 786448 786449 786450 786451 786452 786453 786454 786455

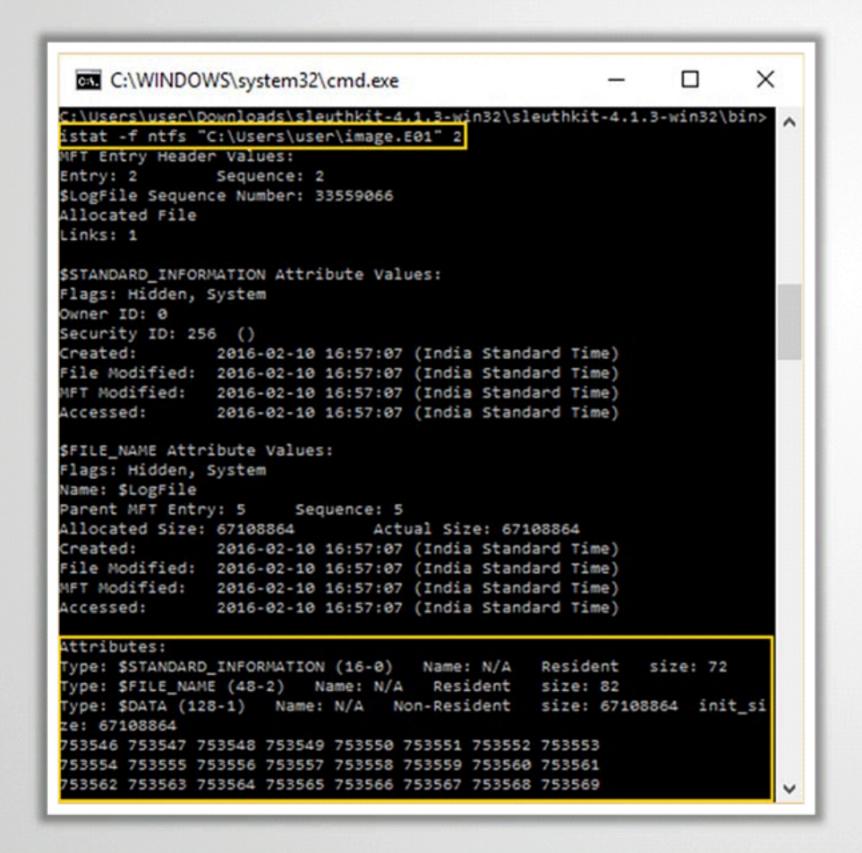
### **MFTMirr File Overview**



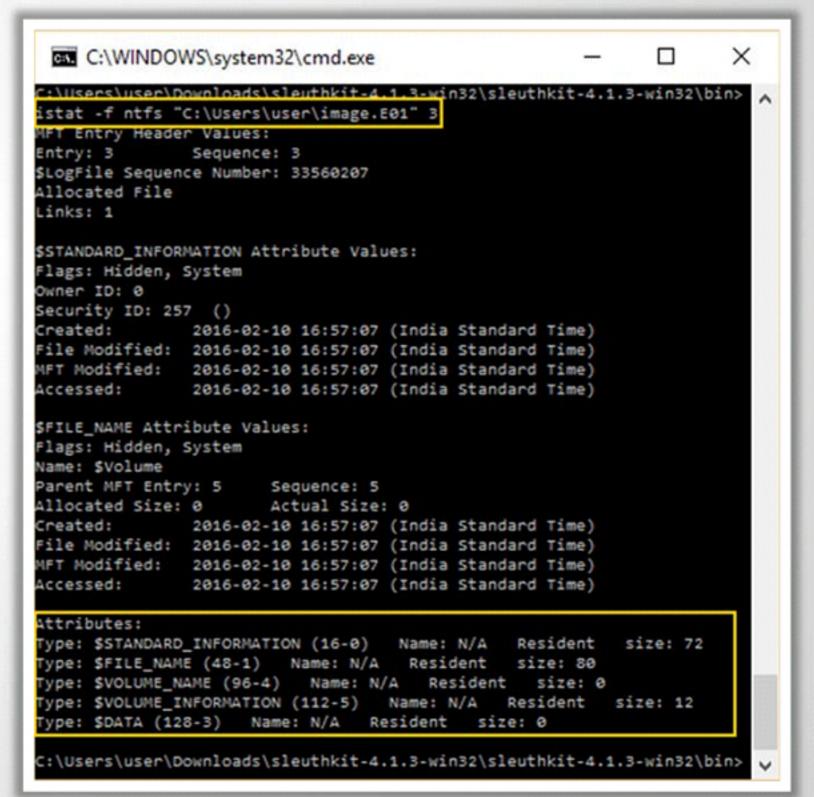
### The Sleuth Kit (TSK): istat (2 of 4)



### **LogFile Overview**



### **Volume File Overview**



### The Sleuth Kit (TSK): istat (3 of 4)



### **AttrDef File Overview**

#### C:\WINDOWS\system32\cmd.exe X C:\Users\user\Downloads\sleuthkit-4.1.3-win32\sleuthkit-4.1.3-win32\bin> istat -f ntfs "C:\Users\user\image.E01" 4 MFT Entry Header Values: Entry: 4 Sequence: 4 \$LogFile Sequence Number: 33560177 Allocated File Links: 1 \$STANDARD INFORMATION Attribute Values: Flags: Hidden, System Owner ID: 0 Security ID: 0 () Created: 2016-02-10 16:57:07 (India Standard Time) File Modified: 2016-02-10 16:57:07 (India Standard Time) MFT Modified: 2016-02-10 16:57:07 (India Standard Time) Accessed: 2016-02-10 16:57:07 (India Standard Time) SFILE NAME Attribute Values: Flags: Hidden, System Name: SAttrDef Parent MFT Entry: 5 Sequence: 5 Allocated Size: 36864 Actual Size: 36000 2016-02-10 16:57:07 (India Standard Time) File Modified: 2016-02-10 16:57:07 (India Standard Time) MFT Modified: 2016-02-10 16:57:07 (India Standard Time) Accessed: 2016-02-10 16:57:07 (India Standard Time) attributes: Type: \$STANDARD\_INFORMATION (16-0) Name: N/A Resident size: 48 Type: \$FILE\_NAME (48-2) Name: N/A Resident size: 82 Type: \$SECURITY\_DESCRIPTOR (80-3) Name: N/A Resident size: 104 Type: \$DATA (128-4) Name: N/A Non-Resident size: 2560 init\_size: 2560 769995

### **Bitmap File Overview**

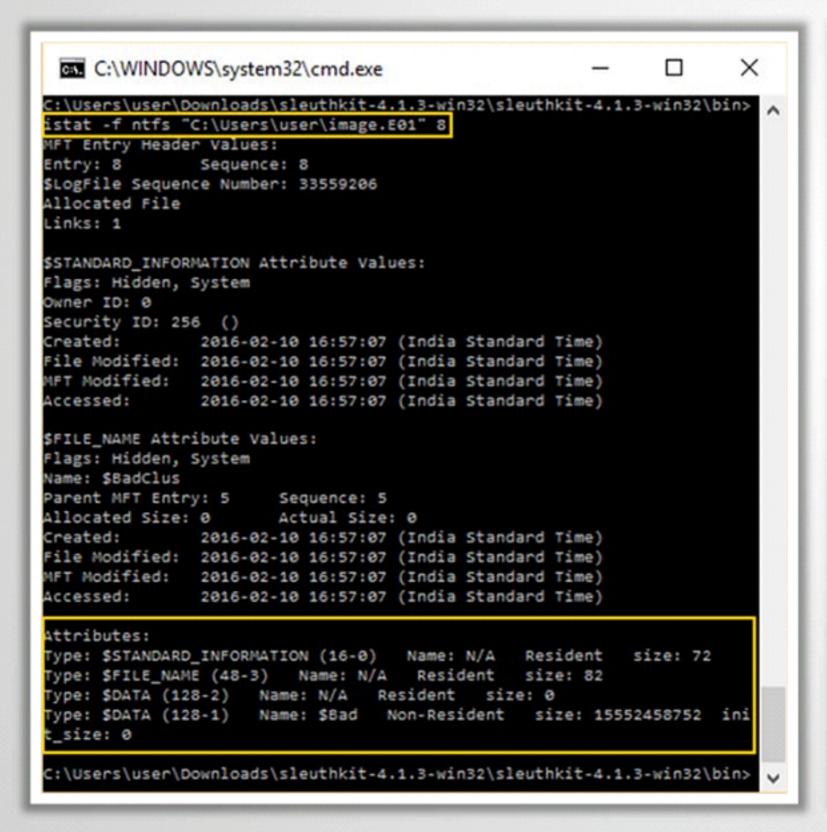
```
C:\WINDOWS\system32\cmd.exe
                                                                     X
C:\Users\user\Downloads\sleuthkit-4.1.3-win32\sleuthkit-4.1.3-win32\bin>
istat -f ntfs "C:\Users\user\image.E01" 6
MFT Entry Header Values:
Entry: 6
               Sequence: 6
$LogFile Sequence Number: 33559136
Allocated File
Links: 1
SSTANDARD INFORMATION Attribute Values:
Flags: Hidden, System
Owner ID: 0
Security ID: 256 ()
               2016-02-10 16:57:07 (India Standard Time)
Created:
File Modified: 2016-02-10 16:57:07 (India Standard Time)
MFT Modified: 2016-02-10 16:57:07 (India Standard Time)
               2016-02-10 16:57:07 (India Standard Time)
SFILE NAME Attribute Values:
Flags: Hidden, System
Name: $Bitmap
Parent MFT Entry: 5
                       Sequence: 5
Allocated Size: 475136
                               Actual Size: 474624
               2016-02-10 16:57:07 (India Standard Time)
File Modified: 2016-02-10 16:57:07 (India Standard Time)
MFT Modified: 2016-02-10 16:57:07 (India Standard Time)
Accessed:
               2016-02-10 16:57:07 (India Standard Time)
Attributes:
Type: $STANDARD_INFORMATION (16-0) Name: N/A
                                               Resident
Type: $FILE_NAME (48-2) Name: N/A Resident
                                               size: 80
Type: $DATA (128-4) Name: N/A Non-Resident size: 474624 init size
786314 786315 786316 786317 786318 786319 786320 786321
786322 786323 786324 786325 786326 786327 786328 786329
786330 786331 786332 786333 786334 786335 786336 786337
```

### The Sleuth Kit (TSK): istat (4 of 4)



### **BadClus File Overview**

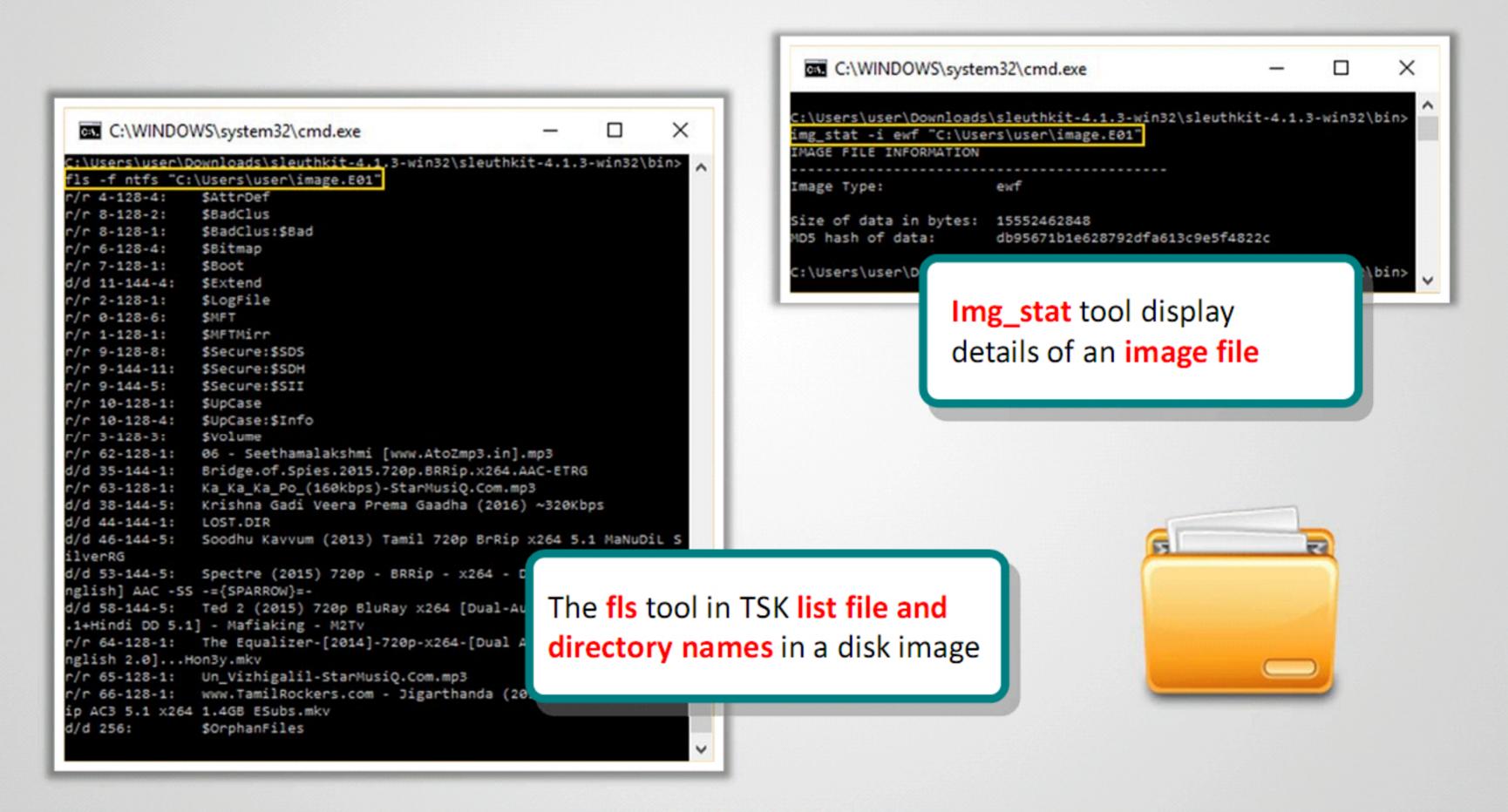
### Secure File Overview



```
C:\WINDOWS\system32\cmd.exe
                                                                     X
C:\Users\user\Downloads\sleuthkit-4.1.3-win32\sleuthkit-4.1.3-win32\bin>
istat -f ntfs "C:\Users\user\image.E01" 9
MFT Entry Header Values:
               Sequence: 9
$LogFile Sequence Number: 33563978
Allocated File
inks: 1
$STANDARD_INFORMATION Attribute Values:
Flags: Hidden, System
Owner ID: 0
Security ID: 257 ()
Created:
               2016-02-10 16:57:07 (India Standard Time)
File Modified: 2016-02-10 16:57:07 (India Standard Time)
MFT Modified: 2016-02-10 16:57:07 (India Standard Time)
Accessed:
               2016-02-10 16:57:07 (India Standard Time)
SFILE_NAME Attribute Values:
Flags: Hidden, System, Index View
Name: $Secure
Parent MFT Entry: 5
                       Sequence: 5
Allocated Size: 0
                       Actual Size: 0
               2016-02-10 16:57:07 (India Standard Time)
File Modified: 2016-02-10 16:57:07 (India Standard Time)
MFT Modified: 2016-02-10 16:57:07 (India Standard Time)
Accessed:
               2016-02-10 16:57:07 (India Standard Time)
Attributes:
Type: $STANDARD_INFORMATION (16-0) Name: N/A
                                                Resident
                                                          size: 72
Type: $FILE_NAME (48-7) Name: N/A Resident
Type: $DATA (128-8) Name: $SDS Non-Resident size: 263008 init_siz
263008
769930 769931 769932 769933 769934 769935 769936 769937
769938 769939 769940 769941 769942 769943 769944 769945
769946 769947 769948 769949 769950 769951 769952 769953
```

# The Sleuth Kit (TSK): fls and img\_stat





## The Sleuth Kit (TSK): fls and img\_stat



